STIC Search Report
EIC 2100

STIC Database Tracking Number: 181657

TO: Thanh-ha Dang
Location: RND 3B15
Art Unit: 2163
Thursday, March 09, 2006

Case Serial Number: 10/618409

From: Emory Damron
Location: EIC 2100
RND 4B19
Phone: 571-272-3520

Emory.Damron@uspto.gov

Search Notes

Dear Thanh-ha,

Please find below your fast and focused results.
References of potential pertinence have been tagged, but please review all the packets in case you like something I didn't.
Of those references which have been tagged, please note any manual highlighting which I've done within the document.
In addition to searching on Dialog, I also searched EPO/JPO/Derwent, IEEE and Inspec.
There may be a few decent references contained herein, but I'll let you determine how useful they may be to you.
Please contact me if I can refocus or expand any aspect of this case, and please take a moment to provide any feedback (on the form provided) so EIC 2100 may better serve your needs. Good Luck!

Sincerely,
Emory Damron
Technical Information Specialist
EIC 2100, US Patent & Trademark Office
Phone: (571) 272-3520
Emory.damron@uspto.gov

```
? d s
Set      Items    Description
S1       34073    S (QUERY? OR QUERIE?)()LANGUAG? OR SQL? ? OR (XML? ? OR EXTEN?()(MARKUP?
OR MARK?()UP)()LANGUAG?)()QUER?
S2         289    S HQL? ? OR QUERYLANGUAG? OR GQL? ?
S3        9652    S (INTERROG? OR RETRIEV? OR OBTAIN? OR REQUIR? OR QUESTION? OR ENQUIR? OR
KEYWORD? OR KEY()WORD?)()LANGUAG? OR FETCH?
S4     6910026    S LINK? OR JOIN? OR COMBINE? OR COMBINING? OR COMBINATION? OR COMBINATOR?
OR RELATIONSHIP?
S5     7990551    S AFFILIAT? OR CONJOIN? OR RELATED OR RELATION? OR BUILD? OR CONCATENAT?
OR COMPIL?
S6     4314563    S KINSHIP? OR CORRESPOND? OR CORRELAT? OR DRAG?(2W)DROP? OR DRAGANDDROP?
OR DRAGNDROP?
S7    11210576    S ADAPT? OR RECONFIGUR? OR CONFIGUR? OR CHANGE? OR CHANGING? OR MODIF? OR
ALTER? OR CONVERT? OR CONVERSION?
S8     3119003    S AMEND? OR UPDAT? OR REVIS? OR REARRANG? OR REORDER? OR RESHUFFL? OR
RESET? OR REFORMAT? OR TRANSFORM?
S9    13781657    S FIELD? OR ARRAY? OR MENU?.? OR DROPDOWN? OR DROP()DOWN? OR GUI? ? OR
(GRAPHIC? OR USER?)(2W)INTERFACE? OR UI OR MODEL?
S10      46506    S INTERFACE?(2N)(APPLICATION? OR SOFTWARE?) OR DIALOG()(BOX OR BOXES) OR
PULLDOWN? OR PULL?()DOWN? OR QAUI? ?
S11      91404    S API?·? OR (APP OR ·APPS OR APPLICATION?)(2N)(INTERFACE? OR SCREEN? OR
·WINDOW? OR PAGE? ? OR BOX OR CELL? ? OR TEMPLAT?)
S12      26980    S S1:S3 AND S4:S8
S13      17783    S S1:S3 AND S9:S11
S14      32526    S S12:S13
S15      12146    S USER? OR CLIENT? OR CUSTOMER? OR ENDUSER? OR NETIZEN? OR OPERATOR? OR
MANAGER?
S16        466    S ACCOUNT?(2N)HOLDER? OR PATRON? OR MEMBER? OR SUBSCRIBER? OR WEBUSER?
S17       1487    S PARTY? OR PERSON? ? OR INDIVIDUAL? OR PARTIE? OR PRINCIPAL?
S18       5276    S DESIRED? OR CHOICE? OR PREFER? OR CHOOS? OR SELECT? OR ELECT? OR OPT? ?
OR OPTING OR OPTED OR OPTION? OR PRECEDEN?
S19        696    S CUSTOMIZ? OR CUSTOMIS? OR PERSONALIS? OR PERSONALIZ? OR INDIVIDUALIZ? OR
INDIVIDUALIS?
S20       2505    S ELECT? OR PICK? OR DESIGNAT? OR DISCRIMINAT? OR ASSIGN? OR SPECIFY? OR
PRESET?
S21        769    S STATEMENT?
S22      10153    S STRING? OR SET? ? OR SEQUENC? OR BYTE? OR BIT? ? OR CHARACTERS OR LINE?
? OR STREAM? OR PREDICATE?
S23       5175    S CHUNK? ? OR TEXT? OR SEGMENT? OR SENTENC? OR SCRIPT? OR EXPRESSION? OR
QUEUE? OR CLAUSE? OR DECLARATION?
S24      13964    S LOGICAL? OR ABSTRACT? OR OBJECT? OR DATAOBJECT? OR METADATA? OR
META()DATA? OR VIRTUAL? OR HTML(2N)DATA? OR HTMLOBJECT? OR CONCEPTUAL?
S25       4555    S PHYSICAL? OR REAL? OR EXECUTAB? OR TANGIB? OR RUNNABLE? OR (RUN? ? OR
LAUNCH? OR START?)(2N)CAPAB? OR LAUNCHAB?
S26        313    S USABL? OR UTILIZAB? OR UTILISAB? OR UTILE OR CONCRETE?
S27      12237    S S14 AND S12 AND S13
S28       1061    S S14 AND S15:S17 AND S18:S20 AND S21:S23 AND S24:S26
S29        529    S S27 AND S28
S30         92    S S29 AND S24 AND S25:S26
S31        147    S S28 AND S24 AND S25:S26
S32        147    S S30:S31
S33        120    S S32 AND PY<2004
S34        122    S S32 NOT PY>2003
S35        122    S S33:S34
S36        117    RD (unique items)
; show files
```

[File 2] **INSPEC** 1898-2006/Feb W4

[File 6] **NTIS** 1964-2006/Feb W3

[File 8] **Ei Compendex(R)** 1970-2006/Feb W4

[File 34] **SciSearch(R) Cited Ref Sci** 1990-2006/Feb W4

[File 35] **Dissertation Abs Online** 1861-2006/Feb

[File 65] **Inside Conferences** 1993-2006/Mar 08

[File 94] **JICST-EPlus** 1985-2006/Dec W2

[File 99] **Wilson Appl. Sci & Tech Abs** 1983-2006/Feb

[File 111] **TGG Natl.Newspaper Index(SM)** 1979-2006/Mar 01

[File 144] **Pascal** 1973-2006/Feb W2

[File 239] **Mathsci** 1940-2006/Apr

[File 256] **TecInfoSource 82-2006/Feb** (c) 2006 Info.Sources Inc

01458959   ORDER NO: AADAA-I1375595
ENHANCED QUERY FORMULATION AND DATABASE MANAGEMENT TOOL DESIGN

Author: BADHE, HARISH
Degree: M.S.C.S.
Year: 1994
Corporate Source/Institution: UNIVERSITY OF LOUISVILLE (0110)
Source: Volume 34/01 of MASTERS ABSTRACTS. of Dissertations Abstracts International.
PAGE 337 .
Year: 1994

...uses. A database requires the ability to manage the various entities associated with it (viz. **users**, accounts, **physical** and **logical** space, privileges, structure etc.). All these actions are effected by **user**-actions generically termed as query **statements**. The act of querying assumes particular importance in **relational** databases which advocate a universal **relationship** concept, since the **user** has to **build** the **relations** each time he/she queries the database. The complexity of each query depends on the underlying data structure and the number of **relations** the **user** wants to **specify**. This thesis developes a **graphic interface** tool which replaces the bare command-**line** interface provided by most database servers can aid in better query formulation and focuses the **user's** attention on problems other than remembering the syntax of fourth generation **query languages** like Structured **Query Languages(SQL)**. The test-bed chosen for developing this interface is ORACLE 6.0. Also keeping in touch with the latest trend in graphics development, Visual Basic 3.0 was used to **build** the **user-interface**. The front-end .query-tool is designed to accommodate **users** at extreme ends of the learning curve as well as everyone in between.

01375367   ORDER NO: AAD94-26045
A FORMAL DEFINITION, QUERY LANGUAGE AND RELATIONAL DATABASE MAPPING FOR THE SEMANTIC DATABASE MODEL (SDM)

Author: AL-MADANI, NABEIL ISMAEIL
Degree: PH.D.
Year: 1994
Corporate Source/Institution: UNIVERSITY OF SOUTH FLORIDA (0206)
A FORMAL DEFINITION, QUERY LANGUAGE AND RELATIONAL DATABASE MAPPING FOR THE SEMANTIC DATABASE MODEL (SDM)
Year: 1994

Conventional database models, particularly the relational model, provided data independence by separating the conceptual representation of data from the physical implementation. However, these models lack the constructs and mechanisms for modeling the complex interrelationships among data items in a real-world application environment. Semantic database models have been developed to provide a richer set of data structuring constructs and mechanisms. The Semantic Database Model (SDM) is one of the most comprehensive published semantic database models. It incorporates a wide range of constructs for expressing data abstractions, relationships among data items, and data derivation mechanisms. The SDM, however, is defined by means of examples and does not have a formally defined semantics or query language. The primary focus of this dissertation is: (1) to formally define the SDM, (2) to develop a query language for the SDM, (3) to develop a scheme that maps the structures of the SDM to relational structures, and (4) to design and develop a prototype system that implements the SDM as a front-end to a relational DBMS. The formal definition of the SDM uses a graph-based formalism for defining the structural components, and first order predicate calculus for specifying the integrity constraints of the model. The SDM query language is an English-like query language that provides commands for schema definition and modification, data retrieval and manipulation, and data update and modification. The mapping scheme provides the necessary theoretical foundations for developing the SDM Schema Management System (SDM/SMS) prototype. The SDM/SMS implements the SDM on top of a relational DBMS. It allows the user to define and manipulate schemata, browse and navigate inside schemata, and specify queries through an intuitive yet powerful graphical user interface.

36/3,K/66 (Item 41 from file: 8) **Links**
Title: Effective utilization of copies in a transparent distributed environment

Author: Orlowska, Maria E.
Corporate Source: Univ of Queensland, St Lucia, Aust
Abstract: In a distributed **relational** database system, the processing of a query involves data transmission among different sites via a computer network. In a distributed database multiple copies of each **relation** can be allocated to different, **physically** distributed sites. In this paper we discuss the query preoptimization problem for **join**-queries. In general, there is a large number of possibilities to use the copies of the data item in a distributed **relational** database when evaluating a **join**-query. We consider the problem of a copy preselection for each **relation** in a **join sequence** of a **join**-query. We show how to express the preselection problem for a given query and data... ...cover problem. It can be treated as a heuristic for the first phase of a **join**-query optimization, and as such as an input to the final stage of optimization, the execution strategy generation for a **join**-query. In this paper we assumed that a distributed system provides fully transparent data management, i.e., data allocation to the network and data replication which is revealed to a **user**. We illustrate the proposed mathematical programming problem through a nontrivial example. (Author abstract) 18 Refs.
Descriptors: *Distributed database systems; **Query languages**; Data communication systems; **Relational** database systems; Computer networks; Linear programming; Heuristic methods; **User interfaces**; Computer **software selection** and evaluation ; Mathematical programming
Identifiers: Copy preselection; Transparent distributed environment; Query optimization; **Join** query; Data allocation; Integer linear programming; Replicated data

36/3,K/72 (Item 2 from file: 34) <u>Links</u>

04320045 **Genuine Article#: RV576   No. References: 20**
AN EXECUTABLE VISUAL FORMALISM FOR OBJECT-ORIENTED CONCEPTUAL MODELING

Author: KUNG DC ·
Corporate Source: UNIV TEXAS,DEPT COMP SCI ENGN,POB 19015/ARLINGTON//TX/76019
Language: ENGLISH   **Document Type:** ARTICLE   (Abstract Available)
**AN EXECUTABLE VISUAL FORMALISM FOR OBJECT-ORIENTED CONCEPTUAL MODELING
, 1995**
Abstract: Conceptual modeling aims at establishing the conceptual knowledge necessary for proper communication between a development team and users. This article presents an executable visual formalism for object-oriented modeling of information systems. This formalism is an integration of the entity-relationship approach, Petri nets, relational calculus, and time temporal logic. It supports integrated and encapsulated modeling of the structural and behavioral aspects of objects, and object evolution. The formalism has textual and graphical representations, allows formal analysis of model properties, and supports rapid prototyping. An environment and a methodology for conceptual modeling also are described.
Research Fronts: 93-0083 001 (DISTRIBUTED DATABASE DESIGN; OBJECT MODEL; EXTENDED
NESTED RELATIONS)
 93-1567 001 (OBJECT -ORIENTED SOFTWARE SYSTEMS; DESIGN EVOLUTION; BEHAVIOR
NETWORK MODEL FOR CONCEPTUAL INFORMATION MODELING)
 93-1976 001 (OBJECT-ORIENTED MODELING; DICTIONARY OF ENTITY- RELATIONSHIP DATA
SCHEMAS; SOFTWARE ENGINEERING ENVIRONMENT)
 93-4292 001 (FUNDAMENTAL GRAPHICAL PRIMITIVES FOR VISUAL QUERY LANGUAGES;
SPECIFYING MULTIPLE-VIEWED SOFTWARE REQUIREMENTS)
 93-5836 001 (REACTIVE SYSTEMS; REAL-TIME SPECIFICATION USING PETRI NETS; ABSTRACT
DESIGN REPRESENTATIONS; COMPOSITIONAL SEMANTICS; DISTRIBUTED PROTOTYPING)

INSPEC
07325423   **INSPEC Abstract Number:** C1999-09-6160M-021
Title: Specifying generic multimedia 3D visualizations and temporal presentations from database queries
Author Baral, C.; Gonzalez, G.

Author Affiliation: Dept. of Comput. Sci., Texas Univ., El Paso, TX, USA
Conference Title: Proceedings IEEE International Conference on Multimedia Computing and Systems   Part vol.1
p. 550-5 vol.1
Publisher: IEEE Comput. Soc , Los Alamitos, CA, USA
Publication Date: 1999 **Country of Publication:** USA   2 vol. (xlix+909+1127) pp.
ISBN: 0 7695 0253 9   **Material Identity Number:** XX-1999-01834
U.S. Copyright Clearance Center Code: 0 7695 0253 9/99/$10.00
Conference Title: Proceedings of ICMCS99: IEEE Multimedia Systems '99: International Conference on
Multimedia Computing and Systems
Conference Sponsor: IEEE Comput. Soc.; IEEE Circuit & Syst. Soc.; IEEE Commun. Soc.; IEEE Signal Process.
Soc
Conference Date: 7-11 June 1999   **Conference Location:** Florence, Italy
Language: English
Subfile: C
Title: Specifying generic multimedia 3D visualizations and temporal presentations from database queries
Abstract: The extended version of SQL+D presented here takes multimedia to its ultimate realm: the third
dimension, allowing the user to explore the data in a virtual reality environment and display VRML files as
data objects. Unlike the limited number of efforts in this area, SQL+D goes beyond constructing 3D bar
charts. It allows truly interactive environments dynamically specified by the user in a query and constructed
from data extracted from the database. Multimedia presentations are usually specified for a fixed set of
elements. When dealing with database queries, the specific elements are unknown until the answer set is
obtained. We present the syntax of an improved version of SQL+D that allows a user to specify complex
generic temporal presentations within the query itself. We have implemented most of our proposed...
Descriptors: ...SQL; ... ...user interfaces; ... ...virtual reality languages
Identifiers: ...virtual reality environment... ...data objects; ... ...SQL+D... ...answer set;
1999

36/3,K/13 (Item 13 from file: 2) **Links**

06056522   **INSPEC Abstract Number:** C9511-6160J-011
Title: METU object-oriented DBMS

Author Dogac, A.; Ozkan, C.; Arpinar, B.; Okay, T.; Evrendilek, C.
Author Affiliation: Centre of Software Res. & Dev., Middle East Tech. Univ., Ankara, Turkey

Title: METU object-oriented DBMS

Abstract: METU object-oriented DBMS includes the implementation of a database kernel, an object-oriented SQL-like language and a graphical user interface. Kernel functions are divided between a SQL interpreter and a C++ compiler. Thus the interpretation of functions are avoided increasing the efficiency of the system. The compiled by C++ functions are used by the system through the function manager. The system is realized on Exodus Storage Manager (ESM), thus exploiting some of the kernel functions readily provided by ESM. The additional functions provided by the MOOD kernel are the optimization and interpretation of SQL statements, dynamic linking of functions, and catalog management. An original query optimization strategy based on the object-oriented features of the language is developed. For this purpose formulas for the selectivity of a path expression, and for the cost of forward and backward path traversals are derived, and join sizes are estimated. New strategies for ordering the joins and path expressions are also developed. A graphical user interface, namely MoodView is implemented on the MOOD kernel. MoodView provides the database programmer with tools... ...for every phase of OODBMS application development. Furthermore, a database administration tool, a full screen text-editor, a SQL based query manager, and a graphical indexing tool for the spatial data, i.e., R Trees are also...

Descriptors: graphical user interfaces; ... ...object-oriented databases... ...object-oriented languages... ...program compilers; ... ...query languages; ... ...SQL
Identifiers: ...object-oriented database... ...object-oriented query language; ... ...graphical user interface; ... ...SQL interpreter... ...C++ compiler; ... ...Exodus Storage Manager; ... ...full screen text-editor
1994

```
?  d s
Set      Items     Description
S1        1555     S (QUERY? OR QUERIE?)()LANGUAG? OR SQL? ? OR (XML? ? OR EXTEN?()(MARKUP?
OR MARK?()UP)()LANGUAG?)()QUER?
S2           7     S HQL?.? OR QUERYLANGUAG?
S3         192     S (INTERROG? OR RETRIEV? OR OBTAIN? OR REQUIR? OR QUESTION? OR ENQUIR? OR
KEYWORD? OR KEY()WORD?)()LANGUAG?
S4     2130087     S LINK? OR JOIN? OR COMBINE? OR COMBINING? OR COMBINATION? OR COMBINATOR?
OR RELATIONSHIP?
S5     1168008     S AFFILIAT? OR CONJOIN? OR RELATED OR RELATION? OR BUILD?
S6     1472627     S KINSHIP? OR CORRESPOND? OR CORRELAT? OR DRAG?(2W)DROP? OR DRAGANDDROP?
OR DRAGNDROP?
S7     3787360     S ADAPT? OR RECONFIGUR? OR CONFIGUR? OR CHANGE? OR CHANGING? OR MODIF? OR
ALTER? OR CONVERT? OR CONVERSION?
S8      628811     S AMEND? OR UPDAT? OR REVIS? OR REARRANG? OR REORDER? OR RESHUFFL? OR
RESET? OR REFORMAT? OR TRANSFORM?
S9        1292     S S1:S3 AND S4:S8
S10    1028343     S FIELD? OR ARRAY? OR MENU? ? OR DROPDOWN? OR DROP()DOWN? OR GUI? ? OR
(GRAPHIC? OR USER?)(2W)INTERFACE? OR UI OR MODEL?
S11      17241     S INTERFACE?(2N)(APPLICATION? OR SOFTWARE?) OR DIALOG()(BOX OR BOXES) OR
PULLDOWN? OR PULL?()DOWN?
S12      11228     S API? ? OR (APP OR APPS OR APPLICATION?)(2N)(INTERFACE? OR SCREEN? OR
WINDOW? OR PAGE? ? OR BOX OR CELL? ?)
S13        371     S S1:S3 AND S10:S12
S14       1379     S S9 OR S13
S15        701     S USER? OR CLIENT? OR CUSTOMER? OR ENDUSER? OR NETIZEN? OR OPERATOR? OR
MANAGER?
S16         41     S ACCOUNT?(2N)HOLDER? OR PATRON? OR MEMBER? OR SUBSCRIBER? OR WEBUSER?
S17         86     S PARTY? OR PERSON? ? OR INDIVIDUAL? OR PARTIE? OR PRINCIPAL?
S18        379     S DESIRED? OR CHOICE? OR PREFER? OR CHOOS? OR SELECT? OR ELECT? OR OPT? ?
OR OPTING OR OPTED OR OPTION? OR PRECEDEN?
S19         24     S CUSTOMIZ? OR CUSTOMIS? OR PERSONALIS? OR PERSONALIZ? OR INDIVIDUALIZ? OR
INDIVIDUALIS?
S20        178     S ELECT? OR PICK? OR DESIGNAT? OR DISCRIMINAT? OR ASSIGN? OR SPECIFY? OR
PRESET?
S21        266     S STATEMENT?
S22        467     S STRING? OR SET? ? OR SEQUENC? OR BYTE? OR BIT? ? OR CHARACTERS OR LINE?
? OR STREAM? OR PREDICATE?
S23        441     S CHUNK? ? OR TEXT? OR SEGMENT? OR SENTENC? OR SCRIPT? OR EXPRESSION? OR
QUEUE? OR CLAUSE?
S24        367     S LOGICAL? OR ABSTRACT? OR OBJECT? OR DATAOBJECT? OR METADATA? OR
META()DATA?
S25        138     S PHYSICAL? OR REAL? OR EXECUTAB? OR TANGIB? OR RUNNABLE? OR (RUN? ? OR
LAUNCH? OR START?)(2N)CAPAB? OR LAUNCHAB?
S26          8     S USABL? OR UTILIZAB? OR UTILISAB? OR UTILE
S27       1298     S IC=G06F?
S28       1061     S MC=T01?
S29        284     S S9 AND S13
S30        284     S S29 AND S15:S28
S31        316     S S14 AND S15:S17 AND S18:S20
S32         51     S S14 AND (S21:S23 OR S10:S12) AND S24 AND S25:S26
S33         80     S S30 AND S31
S34        265     S S30:S31 AND S21:S23 AND S1:S3 AND S4:S6
S35        259     S S34 AND S27:S28
S36        165     S (S34:S35 OR S14) AND S24:S26 AND S15:S17 AND S21:S23
S37         82     S S36 AND S34:S35
S38        175     S S32:S33 OR S37
S39        110     S S38 AND AC=US/PR
S40         95     S S39 AND AY=(1970:2003)/PR
S41         89     S S39 NOT AY=(2004:2006)/PR
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S42      65     S S38 NOT S39
S43      58     S S42 AND PY=1970:2003
S44      53     S S42 NOT PY=2004:2006
S45     156     S S40:S41 OR S43:S44
S46     156     IDPAT (sorted in duplicate/non-duplicate order)
; show files
```

[File 347] **JAPIO** Nov 1976-2005/Nov(Updated 060302)

[File 350] **Derwent WPIX** 1963-2006/UD,UM &UP=200616
File 350: For more current information, include File 331 in your search. Enter HELP NEWS 331 for details.

46/3,K/129 (Item 129 from file: 350) **Links**
Derwent WPIX

008823290   **Image available**
WPI Acc No: 1991-327303/199145
XRPX Acc No: N91-250706
  Object-oriented query language appts. -
  .includes pre-processor to parse, optimise and translate oql
  statements into host language statements which are compiled
  into executable code
Patent Assignee: TEXAS INSTR INC (TEXI)
Inventor: BLAKELY J A; THOMPSON C W; BLAKELEY J A
Number of Countries: 006 Number of Patents: 006
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
EP 455447	A	19911106	EP 91303851	A	19910429	199145	B
EP 455447	A3	19930630	EP 91303851	A	19910429	199405	
US 5761493	A	19980602	US 90516369	A	19900430	199829	
US 5826077	A	19981020	US 90516369	A	19900430	199849	
			US 95473622	A	19950607		
			US 96639808	A	19960429		
			US 97837396	A	19970417		
EP 455447	B1	19990616	EP 91303851	A	19910429	199928	
DE 69131336	E	19990722	DE 631336	A	19910429	199935	
			EP 91303851	A	19910429		

Priority Applications (No Type Date): US 90516369 A 19900430; US 95473622 A
  19950607; US 96639808 A 19960429; US 97837396 A 19970417
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
EP 455447	A		35		

  Designated States (Regional): DE FR GB IT NL

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
EP 455447	A3		35		
US 5826077	A			G06F-017/30	Cont of application US 90516369
					Cont of application US 95473622
					Cont of application US 96639808
EP 455447	B1	E		G06F-017/60	

  Designated States (Regional): DE FR GB IT NL

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
DE 69131336	E			G06F-017/60	Based on patent EP 455447
US 5761493	A			G06F-017/30	

  Object-oriented query language appts...

  ...includes pre-processor to parse, optimise and translate oql
  statements into host language statements which are compiled
  into executable code
  ...Abstract (Basic): The appts. for processing queries includes a
     preprocessor (144, 148) to parse, optimise and translate **object**
     query language statements into efficient lost

language **statements,** connected to a compiler (152) that compiles
such host language **statements** into **executable** code (154).
A programme is inputted containing **statements** using a host
language's type system as well as employing the host language's
expressions, object composition and inheritance in the
formulation of queries...

...The Select-From-Where style of standard **query language** is
combined as a basic structure for query **statements** with
the **set**-valued function of the host language as targets for
those queries. Thus, an associative query...
Title Terms: **OBJECT;**

-continued

```
| function_expression
| parameter_commalist ',' object_ref
;
```

The approach used in OQL to provide a better integration with host languages can also be used by current SQL systems to provide a better integration with object-oriented programming languages.

The invention provides an elegant way to combine the computer industry's most used industrial strength object-oriented programming language, C++, with the industry's standard database query language SQL. The invention complements a C++ programmer's existing programming tools, by adding a query capability. This can result in better, more reliable C++ code and make database programming more accessible to the large number of C++ programmers.

The query processor architecture of OQL 100 is illustrated in FIG. 7. It consists of four main modules: parser 102, query rewriter 104, rule-based query translator and optimizer 106, and code optimizer 108. All these modules use the run-time data dictionary 112 that contains all type information of the application including class information (class names, data member names and types, and member function names types and parameters), class hierarchy information, and set-valued and Boolean-valued variables and functions used in the applications. The query processor takes as input OQL[C++]'s SELECT statements 124, and produces as output optimized C++ or Persistent C++ code 126. In FIG. 7, lightly shadowed boxes represent run-time data structures, darker boxes represent rule-bases, arrows connecting modules represent transfer of control, and continuous lines represent access to data structures from the corresponding modules. Parser 102 takes as input OQL[C++] SELECT statement and, with the help of the data dictionary, checks the syntax of the statement and whether references to class information and user-defined functions are valid. Syntactically incorrect statements are issued syntax errors 122 and syntactically correct statements are mapped into a query graph 116. After parsing, control passes to query rewrite module 104 which reads query graph 116 and transforms it in place after performing semantic and algebraic query optimization with the help of data dictionary 112 and algebraic rule-base 110, respectively. After query rewrite, control passes to rule-based query translator 106 which with the help of storage interface rule-base 118 and data dictionary 112 translates the query graph into a first version of an execution plan 114 written in C++ or Persistent C++. Finally, code optimizer 108 produces an improved sequence of C++ and Persistent C++ code 126 statements corresponding to original query 124.

The complete compilation process 140 of a program containing OQL[C++] statements (i.e., DECLARE, IMPLEMENT, and SELECT statements) is illustrated in FIG. 8. The input to the compilation process is a program containing OQL[C++] statements 142. This input is processed by the OQL preprocessor which, for SELECT statements, performs all the functions described in query processor 100 (FIG. 7). The OQL preprocessor produces a combination of C++ and Persistent C++ code 146 because of the possibility of a program querying transient and persistent data. A persistence C++ preprocessor 148 transforms Persistent C++ code into pure C++ code 150 and C++ code passes this process unchanged. C++ code 150 is then processed by a standard C++ compiler 152 which produces executable code 154.

FIG. 9 illustrates a possible placement of an application code 164 containing OQL[C++] query statements 166,

according to a preferred embodiment of the invention, in the main memory 162 of a computer 160. During the development of the application, a computer programmer enters a program to the computer 160 via a computer terminal 180. The application is compiled according to the process described in FIG. 8 where OQL[C++] statements 142 (FIG. 8)/166 (FIG. 9) are translated into pure C++ code 168 or Persistent C++ code 178 depending on whether queries access transient 170 or persistent data 172. Persistent data 172 is stored in a persistent database 176 on disk 174. Queries on transient data are performed using C++ code 168 in main memory 170. Queries on persistent data are performed using Persistent C++ code 178 which ensures the transfer of persistent data from the database 176 on disk 174 to a special area 172 of main memory 162 where the query is executed. It should be noted that although application code 164, OQL[C++] code 166, C++ code 168, and Persistent C++ code 178 are distinguishable before compilation, they are a single, indistinguishable piece of execution code after compilation.

The description of OQL[C++] as implemented by a a C++ preprocessor is specific to the implemented embodiment. If native C++ compilers (not implemented by preprocessors) were extended to use OQL, similar operations would be implemented by those compilers. In fact, if incremental C++ compilers become available, then OQL[C++] will be able to be similarly incremental and still inherit the "ad hoc" interactive ability to state queries. This feature would be inherited immediately in OQL[CLOS] when applied according to the present invention since CLOS allows incremental compilation.

While a specific embodiment of the invention has been shown and described above in this patent application, various modifications and alternate embodiments will occur to those skilled in the art. Accordingly, it is intended that the invention be limited only in terms of the appended claims.

We claim:

1. A method for enabling an integration of an object query language for an object-oriented data model with a host language, wherein queries for accessing objects in the object oriented data model are in said object query language, comprising the steps of:

constructing a preprocessor including a grammar coupling the object query language with the host language, wherein said grammar includes query statements to use expressions of the host language in accessing the objects and wherein said query statements access both transient and persistent objects by accessing a data dictionary having pointers to said transient and persistent objects to maintain information corresponding to both said transient and persistent objects;

using said preprocessor to preprocess said query statements by parsing, optimizing and translating at least one statement in said object query language into at least one statement of the host language; and

compiling said statement of the host language into executable code.

2. A method for executing a query in an object query language operating in conjunction with a host language, comprising the steps of:

defining the object query language with set declarations;

formulating query statements by combining a type system of the host language as a database model with preselected aspects of a relational query language as a basic structure for said query statements;

preprocessing said query statements into host language statements, said host language statements to access

and Boolean-valued variables and functions used in the applications. The query processor takes as input OQL[C++]'s SELECT statements 124, and produces as output optimized C++ or Persistent C++ code 126. In FIG. 4, lightly shadowed boxes represent run-time data structures, darker boxes represent rule-bases, arrows connecting modules represent transfer of control, and continuous lines represent access to data structures from the corresponding modules. Parser 102 takes as input OQL[C++] SELECT statement and, with the help of the data dictionary, checks the syntax of the statement and whether references to class information and user-defined functions are valid. Syntactically incorrect statements are issued syntax errors 122 and syntactically correct statements are mapped into a query graph 116. After parsing, control passes to query rewrite module 104 which reads query graph 116 and transforms it in place after performing semantic and algebraic query optimization with the help of data dictionary 112 and algebraic rule-base 110, respectively. After query rewrite, control passes to rule-based query translator 106 which with the help of storage interface rule-base 118 and data dictionary 112 translates the query graph into a first version of an execution plan 114 written in C++ or Persistent C++. Finally, code optimizer 108 produces an improved sequence of C++ and Persistent C++ code 126 statements corresponding to original query 124.

The complete compilation process 140 of a program containing OQL[C++] statements (i.e., DECLARE, IMPLEMENT, and SELECT statements) is illustrated in FIG. 5. The input to the compilation process is a program containing OQL[C++] statements 142. This input is processed by the OQL preprocessor which, for SELECT statements, performs all the functions described in query processor 100 (FIG. 4). The OQL preprocessor produces a combination of C++ and Persistent C++ code 146 because of the possibility of a, program querying transient and persistent data. A persistence C++ preprocessor 148 transforms Persistent C++ code into pure C++ code 150 and C++ code passes this process unchanged. C++ code 150 is then processed by a standard C++ compiler 152 which produces executable code 154.

FIG. 6 illustrates a possible placement of an application code 164 containing OQL[C++] query statements 166, according to a preferred embodiment of the invention, in the main memory 162 of a computer 160. During the development of the application, a computer programmer enters a program to the computer 160 via a computer terminal 180. The application is compiled according to the process described in FIG. 5 where OQL[C++] statements 142 (FIG. 5)/166 (FIG. 6) are translated into pure C++ code 168 or Persistent C++ code 178 depending on whether queries access transient 170 or persistent data 172. Persistent data 172 is stored in a, persistent database 176 on disk 174. Queries on transient data are performed using C++ code 168 in main memory 170. Queries on persistent data are performed using Persistent C++ code 178 which ensures the transfer of persistent data from the database 176 on disk 174 to a special area 172 of main memory 162 where the query is executed. It should be noted that although application code 164, OQL[C++] code 166, C++ code 168, and Persistent C++ code 178 are distinguishable before compilation, they are a single, indistinguishable piece of execution code after compilation.

The description of OQL[C++] as implemented by a a C++ preprocessor is specific to the implemented embodiment. If native C++ compilers (not implemented by preprocessors) were extended to use OQL, similar operations would be

implemented by those compilers. In fact, if incremental C++ compilers become available, then OQL[C++] will be able to be similarly incremental and still inherit the "ad hoc" interactive ability to state queries. This feature would be inherited immediately in OQL[CLOS] when applied according to the present invention since CLOS allows incremental compilation.

While a specific embodiment of the invention has been shown and described above in this patent application, various modifications and alternate embodiments will occur to those skilled in the art. Accordingly, it is intended that the invention be limited only in terms of the appended claims.

We claim:

1. A method enabling an integration of an object query language for an object-oriented data model with a host language, wherein queries for accessing objects are in said object query language, comprising the steps of:

   constructing a preprocessor including a grammar coupling the object query language with the host language, wherein said grammar includes query statements to use expressions of the host language in accessing the objects and wherein said query statements access both transient and persistent objects;

   using said preprocessor to preprocess said query statements by parsing, optimizing and translating at least one statement in said object query language into at least one statement of the host language; and

   compiling said statement of the host language into executable code.

2. The method of claim 1 wherein said preprocessing step includes the step of global type checking of said query statements.

3. A method for executing a query in an object query language operating in conjunction with a host language, comprising the steps of:

   defining the object query language with set declarations;

   formulating query statements by combining a type system of the host language as a database model with preselected aspects of a relational query language as a basic structure for said query statements;

   preprocessing said query statements into host language statements, said host language statements being adapted to access transient and persistent objects;

   compiling said host language statements; and

   executing said compiled host language statements to return at least one result in response to said formulated query statement.

4. The method of claim 3, wherein the method further comprises the step of defining said set declarations to include at least one membership function.

5. The method of claim 3, wherein the method further comprises the step of defining said set declarations to include at least one set operation.

6. The method of claim 3, wherein the method further comprises the step of defining said set declarations to include at least one iteration operator.

7. The method of claim 3, wherein the method further comprises the step of defining said set declarations to include at least one indexing function.

8. The method of claim 3, wherein said step of formulating query statements further comprises the step of combining boolean-valued functions of the host language with at least one predicate.

9. The method of claim 3, wherein said step of formulating query statements farther comprises the step of setting targets for said queries by using set-valued functions of the host language.

46/3,K/114 (Item 114 from file: 350) **Links**
·Derwent WPIX

011109053
WPI Acc No: 1997-086978/199708
XRPX Acc No: N97-071730
  **GUI for database retrieval - generates retrieval
  instructions of query language program by set
  creating, drawing, selection, projection, and extraction, and
  graphs quantity of and logical relation between sets
  retrieved**
Patent Assignee: NAGAMORI N (NAGA-I)
Inventor: NAGAMORI N
Number of Countries: 001   Number of Patents: 001
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 5592663	A	19970107	US 93168219	A	19931217	199708	B

Priority Applications (No Type Date): US 93168219 A 19931217
Patent Details:·

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 5592663	A			G06F-017/30	

  GUI for database retrieval...

**...generates retrieval instructions of query language program
  by set creating, drawing, selection, projection, and
  ·extraction, and graphs quantity of and logical relation
  between sets retrieved**

...Abstract (Basic): The program communicates with the database
    query language program through message exchanges between
    the programs or message exchanges at the terminal systems connected
    with the communication network. The retrieval instructions of the
  ' **query language** program are generated by way of **set**
  · creating, **set** drawing, **set selection, set**
    projection, and **set** extraction. The input operation is performed
    in place of the **user.**
    ...

...ADVANTAGE - **Graphical user interface** shows with
    graphs, quantity of and **logical relation** between
    sets of database retrieved, to calculate and manipulate
    sets by pointing area of rectangular shapes of graphs, and to
    extract data by pointing area of rectangular shapes of graphs.
    Facilitates simple retrieval processing and supports **logical**
    thinking of **user** to evaluate and judge quantity of **set**
    and **relation** between **sets.**

...Title Terms: **SET**;
International Patent Class (Main): **G06F-017/30**
Manual Codes (EPI/S-X): **T01-J05B3**

United States Patent [19]

Nagamori

[11] Patent Number: 5,592,663

[45] Date of Patent: Jan. 7, 1997

[54] **GRAPHING METHOD AND APPARATUS FOR DATA BASE RETRIEVAL**

[76] Inventor: **Nobuhiko Nagamori**, 5-5-13, Kunugidaidanchi, 1404, Kawashima-cho, Hodogaya-ku, Yokohama-shi, Kanagawa 240, Japan

[21] Appl. No.: **168,219**

[22] Filed: **Dec. 17, 1993**

[51] Int. Cl.^{6} G06F 17/30

[52] U.S. Cl. 395/605; 364/488; 364/283.4; 364/DIG. 1

[58] Field of Search 395/600, 100, 395/575; 364/468, 401, 488

[56] **References Cited**

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Primary Examiner—Thomas G. Black
Assistant Examiner—C. Pham
Attorney, Agent, or Firm—Ronald R. Snider

[57] **ABSTRACT**

The purpose of the present invention is to provide a graphical user interface which enables to show with graphs, as in FIG. 1, the quantity of and the logical relation between the sets of the database retrieved, to calculate and manipulate the sets by pointing the area of the rectangular shapes of the graphs, and to extract the data by pointing the area of the rectangular shapes of the graphs. The program of the present invention communicates with the database query language program through message exchanges between the programs or message exchanges at the terminal systems connected with the communication network, generates the retrieval instructions of the query language program by way of set creating, set drawing, set selection, set projection, and set extraction, and performs the input operation in place of the user. The present invention facilitates simple retrieval processing and supports logical thinking of the user to evaluate and judge the quantity of the set and the relation between the sets.

20 Claims, 15 Drawing Sheets

RECTANGULAR REGION INDICATING SPACE OBJECTS

REGION OBJECT PROJECTING SET OBJECT
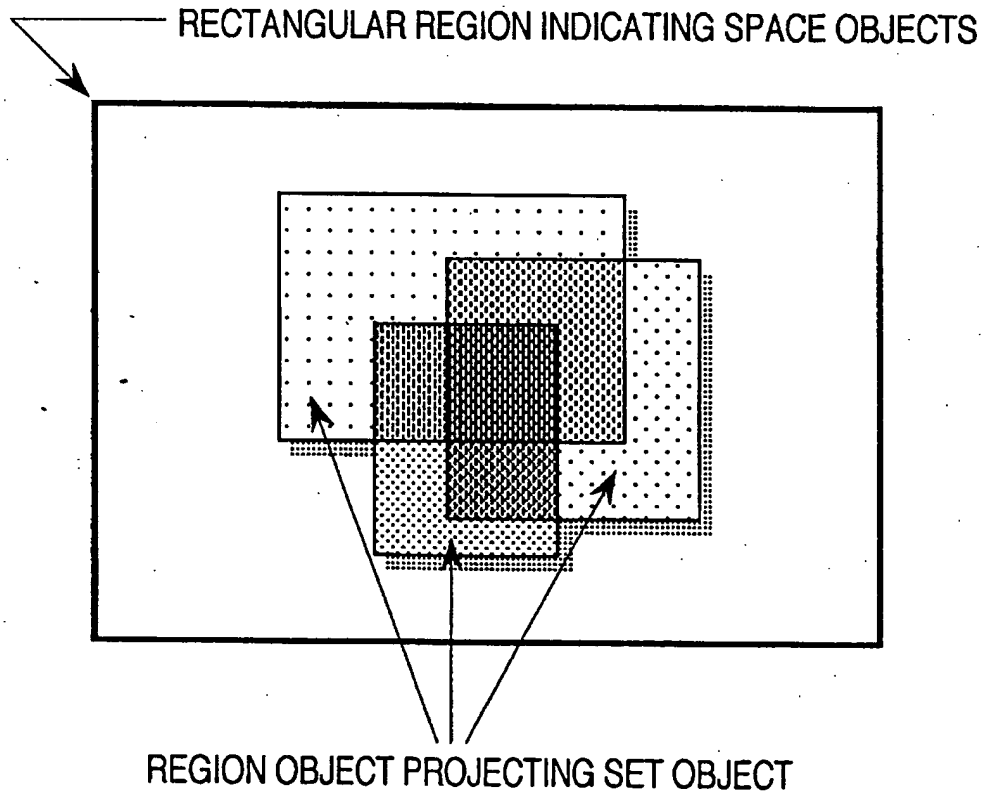
As shown in FIG. 14, the read set button (readSet) changes to "OK?" and a message "Select a group of fields to extract from the field object browser" is shown in the notifying field. Click the fields to be extracted, from the field object browser.

072 001-9I:read Items=:personal_id, name, age, dept_no,
073 001-1C:readSet <+>(('A'*'B'*'D')+('A'*'C'*'D')+('B'*'C'*'D'))<->('A'*'B'*'C'*'D') [personal_id, name, age, dept_no].
074 001-3S:Host:012:"((SELECT PERSONAL_NO, NAME, AGE, DEPT_CODE TECH_HISTORY.EMP_NO FROM EMP, TECH_HISTORY WHERE AGE≧30 AND AGE<35 AND YEARS>12 AND TECH_HISTORY.SYSTEM='UNIX' EMP.PERSONAL_NO=TECH_HISTORY.EMP_NO)
UNION
(SELECT PERSONAL_NO, NAME, AGE, DEPT_CODE, TECH_HISTORY.EMP_NO FROM EMP, TECH_HISTORY WHERE AGE≧30 AND AGE<35 AND PROJ_CODE IS NULL AND TECH_HISTORY.SYSTEM='UNIX' EMP.PERSONAL_NO=TECH_HISTORY.EMP_NO)
UNION
(SELECT PERSONAL_NO, NAME, AGE, DEPT_CODE, TECH_HISTORY.EMP_NO FROM EMP, TECH_HISTORY WHERE YEARS>12 AND PROJ_CODE IS NULL AND TECH_HISTORY.SYSTEM='UNIX' EMP.PERSONAL_NO=TECH_HISTORY.EMP_NO))
MINUS
(SELECT PERSONAL_NO, NAME, AGE, DEPT_CODE, TECH_HISTORY.EMP_NO FROM EMP, TECH_HISTORY WHERE AGE≧30 AND AGE<35 AND YEARS>12 AND PROJ_CODE IS NULL AND TECH_HISTORY.SYSTEM='UNIX' EMP.PERSONAL_NO= TECH_HISTORY.EMP_NO);
075 001-3R:Host:012:"ready for output"

As explained in the above example, "the database retrieval method through set graph" realizes retrieval processing with the use of the graphical user interface to select and point the graphic objects, except the basic conditional formula of the set objects described through the phased operational flow (FIG. 15) which follows the logical way of human thinking.

The present invention brings about the below effects, by way of displaying the sets projected in the geometric rectangular shapes,
(a) to evaluate and recognize visually the quantity of the sets and logical relation between the sets from the size and layout of the geometric rectangular shapes,
(b) to realize data retrieval manipulation by creating 255 different retrieval conditional formulae, with the instruction of selecting combinations of the sets from eight logical sets, including a maximum of seven regional portions and the surrounding mother set, which are drawn with three rectangular regions, and
(c) to use the set graphs for comparing the sets with the time scale as well as for better explanation and presentation, because the quantity and relation of the abstract existence of the sets are recorded and shown as graphs.

In the database retrieval, such as information and documentation retrieval, the object of retrieval is not always clearly grasped, so under an uncertain and unclear tentative condition, retrieval processing is normally repeated on a trial and error basis. In such a retrieval processing, groping for the object itself under various retrieval conditions is meaningful, and retrieval is executed with a comparatively simple

formula, and the sets are created and the object of retrieval is understood from the logical relation of those sets manipulated with a combination of the created sets using the boolean operators.

Furthermore, in case of the database retrieval of the personnel or management information, the best condition for retrieval can be obtained through describing complex conditions against multiple database files and make combinations of the sets obtained from the assumed multiple conditional values.

Such type of retrieval is very important and often found difficult to achieve because the logical challenge to understand the logical relation between the sets is more difficult than the technical challenge in computer manipulation. Therefore, it is still often observed that the user himself draws circles to show the sets on a sheet of paper to try to understand the relation of the sets from the drawing, and that after such manipulation, he can finally return to the keyboard of the computert to enter the retrieval conditions.

The present invention provides an integrated solution for the technical problem in the course of operation as well as for the logical problem in the human thinking, in highly complex retrieval processing or uncertain and unclear retrieval processing.

What is claimed is:

1. A database retrieval method comprising the steps of:
   a) selecting a first space object database;
   b) creating a plurality of set objects for retrieval of data from said space object database;
   c) displaying a visual representation of said plurality of set objects, wherein said display gives overlapping areas of said visual representation when data in said set objects is the same;
   d) selecting one or more areas of said display; and
   e) extracting data from said space object which is within said one or more areas selected.

2. The retrieval method of claim 1 wherein said plurality of sets does not exceed three.

3. The retrieval method of claim 1 further comprising the step of displaying a number of hit elements within one or more areas of said display.

4. The retrieval method of claim 1 further comprising the steps of:
   pre-registering set object definitions, and
   selecting at least one of said created plurality of set objects by selecting said pre-registered set object definitions.

5. The retrieval method in accordance with claim 2 wherein a count of each of said created set objects is displayed on a screen of a graphical display unit.

6. The retrieval method in accordance with claim 1 wherein said visual representation display shows a logical relationship among the created set objects.

7. The retrieval method in accordance with claim 6 wherein said visual representation display is a geometric rectangular region.

8. The retrieval method in accordance with claim 6 further comprising the step of displaying a logical statement of the relationship of the created set objects.

9. The retrieval method in accordance with claim 1 further comprising the steps of:
   creating a first data base set object from said extracted data of said first data base;
   selecting a second space object data base;
   creating a second data base set object which is defined by one or more set objects selected from said second data base;

creating a first visual display of said first data base set object;

creating a second visual display of said second data base set object;

displaying a visual representation of said first visual display of said first data base set object and of said second visual display of said second data base set object; and

extracting data from said first and second data bases which is within one or more areas selected from said first visual display of said first data base set object and of said second visual display of said second data base set object.

10. The retrieval method in accordance with claim 1 wherein said visual representation is drawn in a mode selected from the group consisting of rate, fit, and zoom.

11. A database retrieval apparatus comprising:

a) a means for selecting a space object database;

b) a means for creating a plurality of set objects for retrieval of data from said space object database;

c) a means for displaying a visual representation of said plurality of set objects, wherein said display gives overlapping areas of said visual representation when data in said set objects is the same;

d) a means for selecting one or more area of said display; and

e) a means for extracting data from said space object which is within said one or more areas selected.

12. The retrieval apparatus of claim 11 wherein said plurality of sets does not exceed three.

13. The retrieval apparatus of claim 11 further comprising a means for displaying a number of hit elements within one or more areas of said display, and said space object.

14. The retrieval apparatus of claim 11 further comprising:

a means for pre-registering set objects, and

a means for selecting at least one of said created set of objects by selecting said pre-registered set object.

15. The retrieval apparatus in accordance with claim 11 wherein a count of said created set objects is displayed on a screen of a graphical display unit.

16. The retrieval apparatus in accordance with claim 11 wherein said display shows a logical relationship among the created set objects.

17. The retrieval apparatus in accordance with claim 16 wherein said display is a geometric rectangular region.

18. The retrieval apparatus in accordance with claim 16 further comprising a means for displaying a logical statement of the relationship of the created set objects.

19. The retrieval apparatus in accordance with claim 11 further comprising:

a means for creating another set object from a space object database;

a means for creating another set object which is defined by the one or more areas selected when logically combined;

a means for displaying a visual display of said another set object obtained from a space object database;

a means for displaying a visual display of said object which is defined by one or more areas selected when logically combined;

a means for selecting one or more areas of said second visual display; and

a means for extracting data from said space object which is within one or more areas selected from said second visual display.

20. The retrieval apparatus according to claim 11 wherein said visual representation is drawn in a mode selected from the group consisting of rate, fit, and zoom.

* * * * *

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015213993     **Image available**
WPI Acc No: 2003-274530/200327
XRPX Acc No: N03-217808
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   **Object oriented interface provision method using
   internet for relational database, involves presenting
   object model mapped onto database schema in user
   -friendly graphical user interface to users**

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Patent Assignee: AT & T CORP (AMTT  )
Inventor: NESHATFAR S; WARTY P
Number of Countries: 001  Number of Patents: 001
Patent Family:
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Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 6490581	B1	20021203	US 2000577103	A	20000524	200327	B

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Priority Applications (No Type Date): US 2000577103 A 20000524
Patent Details:
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Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 6490581	B1		15	G06F-017/30	

   **Object oriented interface provision method using
   internet for relational database, involves presenting
   object model mapped onto database schema in user
   -friendly graphical user interface to users**

Abstract (Basic):
...   An **object model** mapped onto a database schema, is
   presented in a **user**-friendly **graphical user**
   interface. The graphical display is **modified** to prompt a
   user for setting an **object** constraint, in response to the
   selections of presented **objects**. A query is generated
   using the **set object** constraint.
...   For providing **object** oriented interface to **users**
   in **relational** database management system, using internet...

....As the **object model** mapped onto the database schema is
   presented in a **user**-friendly **graphical user**
   interface, allows **users** with no knowledge of the database
   query language to perform sophisticated queries with just
   a few mouse clicks...

...The figure shows the **graphical user interface** for
   presenting information...
Title Terms: **OBJECT;**

International Patent Class (Main): **G06F-017/30**
Manual Codes (EPI/S-X): **T01-J05B3...**

...T01-J05B4B...

...T01-J12A...

...T01-N01

(12) **United States Patent**

Neshatfar et al.

(10) Patent No.: **US 6,490,581 B1**

(45) Date of Patent: **Dec. 3, 2002**

(54) **SYSTEM AND METHOD FOR PROVIDING AN OBJECT-ORIENTED INTERFACE TO A RELATIONAL DATABASE**

(75) Inventors: **Shapour Neshatfar**, Manalapan, NJ (US); **Pramod Warty**, Freehold, NJ (US)

(73) Assignee: **AT&T Corp.**, New York, NY (US)

(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 0 days.

(21) Appl. No.: **09/577,103**

(22) Filed: **May 24, 2000**

(51) Int. Cl.^{7} **G06F 17/30**

(52) U.S. Cl. .... **707/4; 707/103 R**

(58) Field of Search .... 707/2, 3, 4, 10, 707/103 R, 103 Z

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Primary Examiner—Jack Choules

(57) **ABSTRACT**

A method provides for an easily understood representation of the elements and relationships captured in a relational database. In accordance with the method an object model specification file and a database schema file are parsed to create a user-friendly graphical representation of object models and their relationship. The user is able perform a query by selecting one or more objects and setting constraints using the graphical representation of object models.

6 Claims, 9 Drawing Sheets
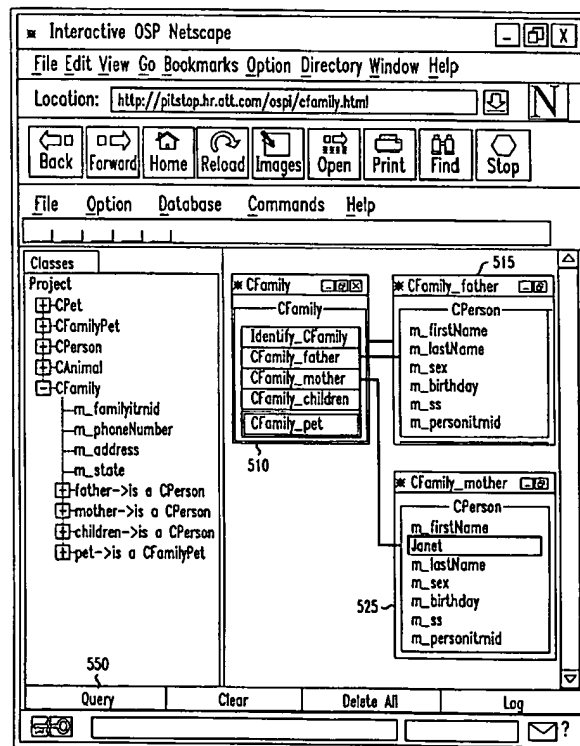
own collection of favorite books, the association of books to a child is easier to visualize when children are shown one at a time with a button labeled favorite Books, where clicking on that button, reveals the books.

Queries can also be made over multiple objects, related or otherwise. Consider the following query:

List pair of families with the same father's name (first name).

The above query can be performed by instantiating (clicking) the CFamily object twice. Expanding the father object on both, and drag the m_firstName from one father and drop it to another one and click the Query button as illustrated in FIG. 7.

Conclusion

The present invention takes advantage of the mapping performed by the OSP tool to present an easily understood description of the relationship between elements in a relational database. Furthermore, this new tool, OSPI, permits casual users to generate queries for complex databases.

What is claimed is:

1. A method of facilitating user operations on a relational database, comprising:

mapping an object model onto a database schema;

presenting the object model in a graphical user interface;

responsive to selections of presented objects, modifying said display to prompt a user to set an object constraint; and

generating a query using said constraint.

2. The method of claim 1 further comprising:

generating a web page based on the results of mapping said object model onto said database schema; and

said step of presenting includes providing said generated web page to a user.

3. The method of claim 1 wherein said step of presenting includes delineating in said graphical user graphical user interface,

a plurality of objects and relationships between the objects in said plurality of objects.

4. A method for presenting relational database information, comprising:

parsing an object model specification file and a database schema file to create a parse tree;

constructing a graphical representation of an object model using said parse tree;

transmitting said graphical representation to a user for display on a graphical user interface.

5. The method of claim 4 wherein said step of constructing comprises generating a web page representation of a plurality of objects and relationships among objects in said plurality.

6. The method of claim 5 further comprising:

receiving an object model selection;

presenting a list of objects corresponding to the selected object model;

receiving a selection of an object selected from said list of objects; and

transmitting constraint parameter information in response to said selection of an object.

* * * * *

46/3,K/75 (Item 75 from file: 350) **Links**
Derwent WPIX

014213980   **Image available**
WPI Acc No: 2002-034678/200204
XRPX Acc No: N02-026663
  Database records relationship display method for
  relational database browsing, involves adding list of possible
  access pathways between related-recorded list based on user
  input and other records in database
Patent Assignee: ORILLION CORP (ORIL-N)
Inventor: LINDSEY T P
Number of Countries: 095 Number of Patents: 005
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
WO 200186504	A2	20011115	WO 2001US14581	A	20010504	200204	B
JP 2001318816	A	20011116	JP 2000314380	A	20001013	200208	
AU 200161215	A	20011120	AU 200161215	A	20010504	200219	
EP 1407382	A2	20040414	EP 2001935091	A	20010504	200426	
			WO 2001US14581	A	20010504		
AU 2001261215	A8	20051020	AU 2001261215	A	20010504	200615	

Priority Applications (No Type Date): US 2000565857 A 20000505
Patent Details:
Patent No Kind Lan Pg   Main IPC   Filing Notes
WO 200186504   A2 E 27 G06F-017/30
    Designated States (National): AE AG AL AM AT AU AZ BA BB BG BR BY BZ CA
    CH CN CO CR CU CZ DE DK DM DZ EE ES FI GB GD GE GH GM HR HU ID IL IN IS
    JP KE KG KP KR KZ LC LK LR LS LT LU LV MA MD MG MK MN MW MX MZ NO NZ PL
    PT RO RU SD SE SG SI SK SL TJ TM TR TT TZ UA UG UZ VN YU ZA ZW
    Designated States (Regional): AT BE CH CY DE DK EA ES FI FR GB GH GM GR
    IE IT KE LS LU MC MW MZ NL OA PT SD SE SL SZ TR TZ UG ZW
JP 2001318816 A   19 G06F-012/00
AU 200161215 A   Based on patent WO 200186504
EP 1407382   A2 E   G06F-017/30   Based on patent WO 200186504
    Designated States (Regional): AT BE CH CY DE DK ES FI FR GB GR IE IT LI
    LU MC NL PT SE TR
AU 2001261215 A8   G06F-017/30   Based on patent WO 200186504

  Database records relationship display method for
  relational database browsing, involves adding list of possible
  access pathways between related-recorded list based on user
  input and other records in database

Abstract (Basic):
...   A list of possible **relationships** defining access pathways
    between a **selected** record and other records in a database, are
    displayed. A list of records **related** to **selected** records
    are added, based on **user** input. A list of possible

relationships defining pathways between the **related** record list and other records are added to the display.

...   a) Computer-readable medium storing record **relationship** display program...

...b) Computer programmed to execute the database records **relationship** display program...

...For browsing **relational** database e.g. Oracle, Sybase, Microsoft Access, using standard **query language** (**SQL**).
        ...

...The data structures of different **relational** databases are viewed in a browsable.format, easily by **choosing** any record as the initial point. The records within a database are easily verified and the **corresponding fields** or tables of various **relational** databases are renamed by the **user**.
        ...

...The figure shows multiple **relational** databases
...Title Terms: **RELATED;**
International Patent Class (Main): **G06F-012/00**...
...**G06F-017/30**
Manual Codes (EPI/S-X): **T01-J05B3**...

...**T01-J05B4B**...

...**T01-J12B**...

...**T01-S03**

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Declarations under Rule 4.17:
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For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

(54) Title: BROWSER ORIENTED METHOD TO VIEW CONTENTS OF A PLURALITY OF DATABASES
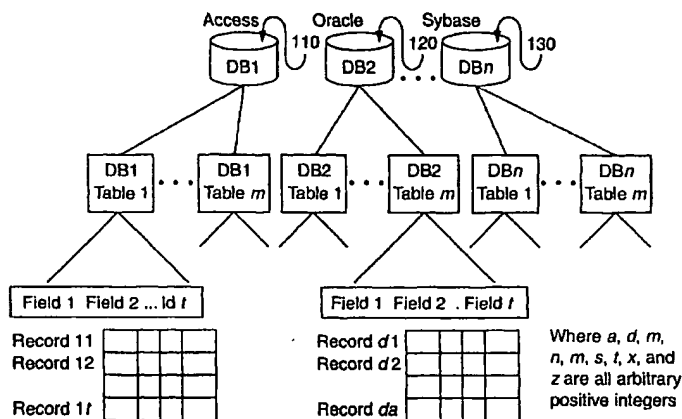
(57) Abstract: A machine-executed method and apparatus for displaying relationships between records in one or more databases in a browseable format is disclosed. Each database comprises one or more tables. Each table includes one or more fields and one or more records. The method entails: (a) displaying, on a display, a field value of a selected record in a table of a selected database; (b) displaying on the display a list of possible relationships between the selected record and other records in the database, where each possible relationship defines a pathway from the selected record to a list of other records; (c) receiving a user input signal designating one of the possible relationships as a selected relationship; (d) adding to the display at least a portion of a list of records that are related to the selected record through the selected relationship, referred to as the related-records list; and (e) adding to the display a list of possible relationships between the related-records list and other records in the database, where each possible relationship defines a pathway from the related-records list to a list of other records.

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WHAT IS CLAIMED IS:

1.   A method of displaying relationships between records in a database, the database comprising a plurality of tables, each table containing one or more records divided into one or more fields, the method comprising:

(a)   determining the respective structures of the plurality of relational databases;

(b)   displaying, on a display, a field value of a selected record in a table of the database;

(c)   displaying on the display a list of possible relationships between the selected record and other records in the database, where each possible relationship defines a pathway from the selected record to a list of other records;

(d)   receiving a user input signal designating one of the possible relationships as a selected relationship;

(e)   adding to the display at least a portion of a list of records that are related to the selected record through the selected relationship, referred to as the related-records list; and

(f)   adding to the display a list of possible relationships between the related-records list and other records in the database, where each possible relationship defines a pathway from the related-records list to a list of other records.

2.   The method of claim 1, further comprising:

(g)   receiving a user input signal designating a record in the related-records list as a new selected record; and

(h)   repeating, for the new selected record, the operations recited in subparagraphs (d) through (f) of claim 1.

3.   A method of displaying relationships between records of a plurality of relational databases of arbitrary structure, each database comprising a plurality of tables, each table containing one or more records divided into one or more fields, each table having a name, and each field having a name, the method comprising:

(a)   determining the respective structures of the plurality of relational databases;

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(b)   integrating the respective structures of the plurality of relational databases into a front-end integrated data structure;

(c)   displaying, on a display, a field value of a selected record in one of the plurality of tables in the plurality of relational databases;

5   (d)   displaying on the display a list of possible relationships between the selected record and other records in the plurality of relational databases, where each possible relationship defines a pathway from the selected record to a list of other records;

(e)   receiving a user input signal designating one of the possible relationships as a selected relationship;

10   (f)   adding to the display at least a portion of a list of records that are related to the selected record through the selected relationship, referred to as the related-records list; and

(g)   adding to the display a list of possible relationships between the related-records list and other records in the plurality of relational databases, where each 15 possible relationship defines a pathway from the related-records list to a list of other records.

4.   The method of claim 3, further comprising:

(a)   receiving a user input signal designating a record in the related-records list as a new selected record; and

(b)   repeating, for the new selected record, the operations recited in 20 subparagraphs (e) through (g) of claim 3.

5.   The method of claim 3, wherein integrating the respective structures comprises:

removing the conflicting names of tables and fields in the front-end integrated data structure; and

25   performing a join operation to establish a link among selected fields in at least two respective tables.

6.   The method of claim 5, wherein removing the conflicting names comprises:

displaying, for a selected name, a cross-reference list of uses of the selected name in the plurality of relational databases;

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receiving at least one editing command from the user, the at least one editing command designating an operation to be performed on a selected table or on a selected field; and

executing the at least one editing command.

7.   The method of claim 6, wherein the at least one editing command comprises: receiving an alias designator which designates an alias for a selected name; and displaying the alias in lieu of the selected name.

8.   The method of claim 7, further comprising: storing a representation of the relationship between the alias and the selected name.

9.   A computer-readable medium encoded with instructions that, when executed by a computer, perform a method in accordance with a specified one of claims 1 through 8.

10.   A computer programmed to perform a method in accordance with a specified one of claims 1 through 8.

46/3,K/119 (Item 119 from file: 350) **Links**
Derwent WPIX

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010443433      **Image available**
WPI Acc No: 1995-344752/199544
XRPX Acc No: N95-257636
```
**Database query system for forming semantically correct
queries - uses query expert system to monitor structure of query and
prevent user from building semantically incorrect queries**

```
Patent Assignee: SOFTWARE AG (SOFT-N); SPEEDWARE LTEE (SPEE-N)
Inventor: SHWARTZ S P; SCHWARTZ S P
Number of Countries: 062  Number of Patents: 010
Patent Family:
Patent No      Kind    Date       Applicat No     Kind    Date      Week
WO 9526003     A1   19950928   WO 95IB517       A    19950323  199544  B
AU 9526816     A    19951009   AU 9526816       A    19950323  199603
US 5584024     A    19961210   US 94217099      A    19940324  199704
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                               WO 95IB517       A    19950323
CA 2186345     C    20000606   CA 2186345       A    19950323  200041
                               WO 95IB517       A    19950323
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Priority Applications (No Type Date): US 94217099 A 19940324; US 96723962 A
   19960926
Patent Details:
Patent No  Kind Lan Pg   Main IPC     Filing Notes
WO 9526003     A1 E 102 G06F-017/30
   Designated States (National): AM AT AU BB BG BR BY CA CH CN CZ DE DK EE
   ES FI GB GE HU JP KE KG KP KR KZ LK LR LT LU LV MD MG MN MW MX NL NO NZ
   PL PT RO RU SD SE SI SK TJ TT UA US UZ VN
   Designated States (Regional): AT BE CH DE DK ES FR GB GR IE IT KE LU MC
   MW NL OA PT SD SE SZ UG
AU 9526816     A       G06F-017/30   Based on patent WO 9526003
US 5584024     A    63 G06F-017/30
EP 803100      A1 E    G06F-017/30   Based on patent WO 9526003
   Designated States (Regional): AT BE CH DE DK ES FR GB IE IT LI NL PT SE
JP 9510565     W   124 G06F-017/30   Based on patent WO 9526003
US 5812840     A       G06F-017/30   Cont of application US 94217099
                                     Cont of patent US 5584024
MX 9604236     A1      G06F-017/30
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EP 803100      B1 E    G06F-017/30   Based on patent WO 9526003
   Designated States (Regional): AT BE CH DE DK ES FR GB IE IT LI NL PT SE
DE 69514123    E       G06F-017/30   Based on patent EP 803100
                                     Based on patent WO 9526003
CA 2186345     C  E    G06F-017/30   Based on patent WO 9526003
```

**... uses query expert system to monitor structure of query
and prevent user from building semantically incorrect
queries**

...Abstract (Basic): The system includes a conceptual layer **manager**
for storing conceptual information about the database including a
predetermined structure and a query assistant which provides the
user with a **set** of permissible **selections** from
which to **build** a semantically correct database query for the
database in an intermediate **query language**. A query
generator receives a query in the intermediate **query
language** from the query assistant and **converts** the query
into the target **query language**.

   ...

...The **user** enters a query through a **set** of dialogue boxes in
an intermediate language. A query expert system monitors the query as
it is built, and using the information about the structure of the
database, prevents the **user** from **building** semantically
incorrect queries by disallowing **choices** in the dialogue box
which would create incorrect queries...

...USE/ADVANTAGE - Guiding **user** to interactively create correct
queries using structured **query language, SQL.**
Permits **user** to enter only queries that are both syntactically
and semantically valid
...Abstract (Equivalent): A database query system for interactively
creating, with a **user,** a syntactically and semantically correct
query for a **relational** database having a plurality of tables,
each of said tables having a plurality of columns and having a
predetermined **relationship** to another of said tables, said
system comprising...

...a conceptual layer **manager** for storing conceptual information
about the **relational** database, said conceptual information
including structural information concerning the identity of each of the
tables and columns and the directionality and cardinality of the
relationships between the tables...
...a query assistant **user interface** (''QAUI'') presenting to
the **user** a **selectable** table **set** of
selectable tables from among the tables in the database, a
selectable column **set** of **selectable** columns from
among the columns of each of said tables in the database, and a
selectable column operations **set** of **selectable**
column operations on the columns, from which the **user** may

select tables, columns, and column operations to construct a
database query for said database in an intermediate **query
language,** said QAUI further accepting from the **user
selections** of tables, columns, and column operations...

...QAUI to receive from said QAUI the identity of each table, column, or
column operation **selected** by the **user,** said QAES
returning to the QAUI after each **selection** by the **user** an
updated version of said **selectable** table **set,** said
selectable column **set,** and said **selectable** column
operations **set,** said QAES excluding from said **selectable
sets** any table, column, or column operation which, if
selected by the **user,** would, based on the then-current
state of the database query and said conceptual information...
...Title Terms: **USER;**
International Patent Class (Main): **G06F-017/30**
International Patent Class (Additional): **G06F-012/00...**

...**G06F-017/27**
Manual Codes (EPI/S-X): **T01-J05B3...**

...**T01-J16A**

US005812840A

[54] DATABASE QUERY SYSTEM

[75] Inventor: Steven P. Shwartz, Orange, Conn.

[73] Assignee: Speedware Ltee./Ltd., Toronto, Canada

[*] Notice: The term of this patent shall not extend beyond the expiration date of Pat. No. 5,584,024.

[21] Appl. No.: 723,962

[22] Filed: Sep. 26, 1996

Related U.S. Application Data

[63] Continuation of Ser. No. 217,099, Mar. 24, 1994, Pat. No. 5,584,024.

[51] Int. Cl.^{6} G06F 17/30
[52] U.S. Cl. 395/604; 395/757; 395/922
[58] Field of Search 395/601, 604, 395/752, 757, 922

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Primary Examiner—Thomas G. Black
Assistant Examiner—Jack M. Choules
Attorney, Agent, or Firm—Morgan, Lewis & Bockius LLP; C. Scott Talbot

[57] **ABSTRACT**

A database query system includes a query assistant that permits the user to enter only queries that are both syntactically and semantically valid (and that can be processed by an SQL generator to produce semantically valid SQL). Through the use of dialog boxes, a user enters a query in an intermediate English-like language which is easily understood by the user. A query expert system monitors the query as it is being built, and using information about the structure of the database, it prevents the user from building semantically incorrect queries by disallowing choices in the dialog boxes which would create incorrect queries. An SQL generator is also provided which uses a set of transformations and pattern substitutions to convert the intermediate language into a syntactically and semantically correct SQL query. The intermediate language can represent complex SQL queries while at the same time being easy to understand. The intermediate language is also designed to be easily converted into SQL queries. In addition to the query assistant and the SQL generator, an administrative facility is provided which allows an administrator to add a conceptual layer to the underlying database making it easier for the user to query the database. This conceptual layer may contain alternate names for columns and tables, paths specifying standard and complex joins, definitions for virtual tables and columns, and limitations on user access.

20 Claims, 26 Drawing Sheets
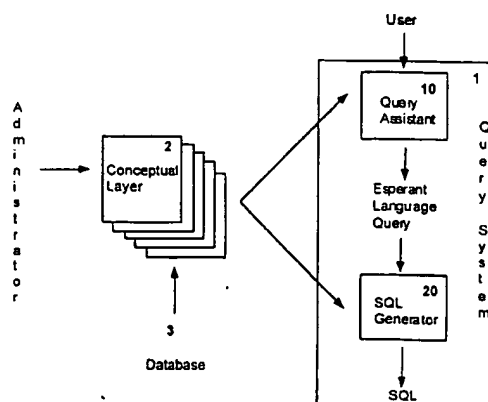
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No new patterns are matched, and there is only one more pattern to match which, when applied, yields:

SHOW CUSTOMERS.NAME OFENTITY!! CUSTOMERS WHERE ORDERS >=1 PRODUCTS ORDERBY CUSTOMERS. CITY

Steps 408–412 are not applicable to this query, since there is no CREATE VIEW command needed for this type of query. If it were one of a specific set of queries which require CREATE VIEW SQL syntax, SQL Generator 20 would be called recursively to create the views. Since CREATE VIEW is not necessary, no new words or phrases for conversion were introduced.

In step 414, the query is broken into SQL components. The query then becomes:

```
SELECT CUSTOMERS.NAME
WHERE ORDERS >= 1 PRODUCTS
ORDER BY CUSTOMERS.CITY
```

The LAST-ENTITY variable is set to CUSTOMERS, since the last table added to the select clause is from the table CUSTOMERS. The OFENTITY!! keyword introduced in the last pattern match is helpful in determining the LAST-ENTITY.

In steps 416–418, the Where clause "ORDERS >=1 PRODUCTS" is applied to the patterns shown above, resulting in one match, with pattern 605. By applying this pattern, the internal SQL structure for the query becomes:

```
SELECT CUSTOMERS.NAME
FROM CUSTOMERS T1
WHERE EXISTS
    SELECT *
    FROM PRODUCTS T2
    WHERE EXISTS
        SELECT *
        FROM ORDERS T3
        JOIN PRODUCTS T2
        JOIN CUSTOMERS T1
ORDER BY CUSTOMERS.CITY
```

The LAST-ENTITY variable is assigned the table PRODUCTS.

In step 420, f there were any table names in the SELECT portion it would expand to include all of the tables columns. Also any virtual table would be expanded. Neither are present in this example, but are performed by simple substitution.

In step 422, any columns in the SORT BY portion are added to SELECT if not present. This step converts the SELECT portion of the internal SQL to:

SELECT CUSTOMERS.CITY, CUSTOMERS.NAME

The date conversion function of step 424 is not applicable, since there are no dates in this example. Similarly, there are no virtual columns for expansion in step 426.

If any aliases need to be specified to the FROM clause, they are made in step 428. This query created the alias during the application of the Where rules, and no other tables were added to the from clause. The aliases are then substituted into the other sections as well. The internal SQL becomes:

```
SELECT T1.CITY T1.NAME
FROM CUSTOMERS TI
WHERE EXISTS
```

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-continued

```
SELECT *
FROM PRODUCTS T2
WHERE EXISTS
    SELECT *
    FROM ORDERS T3
    JOIN PRODUCTS T2
    JOIN CUSTOMERS T1
    ORDER BY T1.CITY
```

In step 430, the ORDER BY clause is converted to:
ORDER BY 1

In step 432, required joins are computed from the internal SQL. They are represented here by the "JOIN Table Alias" statement, and indicates that those tables need to join the table listed in the FROM clause above it. From the prior discussion on join path calculations, the join paths created from the statements:

```
FROM ORDERS T3
JOIN PRODUCTS T2
JOIN CUSTOMERS T1
```

are [CUSTOMERS ORDERS] and [ORDERS LINE_ITEMS PRODUCTS].

Then, in steps 434 and 436, since the join path calculation introduced a new table, LINE_ITEMS, the table needs to be added to the FROM clause with an alias to make:

FROM ORDERS T3, LINE_ITEMS T4

In step 438, the joins are created and added to the where clause from the join paths and foreign key information in the conceptual layer to produce the following Where clause:

```
WHERE      T3.ORDER# = T4.ORDER#
    AND   T2.PRODUCT# = T4.PRODUCT#
    AND   T1.CUSTOMER# = T3.CUSTOMER#
```

Steps 440–444 are not applicable to this example. Finally, in step 446, the internal SQL structure, which is now represented as:

```
SELECT T1.CITY T1.NAME
FROM CUSTOMERS T1
WHERE EXISTS
    SELECT *
    FROM PRODUCTS T2
    WHERE EXISTS
        SELECT *
        FROM ORDERS T3, LINE_ITEMS T4
        WHERE T3.ORDER# = T4.ORDER#
            AND T2.PRODUCT# = T4.PRODUCT#
            AND T1.CUSTOMER# = T3.CUSTOMER#
    ORDER BY 1
```

is converted to textual SQL. The above representation of the internal SQL structure is in proper textual structure for a query. The process of conversion to the textual query from the internal structure is a trivial step of combining the clauses and running through a simple parser.

What is claimed is:

1. A database query system for interactively creating, with a user, a syntactically and semantically correct query for a relational database having a plurality of tables, each of said tables having a plurality of columns and having a predetermined relationship to another of said tables, said system comprising:

a conceptual layer manager storing conceptual information about the relational database, said conceptual

information including structural information concerning the identity of each of the tables and columns and the directionality and cardinality of the relationships between the tables;

a query assistant user interface ("QAUI") presenting to the user a selectable table set of selectable tables from among the tables in the database, a selectable column set of selectable columns from among the columns of each of said tables in the database, and a selectable column operations set of selectable column operations on the columns, from which the user may select tables, columns, and column operations to construct a database query for said database, said QAUI further accepting from the user selections of tables, columns, and column operations;

a query assistant expert ("QAES") coupled to said QAUI to receive from said QAUI the identity of each table, column, or column operation selected by the user, said QAES returning to the QAUI after each selection by the user an updated version of said selectable table set, said selectable column set, and said selectable column operations set, said QAES excluding from said selectable sets any table, column, or column operation which, if selected by the user, would, based on the then-current state of the database query and said conceptual information, produce a semantically incorrect query.

2. The system of claim 1 wherein said QAES includes:

a storage system for maintaining state information about the current state of a database query; and

a query expert logic system specifying to said QAUI said selectable sets by analyzing said state information maintained in said storage system and said conceptual information stored by said conceptual layer manager.

3. The system of claim 2 wherein the database includes at least three tables and if in the then-current state of said database query a first and second of the three tables are selected, the relationship between the first and second of the three tables is one-to-many, and the relationship between the first of the three tables and a third of the three tables is one-to-many, said query expert logic system excludes the third table from said selectable table set.

4. The system of claim 2 wherein said storage system includes:

a set of state variables; and

a set of access routines for adding, deleting, and modifying said state variables.

5. The system of claim 2 wherein said query expert logic system is composed of procedural logic.

6. The system of claim 2 wherein said query expert logic system is a rule-based expert system.

7. The system of claim 2 wherein said conceptual information further comprises one or more of the following: foreign keys, table join paths, table join expression for non-equijoins, virtual table definitions, virtual column definitions, table descriptions, column descriptions, hiddent tables, and hidden columns.

8. The system of claim 2 wherein said conceptual information further compriess table join expression for non-equijoins.

9. The system of claim 2 wherein if said current state of said database query includes an aggregate column operation on a colun in a first table, said query expert logic system excludes from said selectable table set any other of said tables that is more detailed than said first table or is joinable with said first table only through another more detailed table.

10. The system of claim 2 wherein said query expert logic system excludes from said selectable column set for any

selected one of said tables any numeric column for which, if in said current state of said database query an aggregate column operation is applied to another column based on said selected table or based on another table having a one-to-one relationship with said selected table.

11. The system of claim 2 wherein said conceptual information further comprises virtual column definitions.

12. The system of claim 1 wherein each of said selectable table set, said selectable column set, and said selectable column operation set is mutually exclusive to a corresponding nonselectable table set, nonselectable column set, and nonselectable column operation set and is a subset of all tables, columns, and column operations, respectively, maintained by said conceptual layer manager which the user may next select in building a semantically correct database query.

13. The system of claim 12 wherein said QAUI displays said nonselectable sets and visually differentiates said selectable sets from said nonselectable sets.

14. The system of claim 1 wherein said database query is constructed in an intermediate query language and further comprising a query generator coupled to said QAUI to receive from said QAUI a completed database query in said intermediate query language, said query generator converting said query from said intermediate query language into a target query language, different from said intermediate query language, by a set of successive transformations by pattern substitution, said pattern substitutions including at lease one pattern that produces a correlated subquery when said target language is Structured Query Language.

15. A method for interactively building a syntactically and semantically correct query of a relational database from selections by a user, said database having a plurality of tables, each of said tables having a plurality of columns and having a predetermined relationship to another of said tables, said method comprising the steps of:

presenting to the user a selectable table set of selectable tables from among the tables in the database;

receiving from the user a selection of a first one of said selectable tables;

presenting to the user a selectable column set of selectable columns from among the columns based on said first selected table;

receiving from the user a selection of a first selected column from among said selectable columns'

presenting to the user a selectable column operation set of selectable column operations applicable to said first selected column;

receiving from the user a selection of one of said column operations;

determining from the selected table, column, and column operation and from information about the directionality and cardinality of the relationships between the tables the identity of unallowed tables, columns, or column operations which would, if selected by the user would produce a semantically incorrect query;

generating a first updated version of said selectable table set, a first updated version of said selectable column set, and a first updated version of said selectable column operation set, said sets excluding said unallowed tables, columns, and column operations, respectively; and

presenting to the user said first updated version of selectable table set.

16. The method of claim 15 wherein:

the database includes at least three tables; and

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in said determining step if in the then-current state of said database query a first and second of the three tables are selected, the relationship between the first and second of the three tables is one-to-many, and the relationship between the first of the three tables and a third of the three tables is one-to-many, the third table is determined to be an unallowed table.

17. The method of claim 15 further comprising the steps of:

constructing from said user selections a database query in an intermediate query language;

converting said database query from said intermediate query language to a target query language, different from said intermediate query language.

18. The method of claim 17 wherein said target language is Structured Query Language and in said converting step

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said query is converted by a set of successive transformations by pattern substitution, said pattern substitutions including at lease one pattern that produces a correlated subquery.

19. The method of claim 15 wherein said selectable table set is mutually exclusive to a corresponding nonselectable table set and is a subset of all tables which the user may next select in building a semantically correct database query.

20. The method of claim 19 wherein said step of presenting to the user said first updated version of said selectable table set includes displaying said nonselectable sets and visually differentiating said selectable table set from said nonselectable table set.

*   *   *   *   *

US005584024A

United States Patent [19]

Shwartz

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[45] **Date of Patent:** Dec. 10, 1996

[54] **INTERACTIVE DATABASE QUERY SYSTEM AND METHOD FOR PROHIBITING THE SELECTION OF SEMANTICALLY INCORRECT QUERY PARAMETERS**

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[73] Assignee: **Software AG**, Germany

[21] Appl. No.: **217,099**

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364/274.2; 364/275.4; 364/283.3; 364/DIG. 1;
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364/419.01, 419.07, 974.6, 972.2, 274.2,
274.7, 275.1, 275.4, 283.3, 282.1

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Primary Examiner—Wayne Amsbury
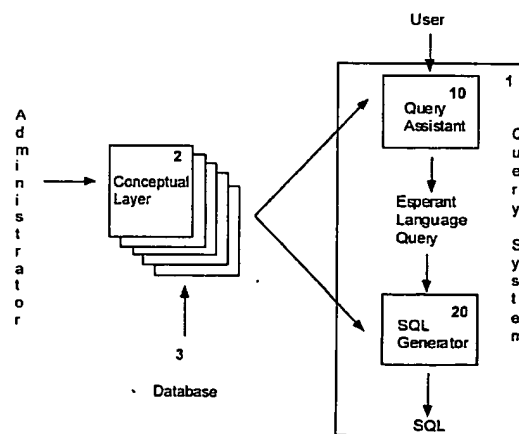*Assistant Examiner*—Jack M. Choules
Attorney, Agent, or Firm—Howrey & Simon; C. Scott Talbot; Thomas G. Woolston

[57] **ABSTRACT**

A database query system includes a query assistant that permits the user to enter only queries that are both syntactically and semantically valid (and that can be processed by an SQL generator to produce semantically valid SQL). Through the use of dialog boxes, a user enters a query in an intermediate English-like language which is easily understood by the user. A query expert system monitors the query as it is being built, and using information about the structure of the database, it prevents the user from building semantically incorrect queries by disallowing choices in the dialog boxes which would create incorrect queries. An SQL generator is also provided which uses a set of transformations and pattern substitutions to convert the intermediate language into a syntactically and semantically correct SQL query.

The intermediate language can represent complex SQL queries while at the same time being easy to understand. The intermediate language is also designed to be easily converted into SQL queries. In addition to the query assistant and the SQL generator, an administrative facility is provided which allows an administrator to add a conceptual layer to the underlying database making it easier for the user to query the database. This conceptual layer may contain alternate names for columns and tables, paths specifying standard and complex joins, definitions for virtual tables and columns, and limitations on user access.

27 Claims, 26 Drawing Sheets

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What is claimed is:

1. A database query system for interactively creating, with a user, a syntactically and semantically correct query for a relational database having a plurality of tables, each of said tables having a plurality of columns and having a predetermined relationship to another of said tables, said system comprising:

a conceptual layer manager for storing conceptual information about the relational database, said conceptual information including structural information concerning the identity of each of the tables and columns and the directionality and cardinality of the relationships between the tables;

a query assistant user interface ("QAUI") presenting to the user a selectable table set of selectable tables from among the tables in the database, a selectable column set of selectable columns from among the columns of each of said tables in the database, and a selectable column operations set of selectable column operations on the columns, from which the user may select tables, columns, and column operations to construct a database query for said database in an intermediate query language, said QAUI further accepting from the user selections of tables, columns, and column operations;

a query assistant expert ("QAES") coupled to said QAUI to receive from said QAUI the identity of each table, column, or column operation selected by the user, said QAES returning to the QAUI after each selection by the user an updated version of said selectable table set, said selectable column set, and said selectable column operations set, said QAES excluding from said selectable sets any table, column, or column operation which, if selected by the user, would, based on the then-current state of the database query and said conceptual information, produce a semantically incorrect query.

2. The database query system of claim 1 wherein said QAES includes:

a storage system for maintaining state information about the current state of a database query; and

a query expert logic system specifying to said QAUI said selectable sets by analyzing said state information maintained in said storage system and said conceptual information stored by said conceptual layer manager.

3. A database query system according to claim 2 wherein said storage system includes:

a set of state variables; and

a set of access routines for adding deleting and modifying said state variables.

4. A database query system according to claim 2 wherein said storage system includes:

a state database, said state database containing said state information; and

a set of database access routines for adding to, deleting from and modifying said state database.

5. A database query system according to claim 2 wherein said query expert logic system is composed of procedural logic.

6. A database query system according to claim 2 wherein said query expert logic system is a rule-based expert system.

7. A database query system according to claim 2 wherein said conceptual information further comprise one or more of the following: foreign keys, table join paths, table join expression for non-equijoins, virtual table definitions, virtual column definitions, table descriptions, column descriptions, hidden tables, and hidden columns.

8. A database query system according to claim 2 wherein said conceptual information further comprises table join expression for non-equijoins.

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9. The database query system of claim 2 wherein if said current state of said database query includes an aggregate column operation on a column in a first table, said query expert logic system excludes from said selectable table set any other of said tables that is more detailed than said first table or is joinable with said first table only through another more detailed table.

10. The system of claim 2 wherein the database includes at least four tables and wherein if in the then-current state of said database query two of the tables are selected, said query expert logic system excludes from said selectable table set any other of said tables that does not form in combination with said two selected tables a navigable set.

11. The system of claim 2 wherein said query expert logic system excludes from said selectable column set for any selected one of said tables any numeric column for which, if in said current state of said database query an aggregate column operation is applied to another column based on said selected table or based on another table having a one-to-one relationship with said selected table.

12. A database query system according to claim 2 wherein said conceptual information further comprises virtual column definitions.

13. A database query system according to claim 12 wherein said virtual column definition contains primary key and foreign key references to define a join operation.

14. A database query system according to claim 1 wherein said each of said selectable table set, said selectable column set, and said selectable column operation set is mutually exclusive to a corresponding nonselectable table set, nonselectable column set, and nonselectable column operation set and is a subset of all tables, columns, and column operations, respectively, maintained by said conceptual layer manager which the user may next select in building a semantically correct database query.

15. A database query system according to claim 14 wherein said QAUI indicates to the user said selectable sets set of permissible selections and not displaying said nonselectable sets.

16. A database query system according to claim 14 wherein said QAUI displays said nonselectable sets and visually differentiates said selectable sets from said nonselectable sets.

17. A database query system according to claim 16 wherein said QAUI visually differentiates said sets by color.

18. A database query system according to claim 16 wherein said QAUI visually differentiates said selectable and nonselectable sets by type characteristic.

19. The database query system of claim 1 further comprising a query generator coupled to said QAUI to receive from said QAUI a completed database query in said intermediate query language, said query generator converting said query from said intermediate query language into a target query language different from said intermediate query language.

20. A database query system according to claim 19 wherein said target query language is Structured Query Language (SQL).

21. A database query system according to claim 19 wherein said query generator converts said intermediate language query into said target language by a set of successive transformations.

22. A database query system according to claim 21 wherein at least one of said set of successive transformations is transformation by pattern substitution.

23. A database query system according to claim 21 wherein said set of transformations comprises:

a set of structural transformations;

a set of transformation to include inferred information; and

a set of transformations by pattern substitution.

24. A method for interactively building a syntactically and semantically correct query of a relational database from selections by a user, said database having a plurality of tables, each of said tables having a plurality of columns and having a predetermined relationship to another of said tables, said method comprising the steps of:

presenting to the user a selectable table set of selectable tables from among the tables in the database;

receiving from the user a selection of a first one of said selectable tables;

presenting to the user a selectable column set of selectable columns from among the columns based on said first selected table;

receiving from the user a selection of a first selected column from among said selectable columns;

presenting to the user a selectable column operation set of selectable column operations applicable to said first selected column;

receiving from the user a selection of one of said column operations; .

presenting to the user a first updated version of said selectable table set from which the user may select a second selected table, said selectable table set excluding any table that is more detailed than said first selected table on which said selected column operation is applied or that is joinable with said first selected table

only through a more detailed table if said selected column operation is an aggregate operation.

25. The method of claim 24 wherein said first updated version of said selectable table set further excludes any table that is not joinable with said first selected table.

26. The method of claim 25 wherein the database includes at least four tables and further comprising the steps of:

receiving from the user a selection of a second selected table from said first updated version of said selectable table set; and

presenting to the user a second updated version of said selectable table set from which the user may select a third selected table, said selectable table set excluding any table the selection of which does not form in combination with said first and second selected tables a navigable set.

27. The method of claim 26 further comprising the steps of:

receiving from the user a selection of a second selected table from said first updated version of said selectable table set; and

presenting to the user an updated version of said selectable column set based on said second selected table, said selectable columns set excluding any column that contains non-numeric information if said second selected table is the same as said first selected table or has a one-to-one relationship with said first selected table.

*   *   *   *   *

·Derwent WPIX

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012238721      **Image available**
WPI Acc No: 1999-044829/199904
Related WPI Acc No: 2002-033185
XRPX Acc No: N99-032727
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   **User interface creating method for constraint
   handling in product configurator software - involves using control
   palette for placing user control at predefined position on
   configurator screen display template and several parameters
   associated with user control**
```
Patent Assignee: BT SQUARED TECHNOLOGIES INC (BTSQ-N)
Inventor: GELLER S D; HEYDA M S; NEES R
Number of Countries: 001  Number of Patents: 001
Patent Family:
Patent No     Kind   Date     Applicat No    Kind   Date      Week
US 5844554    A   19981201   US 96718645      A   19960917  199904  B

Priority Applications (No Type Date): US 96718645 A 19960917
Patent Details:
Patent No  Kind Lan Pg   Main IPC     Filing Notes
US 5844554    A     -79 G06F-015/00
```

   **User interface creating method for constraint
   handling in product configurator software...**

   **...involves using control palette for placing user control at
   predefined position on configurator screen display template and
   several parameters associated with user control**

   ...Abstract (Basic): involves displaying a form creation region on the
      display. The form creation region includes a **configuration**
      screen display template for display of **user** controls and
      information in the **user** product **configuration** program. A
      control palette containing several control **objects** which upon
      activation, places and displays **corresponding user**
      control at predefined position on **configurator** screen display
      template. A developer command issued with respect to the **user**
      control causes a parameter **selector** window to be displayed...

   ...The window contains several parameters associated with the **user**
      control. One such parameter is **selected** and associated with the
      user control. The information relating to the **selected**
      parameter is retrieved from the memory where it is stored, and used
      corresponding to the **user** control for computing product
      configurations.

         . . .

...ADVANTAGE - Allows external data tables to be **updated** independent
   of **configuration** software. Enables **updating**
   configurator software by merely executing **SQL** queries,
   avoiding rewriting of program code. Features very rapid execution time
   and quick response. Features **builder** windows and editors that
   allow tree-type hierarchical navigation and double-click
   selection, and well-organised queries
Title Terms: **USER;**
International Patent Class (Main): **G06F-015/00**
Manual Codes (EPI/S-X): **T01-J12**

United States Patent [19]

Geller et al.

[54] **METHODS AND SYSTEMS FOR USER INTERFACES AND CONSTRAINT HANDLING CONFIGURATIONS SOFTWARE**

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[73] Assignee: **BT Squared Technologies, Inc.**, Atlanta, Ga.

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(List continued on next page.)

[57] **ABSTRACT**

A computer implemented method of generating a user product configuration program module from a development environment. The user product configuration program module includes user controls that allow user input of information for use in configuration computations. Methods are disclosed for creating and maintaining the logic for a configuration program module in the form of configuration parameters; creating and maintaining the visual controls and user interface; and linking created visual controls with underlying structure represented by the parameters. Parameters assume values and are indicated as valid or invalid through operation of constraints and queries. Further methods include display of parameter creation and selection windows, and query creation and selection windows for creating SQL queries to access data in external tables. Parameters, queries, and constraints can be displayed in expandable and collapsible hierarchies, and quickly utilized in formulas, queries, and logical expressions by clicking in an expanded hierarchical display. Dependencies occurring in the underlying configuration logic are evaluated in both the forward and reverse direction so as to provide for very fast execution of the resultant configuration program module when a user provides new data via a user control.

47 Claims, 49 Drawing Sheets
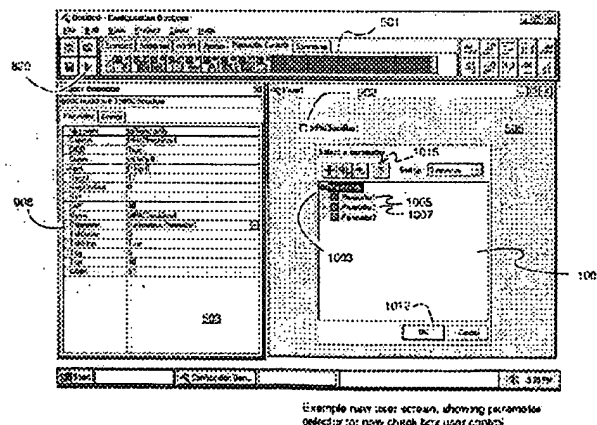
FIG. 45 illustrates the process 4500 executed within a compiled configuration program module to evaluate and apply constraints to a parameter, for purposes of determining whether constraints result in a valid or invalid condition for a parameter, as well as determine the values that a parameter can assume and the conditions for applying a constraint to parameter. It will be understood that the steps shown in FIG. 45 are executed within the compiled configuration program module each time that a user, not a developer, of the compiled configuration program module provides information input to the configuration program by activating a user control such as checking a check box, entering text into an edit box, activating a radio button, etc.

Starting at 4501, the first inquiry made is whether there are one or more constraints associated with a particular parameter whose value is being evaluated in response to input of information by a user of the compiled configuration program module. If so, control passes to process 4503. Process 4503 results in the execution of the conditions set forth in the "UsedWhen" property of the constraint. It will be recalled from previous discussion that the UsedWhen property, as shown at 1634 in FIG. 16A, determines conditions for applying constraints. These properties are shown for the various types of constraints at 1634 for a query constraint (FIG. 16A), 1734 for a discrete constraint (FIG. 17), 1834 for a string list constraint (FIG. 18), 2134 for a range constraint (FIG. 21), and 2310 for a formula constraint (FIG. 23). Further, it will be recalled from the discussion in conjunction with FIG. 20 how this conditional logic is employed by the UsedWhen property for application of a constraint to a parameter.

After process 4503, the inquiry is made at decision 4505 whether the particular constraint is used. If not, control passes to decision 4509, and the inquiry made whether there are more constraints associated with the particular parameter. It will be recalled that a number of constraints can be applied to a given parameter, such as the constraints 1622, 1624, 1626, all of which are applied to the selected Sound-System parameter 1608 in FIG. 16A. If there are more constraints, control passes to step 4513, the next constraint is evaluated, and control returns to process 4503. If there are no more constraints at this point, the no branch is taken from decision 4509 and the process executes. In this case, a Valid condition for the parameter being evaluated is indicated, as no constraints have been applied that might cause the condition to be invalid.

Return new to step 4505 where the inquiry is made whether a constraint is used. If so, the yes branch is taken to process 4520, and a validate process is executed to determine whether the application of the particular constraint validates the associated parameter or results in an invalid condition. This results from evaluation of any formulas or other computations in conjunction with the constraint.

At decision 4522, if the constraint validates the parameter, the yes branch is taken to step 4525, and the process exits with an indication of a valid condition. If not, the no branch is taken to decision 4530, and the inquiry is made whether there are more constraints to be applied to this parameter. If not, the no branch is taken to step 4535, and an invalid condition is indicated.

An invalid condition is indicated at this stage because it has been determined that a particular constraint is to be employed, but such constraint did not result in validating the associated parameter, and there are no more constraints that might further affect the validity or invalidity of the parameter. In this case, an invalid condition is indicated for the parameter.

However, if there are more constraints that might result in validation of the parameter, the yes branch is taken from decision 4530 to step 4537. Here, the next constraint is evaluated, and at process 4539 the UsedWhen condition of the constraint is evaluated in a manner as described above. At decision 4541, an inquiry is made whether the constraint being evaluated has been used. If not, the program loops back to 4530. If at decision 4541 the constraint is used, control passes to step 4544, and the Execute Validate process, similar to that of step 4520, is executed. The next step is 4548, and the inquiry made whether the constraint being evaluated validates the parameter. If so, the parameter assumes a valid condition at 4525 and the process exits. On the other hand, if not, control returns to step 4530 and further inquiries are made whether there are more constraints that might validate the parameter.

From the foregoing, those skilled in the art will appreciate that the preceding steps allow the very quick determination of the effect of the input of information via a user control in the executed configuration program module.

Finally, it will be understood that the preferred embodiment has been disclosed by way of example, and that other modifications may occur to those skilled in the art without departing from the scope and spirit of the appended claims.

What is claimed is:

1. A computer-implemented method of generating a user product configuration program module that executes product configuration computations, the user product configuration program module including user controls that allow user input of information for use in the product configuration computations, the method being carried out in a developer environment on a computer system including a display, a command entry device, a pointing device, and a memory, comprising the steps of:

displaying a form creation region on the display, the form creation region comprising a configurator screen display template for display of user controls and information in the user product configuration program;

displaying a controls palette comprising a plurality of selectable controls in a palette region on the display, each of the plurality of controls comprising a control object that, when activated by a command, causes placement and display of a user control corresponding to the activated control at a predetermined position on the configurator screen display template;

in response to a developer command in the developer environment with respect to the user control on the configurator screen display template, displaying a parameter selector window, the parameter selector window comprising a plurality of parameters that can be associated with the user control, information associated with the parameters being stored in the memory;

in response to a developer command in the developer environment with respect to a selected one of the plurality of parameters in the parameter selector window, associating the selected parameter with the user control,

whereby information stored in memory associated with the selected parameter, as influenced by a user command with respect to the user control as utilized in the user product configuration program, is utilized in product configuration computations.

2. The computer-implemented method of claim 1, further comprising the step of displaying an object inspector region on the display comprising information corresponding to one or more predetermined properties of the user control, the

predetermined properties including a parameter property corresponding to information utilized in a product configuration computation; and

wherein the parameter selector window is displayed in response to selection of the parameter property of the user control in the object inspector region on the display by the developer.

3. The computer-implemented method of claim 1, further comprising the steps of:

in response to a compile command by the developer, compiling program code corresponding to the user control and information corresponding to the selected parameter into an executable configuration program module; and

providing the executable configuration program module as an output file, the executable configuration program module, when installed on a computer system, causing display of a user configurator screen display corresponding to the configurator screen display template and including the user control.

4. The computer-implemented method of claim 1, wherein the selected parameter is operative to assume an invalid state in response to predetermined conditions during execution of the configuration program module,

whereby during execution of the configuration program module, in response to user input of information via the user control that causes the predetermined condition, providing a visual indication of the invalid condition on the user configurator screen display.

5. The computer-implemented method of claim 1, wherein the parameter selector window includes a control that expands the display of parameters into a hierarchical arrangement of parameters.

6. The computer-implemented method of claim 5, wherein the parameter selector window includes a parameter group icon that when activated by a developer command, expands to display one or more parameters or other group icons associated with the parameter group icon.

7. The computer-implemented method of claim 1, wherein the control object in the controls palette is activated by the developer pointing to a predetermined region on the controls palette and activating a command entry device.

8. The computer-implemented method of claim 1, further comprising the steps of:

in response to a developer command in the developer environment, displaying a parameter explorer window on the display, the parameter explorer window comprising an editable display of the plurality of selectable parameters stored in the memory.

9. The method of claim 8, further comprising the step of allowing developer modification of a parameter property by direct data entry with the command entry device.

10. The computer-implemented method of claim 9, wherein the developer modification of the selected parameter property comprises use of a formula.

11. The computer-implemented method of claim 10, wherein the modification of the selected parameter property by use of a formula comprises the steps of:

in response to a developer command in the developer environment, displaying a parameter explorer window on the display, the parameter explorer window comprising an arrangement of the plurality of selectable parameters;

in response to selection of one of the parameters in the parameter explorer window, displaying the selected parameter's properties in an object inspector region on

the display, the parameter properties including a formula property;

in response to selection of the formula property, displaying a formula builder window on the display;

allowing developer entry of a functional expression for assignment of a data value to the selected parameter.

12. A computer-implemented method for creating and maintaining configuration logic for a user product configuration program module that executes product configuration computations, the method being carried out in a developer environment on a computer system including a display, a command entry device, a pointing device, and a memory, comprising the steps of:

displaying a user control in the developer environment, the user control allowing user input of information for use in the configuration computations during execution of the user product configuration program module;

in response to a first developer command in the developer environment, displaying a parameter window on the display, the parameter window comprising a developer interface for creating and maintaining a plurality of selectable parameters stored in the memory;

in response to a second developer command in the developer environment with respect to a selected parameter in the parameter window, displaying a constraint item in association with the selected parameter, the constraint item being operative for constraining information that can be associated with the selected parameter; and

in response to a third developer command in the developer environment, associating a selected parameter with the user control,

whereby the user product configuration program module is operative during execution of the user product configuration program module to receive user input via the user control and apply information associated with the parameter associated with the user control in the product configuration computations.

13. The computer-implemented method of claim 12, wherein the constraint item comprises a query type constraint.

14. The computer-implemented method of claim 13, further comprising the steps of:

in response to a fourth developer command in the developer environment with respect to the query type constraint, displaying a query selector window on the display, the query selector window comprising a plurality of queries that can be associated with parameter information associated with the query being obtained from an external database and stored in the memory.

15. The computer-implemented method of claim 14, further comprising the steps of:

in response to a fifth developer command in the developer environment, displaying a query explorer window on the display, the query explorer window comprising a developer interface for creating and maintaining a plurality of selectable queries in the memory, each of the queries comprising a set of external database commands executable by the computer for retrieving information from an external database for association with one of the plurality of parameters.

16. The computer-implemented method of claim 15, wherein the query selector window and the query explorer window comprise separate windows.

17. The computer-implemented method of claim 12, wherein the constraint item is one or more selected from the

group comprising: a query constraint, a discrete constraint, a formula constraint, a string list constraint, and a range constraint.

18. The computer-implemented method of claim 12, further comprising the steps of:

displaying a parameter group icon in the parameter window, the parameter group icon being associated with an organizational concept with which a plurality of parameters can be associated; and

in response to a developer command with respect to the parameter group icon, expanding the display of parameters associated with the parameter group icon, thereby allowing creation of a hierarchically expandable and collapsible display of parameters.

19. The computer-implemented method of claim 12, wherein the parameter window is a parameter explorer window, and further comprising the steps of:

in response to a sixth developer command in the developer environment, displaying a parameter selector window on the display, the parameter selector window comprising a developer interface for associating a selected parameter stored in the memory with a selected user control.

20. The computer-implemented method of claim 19, wherein the parameter selector window is displayed in response to a developer command with respect to a control placed on a configuration screen display template, and further comprising the steps of:

displaying a form creation region on the display, the form creation region comprising a configurator screen display template for display of user controls and information in the user product configuration program;

displaying a controls palette comprising a plurality of selectable controls in a palette region on the display, each of the plurality of controls comprising a control object that, when activated by a command, causes placement and display of a user control corresponding to the activated control on the configurator screen display template; and

in response to a developer command with respect to a selected parameter in the parameter selector window, associating the selected parameter with the user control on the configurator screen display template,

whereby information stored in memory associated with the selected parameter, as influenced by a user command with respect to the user control as utilized in the user product configuration program, is utilized in product configuration computations.

21. The computer-implemented method of claim 20, further comprising the step of displaying an object inspector region on the display comprising information corresponding to one or more predetermined properties of the user control on the configurator screen display template, the predetermined properties including a parameter property corresponding to information utilized in a product configuration computation; and

wherein the parameter selector window is displayed in response to selection of the parameter property of the user control in the object inspector region on the display by the developer.

22. The computer-implemented method of claim 12, further comprising the steps of:

in response to a compile command by the developer, compiling program code corresponding to the user control and information corresponding to the selected parameter into an executable configuration program module; and

providing the executable configuration program module as an output file, the executable configuration program module, when installed on a computer system, causing display of a user configurator screen display including the user control.

23. The computer-implemented method of claim 12, wherein the selected parameter is operative to assume an invalid state in response to predetermined conditions during execution of the configuration program module,

whereby during execution of the configuration program module, in response to user input of information via the user control that causes the predetermined condition, providing a visual indication of the invalid condition on the user configurator screen display.

24. A method of generating a user product configuration program module that executes product configuration computations, the method being carried out in a developer environment on a computer system including a display, a command entry device, a pointing device, and a memory, comprising the steps of:

displaying a form creation region on the display, the form creation region comprising a configurator screen display template for display of user controls and information in the user product configuration program module;

displaying a controls palette comprising a plurality of selectable controls in a palette region on the display, each of the plurality of controls comprising a control object that, when activated by a command by a developer, causes placement and display of a user control corresponding to the activated control on the configurator screen display template;

displaying an object inspector region on the display comprising information corresponding to one or more predetermined properties of the user control, the predetermined properties including a parameter property corresponding to information utilized in a product configuration computation;

in response to selection of the parameter property of the user control in the object inspector region on the display by the developer, displaying a parameter selector window on the display, the parameter selector window comprising an arrangement of a plurality of selectable parameters;

in response to selection of one of the selectable parameters in the parameter selector window by the developer, associating information corresponding to the selected parameter with the user control;

in response to a compile command by the developer, compiling program code corresponding to the user control and information corresponding to the selected parameter into an executable configuration program module; and

providing the executable configuration program module as an output file, the executable configuration program module, when installed on a computer system, causing display of a user configurator screen display corresponding to the configurator screen display template and including the user control.

25. The method of claim 24, wherein the control control in the controls palette is activated by pointing to a predetermined region on the controls palette and activating the command entry device.

26. The method of claim 24, further comprising the step of allowing developer modification of the selected parameter property by direct data entry with the command entry device.

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27. The method of claim 24, wherein the arrangement of the plurality of selectable parameters displayed in the parameter selector window comprises a hierarchical display.

28. The method of claim 24, further comprising the step of allowing developer modification of a parameter property of the user control on the configurator screen display template.

29. The method of claim 28, wherein the developer modification of the selected parameter property comprises use of a formula.

30. The method of claim 29, wherein the modification of the selected parameter property by use of a formula comprises the steps of:

in response to a developer command, displaying a parameter explorer window on the display, the parameter explorer window comprising an arrangement of the plurality of selectable parameters;

in response to selection of one of the parameters in the parameter explorer window, displaying the selected parameter's properties in the object inspector region on the display, the parameter properties including a formula property;

in response to selection of the formula property, displaying a formula builder window on the display;

allowing developer entry of a functional expression for assignment of a data value to the selected parameter.

31. The method of claim 30, wherein the formula builder window comprises a variables pane, a functions pane, and a formula pane; and

in response to selection by the developer of a function in the functions pane, utilizing the selected function as a part of a formula displayed in the formula pane.

32. The method of claim 31, wherein the functions in the functions pane are selected from the group comprising Assign, AssignWhen, AssignWhenElse, Lookup, Display, and DisplayWhen.

33. The method of claim 31, wherein one of the functions is a Lookup function, wherein an external data table including information for use in the function is stored in the memory, and further comprising the steps of:

displaying a function template in the formula pane, the function template comprising editable text string corresponding to a mathematical expression;

allowing developer editing of the function template in the formula pane;

displaying a hierarchical arrangement of parameters and queries in the variables pane;

in response to selection of a query in the variables pane with the pointing device, displaying the selected query in the formula pane; and

allowing developer editing of selected query.

34. A method of representing a constraint to a configuration parameter in the user interface of a configuration program module that executes product configuration computations, the method being carried out in a developer environment on a computer system including a display, a command entry device, and a memory, comprising the steps of:

storing a plurality of selectable parameters in a predetermined arrangement in the memory, each of the parameters comprising information utilized in a product configuration computation;

displaying on the display a configurator screen display template for display of user controls and information in the configuration program module;

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displaying a user control on the screen display template, the user control being operative, in the configuration program module, for receiving user information input;

displaying an object inspector region on the display comprising information corresponding to properties of the user control displayed on the screen display template;

in response to developer selection of a parameters property of the user control in the object inspector region, displaying a first parameter window on the display for selection of a parameter, the first parameter window comprising a display of the plurality of selectable parameters stored in the memory;

in response to developer selection of one of the parameters in the first parameter window, associating the selected parameter with the user control in the memory;

in response to a first developer command, displaying a second parameter window on the display, the second parameter window comprising an editable display of the plurality of selectable parameters stored in the memory;

in response to a second developer command relating to a selected one of the parameters in the second parameter window, displaying a constraint item on the display in association with the selected parameter;

in response to a third developer command relating to the displayed constraint item, receiving developer input regarding the application of the constraint item to the selected parameter; and

in response to a developer compile command, providing the configuration program module as an output, the configuration program module being responsive, when run on a computer system, to information input by a user via the user control for determining application of the constraint to the selected parameter.

35. The method of claim 34, wherein the step of receiving developer input regarding application of the constraint item to the selected parameter comprises displaying properties of the constraint item in the object inspector region, the properties of the constraint item including a conditional application property, and

wherein in response to a developer command relating to the conditional application property, displaying a formula builder window on the display, the formula builder window comprising a region for developer input of a logical condition for application of the constraint represented by the constraint item to the parameter.

36. The method of claim 35, wherein the conditional application property is selected from the group comprising: UsedWhen, UsedWhenExcept, Display, DisplayWhen, and Lookup.

37. The method of claim 35, wherein the condition for application of the particular constraint represented by the constraint item to the parameter comprises a Boolean function that evaluates to true or false at run time, based on user input.

38. The method of claim 34, wherein the constraint item is selectable from the group comprising a discrete constraint, a range constraint, a string constraint, a query constraint, and a formula constraint.

39. The method of claim 34, wherein first parameter window for displaying the plurality of selectable parameters includes a parameter group icon, the parameter group icon comprising display for organizing a collection of parameters into a logical grouping on the display.

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40. The method of claim **34**, wherein the first parameter window and the second parameter window are separate windows on the display.

41. The method of claim **34**, further comprising the steps of:

   displaying a constraint list icon in association with the selected parameter in the second parameter window:

   in response to developer selection of the constraint list icon in the parameter explorer window, displaying a constraint menu comprising a plurality of selectable constraint type controls;

   in response to developer selection of one of the selectable constraint type controls, displaying a new constraint item in association with the selected parameter in the second parameter window.

42. The method of claim **34**, further comprising the steps of:

   in response to developer selection of a constraint item in the second parameter window, displaying properties of the selected constraint item in the object inspector region; and

   allowing developer modification of the properties of the selected constraint item in the object inspector region.

43. The method of claim **34**, further comprising the steps of:

   receiving user input of information into a computer system running the configuration program module via the user control;

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   in response to determination that the information provided by the user input violates a constraint represented by a constraint item, providing a visual indication of the constraint violation on the user configurator screen display.

44. The method of claim **43**, wherein the visual indication of the constraint violation comprises display of a red indication on a displayed user control associated with a parameter including the violated constraint.

45. The method of claim **44**, wherein the visual indication of the constraint violation is provided on all user controls that are associated with the parameter having the violated constraint.

46. The method of claim **34**, wherein the selected parameter is a first parameter whose value is dependent on the value of a second parameter,

   wherein the assignment of a value to the first parameter comprises evaluation of a formula containing the second parameter,

   wherein the value of the second parameter is a function of a user control, and

   wherein information input by a user via the user control is reflected in real time to the first parameter.

47. The method of claim **34**, wherein a plurality of constraint items may be associated with the selected parameter, and wherein the effect of the plural constraint items on the parameter is additive.

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015624881   **Image available**
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Related WPI Acc No: 2005-793721
XRPX Acc No: N03-548739
  Machine-readable medium for managing functions that express
  business rules to allow calling, maintaining functions and providing an
  execution framework for functions
Patent Assignee: DAMAN INC (DAMA-N)
Inventor: GHATATE B
Number of Countries: 001 Number of Patents: 001
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 6338069	B1	20020108	US 98210340	A	19981211	200365	B

Priority Applications (No Type Date): US 98210340 A 19981211
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 6338069	B1		60	G06F-017/30	

Abstract (Basic):
         Function **metadata**' is stored and allows the dynamic
    execution of a function e.g. execution of a command. When the engine
    receives a request e.g. for a **set** of **fields,** an execution
    plan can be created by determining functions to be processed based upon
    their inputs and outputs identified by the function **metadata.**
         An **object** technology infrastructure is formed to store
    data and **metadata** fro the functions to be maintained.
    Metadata about a function can include data describing what a
    function does, a 'cost' associated with...

...function and the input and output parameters required by the function.
    The exposure of the **metadata** regarding the functions' input and
    output parameters allows an engine to track input/output
    relationship between the functions and therefore define the
    order of execution...

...allow calling functions, maintaining functions, providing an execution
    framework for functions e.g. C++ code, **set** of **SQL**
    statements to be maintained...

...Exposure of **metadata** regarding a functions' input and output
    parameters allows an engine to track input/output **relationship**
    between the functions and define order of execution...
...The drawing shows a block diagram illustrating the **relationship**
    between functions being tracked, **metadata objects,**
    execution **objects** and **manager objects** according to

an embodiment of the invention...

...**Metadata objects** (120...

...Execution **objects** (125...

...**Manager object** (130
International Patent Class (Main): **G06F-017/30**
Manual Codes (EPI/S-X): **T01-F05A**...

...**T01-J20B**...

...**T01-S03**

(12) **United States Patent**

Ghatate

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(45) Date of Patent: **Jan. 8, 2002**

(54) **METHOD AND APPARATUS FOR MANAGING FUNCTIONS**

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(73) Assignee: **Daman, Inc.**, Austin, TX (US)

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(52) U.S. Cl. 707/103
(58) Field of Search 707/103

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(57) **ABSTRACT**

A method and apparatus for managing functions (e.g., that express business rules) to allow calling functions, maintaining functions, and providing of an execution framework for functions. In one embodiment, there are a number of functions to be maintained. An object technology infrastructure is formed to store data and metadata for the functions. For example, metadata about a function can include data describing what that function does, a "cost" associated with that function, how to execute that function, the input and output parameters required by that function. The exposure of the metadata regarding the functions' input and output parameters allows an engine to track input/output relationships between the functions and, in essence, define the order of execution.

46 Claims, 21 Drawing Sheets
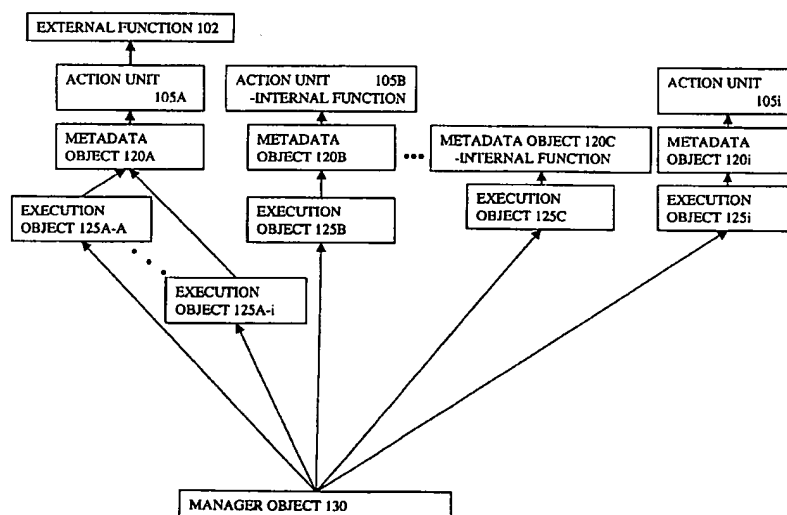
execute methods of the execution objects in the level objects of the EXECUTION_PATH_COLLECTION structure are applied in level order.

It should be noted that application of each EXECUTE method requires the operations illustrated in FIG. 19 to be performed. It should be understood that application of one or more of these EXECUTE methods could also result in a determination that the values for one or more required inputs of an underlying function is missing. In other words, this operation can be recursive in nature.

According to the embodiment shown in FIG. 20, for each key in the NEED_COLLECTION (see block 2015) it is verified that a value was acquired. In particular, block 2020 determines if a value was returned. If a value was not returned, control passes to block 2030 in which a flag is set for debugging purposes. However, if a value was returned, then the SET_PARAM method is applied with the returned value to associate the value with the EXECUTION object. While one embodiment is described in which verification that a value was returned for each key in the need collection is performed, alternative embodiments can avoid this verification and/or provide for this verification in other areas.

In this manner, values for any of the missing input parameters to the underlying function are located as part of applying the EXECUTE method from the EXECUTION object.

Relationship Objects

In one embodiment, the integration layer also includes two types of objects, base objects and relationship objects. The base objects describe the disparate integration sources A-i. The relationship objects describe relationships between the base objects, as well as relationships between the relationship objects themselves.

Specifically, the common concept of encapsulation in object technology is to place data with its' associated methods together in a single object. In contrast, the integration layer is developed such that those methods that express a relationship (referred to herein as "relationship methods") are placed in separate objects—the relationship objects. In this manner, a higher degree of encapsulation is provided. While relationship methods within relationship objects are similar to methods commonly found in object technology in that they can be applied to the objects that contain them, relationship methods within relationship objects are different in that they often are applied to other objects, including base objects and other relationship objects. For a further description of relationship objects, see "A Method and Apparatus for Providing Relationship Objects and Various Features to Relationship and Other Objects," filed Sep. 30, 1998, by Bhalchandra Ghatate, which is herein incorporated by reference.

As previously indicated, in one embodiment of the invention, relationship object(s) are used to provide for one or more SATISFY_RELATIONSHIP routines. In this manner, different types of satisfy relationship routines can be provided and selected from using the relationship object techniques.

Alternative Embodiments

While the invention has been described in terms of several embodiments, those skilled in the art will recognize that the invention is not limited to the embodiments described. The method and apparatus of the invention can be practiced with modification and alteration within the spirit and scope of the appended claims. The description is thus to be regarded as illustrative instead of limiting on the invention.

What is claimed is:

1. A machine readable medium having stored thereon:
a function requiring a set of one or more input parameters;
a first object including,
  a first structure storing a key for each of said set of input parameters, and
  an action method, which when applied by a processor, causes that processor to invoke an action unit; and
a second object including,
  a first structure to store data for identifying, for each of said set of input parameters, the corresponding key and a value for that input parameter,
  a second structure identifying said first object, and
  an execute method, which when applied by a processor, causes that processor to apply said action method;
said action unit including instructions, which when executed by a processor, cause that processor to,
  access said values, and
  invoke said function using said values as input parameters.

2. The machine readable medium of claim 1, wherein said function is contained within said action unit.

3. The machine readable medium of claim 1, wherein said first structure in said second object has stored therein, for at least one of said set of input parameters, the corresponding key and a value for that input parameter.

4. The machine readable medium of claim 1, wherein:
the data in said first structure of said second object identifies a third object; and
said third object includes a structure to store, for one or more of said set of input parameters, the corresponding key and value for that input parameter.

5. The machine readable medium of claim 1, wherein:
the data in said first structure of said second object identifies a third object; and
said third object includes,
  a first structure to store a plurality of context objects, each of said plurality of context objects to store, for one or more of said set of input parameters, the corresponding key and a value for that input parameter, and
  a second structure to store data identifying one of said plurality of context objects as a default context object.

6. The machine readable medium of claim 1, wherein:
said function provides a set of one or more output parameters;
said first structure in said first object also storing a key for each of said set of output parameters.

7. A machine readable medium having stored thereon:
a function requiring a set of one or more input parameters;
a first object including,
  a first structure storing a key for each of said set of input parameters, and
  an action method, which when applied by a processor, causes that processor to invoke said function; and
a second object including,
  a first structure identifying a third object,
  a second structure identifying said first object, and
  an execute method, which when applied by a processor, causes that processor to apply said action method; and
said third object including,
  a first structure to store a plurality of context objects, each of said plurality of context objects to store, for

one or more of said set of input parameters, the corresponding key and a value for that input parameter, and

a second structure to store data identifying one of said plurality of context objects as a default context object.

8. The machine readable medium of claim 7 further having stored thereon:

an action unit including said function as a method.

9. The machine readable medium of claim 7 further having stored thereon:

an action unit including instructions, which when executed by a processor, cause that processor to invoke said function; and

wherein said action method, when applied, causes the execution of the instructions in said action unit.

10. The machine readable medium of claim 7 further having stored thereon:

an action unit including instructions, which when executed by a processor, cause that processor to,

access values for one or more of said set of input parameters from a selected one of said plurality of context objects that is passed to said action unit, and

invoke said function using said values as input parameters; and

wherein said action method, when applied, causes said function to be invoked through the instructions in said action unit.

11. The machine readable medium of claim 10, wherein:

said second object also includes a first structure storing data identifying, for one or more of said set of input parameters, the corresponding key and a value for that input parameter; and

said instructions of said action unit, when executed by a processor, also cause that processor to access values for one or more of said set of input parameters from said first structure.

12. The machine readable medium of claim 11, wherein:

said instructions of said action unit, when executed by a processor, also cause that processor to access values for one or more of said set of input parameters from said default context object.

13. The machine readable medium of claim 11, wherein:

said instructions of said action unit, when executed by a processor, also cause that processor to access values for one or more of said set of input parameters from said default context object.

14. The machine readable medium of claim 7 further having stored thereon:

said second object also includes a first structure to store data identifying, for one or more of said set of input parameters, the corresponding key and a value for that input parameter; and

an action unit including instructions, which when executed by a processor, cause that processor to,

access values for said set of input parameters from said first structure of said second object, a selected one of said plurality of context objects that is passed to said action unit, and said default context object, and

invoke said function using said values as input parameters.

15. A machine readable medium having stored thereon:

a plurality of functions for one or more applications, each of said plurality of functions requiring one or more input parameters, the input parameters required by said

plurality of functions collectively defining a set of parameter kinds irrespective of data type, each parameter kind in said set being assigned a unique key;

a metadata object corresponding to each of said plurality of functions, each said metadata object storing data to locate the corresponding one of said plurality of functions, each said metadata object also storing the unique key for each input parameter required by the corresponding one of said plurality of functions; and

each metadata object having one or more corresponding execution objects, each execution object including a structure storing data to identify a value for each input parameter of the one of said plurality of functions identified by the corresponding metadata object.

16. The machine readable medium of claim 15, wherein each parameter of said plurality of functions is of one of a plurality of data types each supporting a range of values, wherein different data is categorized irrespective of data type, and wherein each category of data defines one of the parameter kinds.

17. The machine readable medium of claim 15 further having stored thereon:

an action unit for each of said plurality of functions, each of said action units including instructions, which when executed by a processor, cause that processor to invoke the corresponding one of said plurality of functions; and

wherein the data in each of said metadata objects for locating the corresponding one of said plurality of functions identifies the corresponding one of said action units; and

wherein each of said metadata objects includes an action method, which when applied, causes the execution of the instructions in the corresponding one of said action units.

18. The machine readable medium of claim 17, wherein:

each execution object includes a method, which when applied, causes said action method of the corresponding metadata object to be applied for that business rule;

each action method, when applied responsive to the method of an execution object, causes the instructions in the corresponding action unit to be executed for that business rule; and

the instructions in each action unit, when applied responsive to an action method responsive to the method of an execution object, causes the values of that business rule to be accessed and said function of that action unit to be invoked with said values as input parameters.

19. The machine readable medium of claim 15, wherein at least one of said execution objects stores the key and a corresponding value for at least one input parameter of the corresponding one of said plurality of functions.

20. The machine readable medium of claim 15, wherein:

at least one of said execution objects includes a structure identifying a manager object; and

said manager object includes a structure to store the key and a corresponding value for at least on input parameter of the one of said plurality of functions corresponding to the at least one of said execution objects.

21. The machine readable medium of claim 15, wherein:

at least one of said execution objects includes a structure identifying a manager object; and

said manager object includes,

a first structure to store a plurality of context objects, each of said plurality of context objects to store the

key and a corresponding value for one or more input parameters to at least certain of the plurality of functions, and

a second structure to store data identifying one of said plurality of context objects as a default context object.

22. The machine readable medium of claim 15, wherein:

each of said plurality of functions provides one or more output parameters, the input and output parameters of said plurality of functions collectively defining said set of parameter kinds irrespective of data type;

each said metadata object also storing the unique key for each input and output parameter of the corresponding one of said plurality of functions.

23. A machine readable medium having stored thereon sequences of instructions, which when executed by a set of one or more processors, cause said set of one or more processors to perform the acts of:

applying a first method from a first execution object, said first execution object identifying a first metadata object corresponding to a first function, said first function requiring one or more input parameters, said first metadata object storing data describing each input parameter of said first function, said first method causing the acts of,

accessing the data describing each input parameter of said first function from said first metadata object;

attempting to match values associated with said first execution object to each input parameter of said first function as described by the data; and

determining a value is missing for at least a first input parameter to said first function.

24. The machine readable medium of claim 23, wherein said first method further causes the acts of:

locating a second metadata object corresponding to a second function having one or more output parameters, said second metadata object storing data describing each output parameter of said second function;

determining the missing first input parameter is an output parameter of said second function; and

executing said second function to acquire the missing value.

25. The machine readable medium of claim 24, wherein said first method further causes the act of:

associating the acquired value with said first execution object; and

executing said first function using the acquired value now associated with said first execution object as the first input parameter.

26. The machine readable medium of claim 23, wherein said first method further causes the acts of:

determining a set of metadata objects that each have stored therein data describing an output parameter that matches one or more of the missing input parameters, wherein each metadata object in said set corresponds to a different function having a set of one or more output parameters, each metadata object in said set storing data describing said set of output parameters for the corresponding function; and

executing the functions corresponding to the set of metadata objects to acquire said set of missing values;

associating the acquired values with the first execution object; and executing said first function using the values associated with the first execution object as input parameters.

27. The machine readable medium of claim 23, wherein said attempting further includes:

accessing a structure in said first execution object, said structure in said first execution object to store values for one or more said set of input parameters to said first function.

28. The machine readable medium of claim 23, wherein said attempting further includes:

accessing a structure in a manager object identified by a structure in said first execution object, said structure in said manager object identifying a default one of a plurality of context objects, each of said plurality of context objects to store values for one or more of the input parameters to said first function; and

accessing said values from said default context object.

29. The machine readable medium of claim 23, wherein said sequences of instructions, when executed, cause said set of processors to further perform the acts of:

applying a first method from a second execution object, said second execution object also identifying said first metadata object, said first method of said second execution object causing the acts of,

accessing the data describing each input parameter of said first function from said first metadata object;

attempting to match values associated with said second execution object to each input parameter of said first function as described by the data in said first metadata object; and

determining one or more values are missing for at least certain input parameters of said first function.

30. A machine readable medium having stored thereon sequences of instructions, which when executed by a set of one or more processors, cause said set of one or more processors to perform the acts of:

applying a first method from a first execution object, said first execution object identifying a first of a plurality of metadata objects, each of said plurality of metadata objects identifying a different function, each of said functions having input and output parameters, wherein one or more parameters for different functions are the same, the parameters for the different functions collectively defining a set of parameter kinds, each parameter kind in said set of parameter kinds being assigned a unique key, each of said plurality of metadata objects storing the unique keys assigned the input and output parameters of the function they identify, said first method causing the acts of,

accessing the key for each input parameter stored in said first metadata object;

attempting to match parameter values associated with said first execution object to each of the accessed keys; and

determining parameter values are missing for a set including at least one of the accessed keys.

31. The machine readable medium of claim 30, wherein each parameter of said functions is of one of a plurality of data types each supporting a range of values, wherein different data is categorized irrespective of data type, and wherein each category of data defines one of the set of parameter kinds.

32. The machine readable medium of claim 30, wherein said first method further causes the acts of:

locating a set of one or more of said plurality of metadata objects that collectively store each of the set of keys; and

executing the functions corresponding to the set of metadata objects to acquire said set of missing parameter values as outputs of the functions;

associating the acquired parameter values with the first execution object; and

executing the function identified by the first metadata object using the parameter values associated with the first execution object as input parameters.

33. The machine readable medium of claim 30, wherein said attempting further includes:

accessing a structure in said first execution object, said structure in said first execution object to store values for one or more said set of input parameters to the function identified by the first metadata object.

34. The machine readable medium of claim 30, wherein said attempting further includes:

accessing a structure in a manager object identified by a structure in said first execution object, said structure in said manager object identifying a default one of a plurality of context objects, each of said plurality of context objects to store values for one or more of the input parameters to said first function; and

accessing one or more of said values from said default context object.

35. The machine readable medium of claim 30, wherein said sequences of instructions, when executed, cause said set of processors to further perform the acts of:

applying a first method from a second execution object, said second execution object also identifying said first metadata object, said first method of said second execution object causing the acts of,

accessing the key for each input parameter stored in said first metadata object;

attempting to match values associated with said second execution object to each of the accessed keys; and

determining parameter values are missing for a set including at least one of the accessed keys.

36. The machine readable medium of claim 30, wherein said attempting further includes:

accessing from a first structure in said first execution object a value for a first input parameter to said function identified by said first metadata object;

accessing, from a first of a set of context objects that was passed, a value for a second input parameter to said first function, said set of context objects being stored in a first structure of a manager object, each of said set of context objects to store values for one or more of the input parameters to said first function, said manager object identifying one of said set of context objects as a default context object; and

accessing from said default context object a value for a third input parameter to said first function.

37. A machine readable medium having stored thereon sequences of instructions, which when executed by a set of one or more processors, cause said set of one or more processors to perform the acts of:

receiving a request to locate a function that provides a particular parameter kind as an output;

locating a metadata object having stored therein data identifying said particular parameter kind, said metadata object identifying a function and storing said data to indicate the particular parameter kind is an output parameter of said function; and

providing an execution object that identifies said metadata object, wherein said execution object includes,

structure to identify values for a set of one or more input parameters to said function, and

a method, which when applied, causes said function to be invoked using the values identified by said structure as input parameters.

38. The machine readable medium of claim 37, wherein:

said function has a plurality of parameters, said metadata object stores data identifying each kind of said plurality of parameters.

39. The machine readable medium of claim 38, wherein each parameter of said function is of one of a plurality of data types each supporting a range of values, wherein different data is categorized irrespective of data type, and wherein each category of data defines one of the parameter kinds.

40. The machine readable medium of claim 37, wherein:

said metadata object includes an action method, which when applied by a processor, causes said processor to invoke said function; and

said method in said execution object, when applied, causes said action method to be applied.

41. A machine readable medium having stored thereon sequences of instructions, which when executed by a set of one or more processors, cause said set of one or more processors to perform the acts of:

receiving a request to locate a function that provides a particular output parameter, wherein each of a plurality of metadata objects identify a different function having one or more output parameters, said output parameters for the different functions collectively defining a set of parameter kinds, each parameter kind in said set being assigned a unique key, each of said plurality of metadata objects storing the unique keys assigned the output parameters of the function they identify;

locating a first of said plurality of metadata objects that stores the unique key for the particular output parameter; and

creating an execution object that identifies said first metadata object, wherein said execution object includes,

a structure to identify values for a set of one or more input parameters to said function identified by said first metadata object, and

a method, which when applied, causes said function identified by said first metadata object to be invoked using the values identified by said structure as input parameters.

42. The machine readable medium of claim 41, wherein each parameter of said functions is of one of a plurality of data types each supporting a range of values, wherein different data is categorized irrespective of data type, and wherein each category of data defines one of the parameter kinds.

43. The machine readable medium of claim 41, wherein said sequences of instructions, when executed, cause said set of processors to further perform the acts of:

applying said method from said execution object, wherein both said input and output parameters for the different functions collectively define said set of parameter kinds, each of said plurality of metadata objects storing the unique keys assigned the input and output parameters of the function they identify, said method from said execution object causing the acts of,

accessing the key for each input parameter stored in said first metadata object;

attempting to match parameter values associated with said execution object to each of the accessed keys; and

executing the function identified by said first metadata object using the parameter values associated with the execution object as input parameters.

44. The machine readable medium of claim 41, wherein said attempting further causes the acts of:

determining parameter values are missing for a set including at least one of the accessed keys.

locating a set of one or more of said plurality of metadata objects that collectively store each of the set of keys for output parameters; and

executing the functions corresponding to the plurality of metadata objects to acquire said set of missing parameter values as outputs of the functions; and

associating the acquired parameter values with first execution object.

45. A machine readable medium having stored thereon sequences of instructions, which when executed by a set of one or more processors, cause said set of one or more processors to perform the acts of:

applying a first method from a first execution object, said first execution object identifying a first of a plurality of metadata objects, each of said plurality of metadata objects identifying a different function, each of said functions having input and output parameters, wherein one or more parameters for different functions are the same, said parameters for the different functions collectively defining a set of parameter kinds, each parameter kind in said set of parameter kinds being assigned

a unique key, each of said plurality of metadata objects storing the unique keys assigned the input and output parameters of the function they identify, said first method causing the acts of,

accessing the key for each input parameter stored in said first metadata object;

associating with said first execution object a value stored as part of a first of a set of context objects that was passed, said value stored for use as a first input parameter, said set of context objects being stored in a first structure of a manager object, each of said set of context objects to store values for one or more of the input parameters to said first function; and

executing said first function using the parameter values associated with the first execution object as input parameters.

46. The machine readable medium of claim 45, wherein said first method further causes the acts of:

associating with said first execution object a value stored as part of one of said set of context objects identified by said manager object as a default context object, said value stored for a second input parameter to said first function.

* * * * *

46/3,K/37 (Item 37 from file: 350) <u>Links</u>
Derwent WPIX

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015684733     **Image available**
WPI Acc No: 2003-746922/200370
Related WPI Acc No: 2005-272127; 2005-272128
XRPX Acc No: N03-598603
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**Schema generation method in computer system, involves
generating schema to represent hierarchies of inter-object
relationships between objects in data polyarchy, based on
attribute values of objects**

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Number of Countries: 001  Number of Patents: 002
Patent Family:
```

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 20030177139	A1	20030918	US 2001995415	A	20011126	200370	B
US 6944626	B2	20050913	US 2001995415	A	20011126	200560	

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Priority Applications (No Type Date): US 2001995415 A 20011126
Patent Details:
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Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 20030177139	A1		35	G06F-007/00	
US 6944626	B2			G06F-017/00	

**Schema generation method in computer system, involves
generating schema to represent hierarchies of inter-object
relationships between objects in data polyarchy, based on
attribute values of objects**

Abstract (Basic):
          A schema is generated to represent multiple hierarchies of
    inter-**object relationships** between various **objects**
    in data polyarchy, based on attribute values of the **objects.**
          3) polyarchical **query language** data structure...

...including personal computer (claimed), server computer, laptop computer,
    handheld computer, multiprocessor system, microprocessor based system,
    set-top box, programmable consumer electronics, minicomputer and
    mainframe computer connected to network such as Internet...

...Since the data polyarchy complex inter-**object
    relationships** are dynamically generated, the manual intervention
    from a network administrator to **configure** inter-**object
    relationship**, is not required...
...**optional** data servers (106...

...**client** computer (110...

...computer **executable** instructions (116...

...polyarchical data **set** (122
...Title Terms: **OBJECT**;
International Patent Class (Main): **G06F-007/00**...

...**G06F-017/00**
Manual Codes (EPI/S-X): **T01-F05E**...

...**T01-F07**...

...**T01-N02B1A**...

...**T01-S03**

(12) **United States Patent**
Cameron et al.

(10) **Patent No.:** **US 6,944,626 B2**
(45) **Date of Patent:** **Sep. 13, 2005**

(54) **DYNAMICALLY GENERATED SCHEMA REPRESENTING MULTIPLE HIERARCHIES OF INTER-OBJECT RELATIONSHIPS**

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(51) **Int. Cl.**[7] **G06F 17/00**
(52) **U.S. Cl.** .... **707/103 R**; 707/100; 707/102
(58) **Field of Search** .... 707/100, 102, 707/103 R, 103 Y

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(57) **ABSTRACT**

The described arrangements and procedures provide for interfacing (e.g., managing, presenting, etc.) with complex and often elastic inter-object relationships between objects in a data polyarchy. Specifically, a schema is dynamically generated by a server to represent multiple hierarchies of inter-object relationships between objects in a data polyarchy. The schema indicates or lists each attribute or element of interest in the data polyarchy. The schema further indicates any of one or more dimensions of inter-object relationships within which objects that comprise at least a subset of the listed the elements of interest participate. Thus, the schema indicates how to interface with the data polyarchy, which represents multiple hierarchies of inter-object relationships based on the values of attributes of the represented objects.

70 Claims, 14 Drawing Sheets
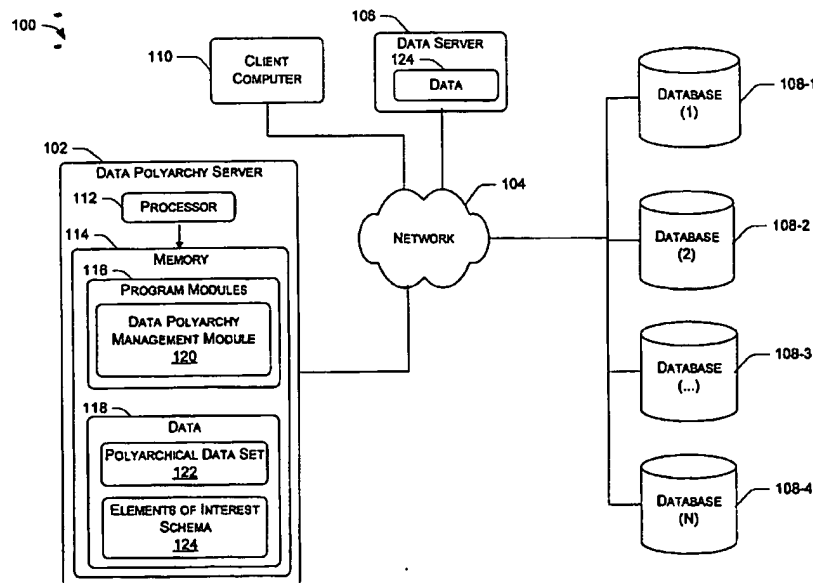
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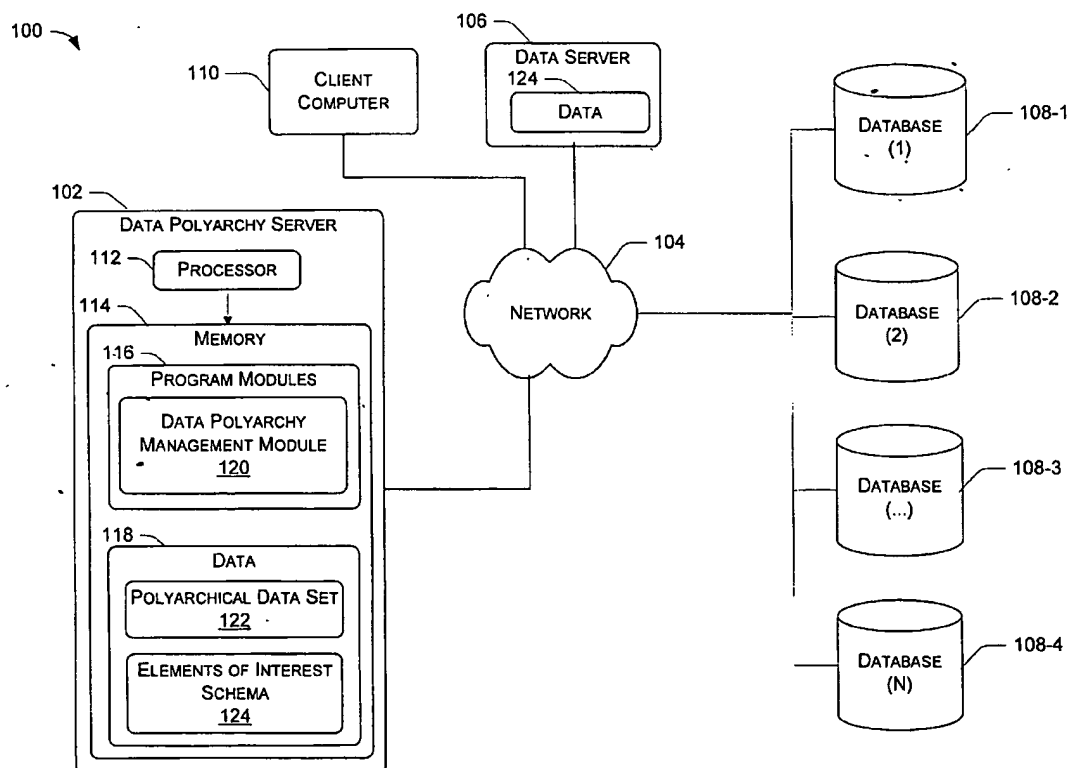
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(57) **ABSTRACT**

The described arrangements and procedures provide for
interfacing (e.g., managing, presenting, etc.) with complex
and often elastic inter-object relationships between objects
in a data polyarchy. Specifically, a schema is dynamically
generated by a server to represent multiple hierarchies of
inter-object relationships between objects in a data pol-
yarchy. The schema indicates or lists each attribute or
element of interest in the data polyarchy. The schema further
indicates any of one or more dimensions of inter-object
relationships within which objects that comprise at least a
subset of the listed the elements of interest participate. Thus,
the schema indicates how to interface with the data pol-
yarchy, which represents multiple hierarchies of inter-object
relationships based on the values of attributes of the repre-
sented objects.

objects. In this manner, complex real-world objects are represented with respect to the particular objects themselves, with respect to any set of decomposed sub-entities, or sub-objects that are related to the particular objects. These inter-object relationships are managed and navigated using a data polyarchy schema 124 that has been generated to access elements of interest in the data polyarchy 122.

[0155] Although the described subject matter to generate and manage polyarchies of data relationships has been described in language specific to structural features and/or methodological operations, it is to be understood that the subject defined in the appended claims is not necessarily limited to the specific features or operations described. Rather, the specific features and steps are disclosed as preferred forms of implementing the claimed present invention.

1. In a computer system, a method comprising dynamically generating a schema to represent multiple hierarchies of inter-object relationships between a plurality of objects in a data polyarchy, the schema being generated based on values of attributes of the objects.

2. A method as recited in claim 1, wherein the inter-object relationships comprise a flat relationship, a hierarchical relationship, and multiple intersecting hierarchies of relationships.

3. A method as recited in claim 1, where the objects comprise enterprise resource planning (ERP) objects, directory based objects, or database objects.

4. A method as recited in claim 1, wherein generating the schema further comprises:

identifying a plurality of attributes of interest based on values of attributes of the objects;

identifying one or more dimensions of inter-object relationships within which objects that comprise at least a subset of the attributes of interest participate.

5. A method as recited in claim 1, wherein the schema is designed to provide access control to organizational resources.

6. A method as recited in claim 1, further comprising communicating the schema to a client, the schema identifying how the client can access objects in the data polyarchy.

7. A method as recited in claim 1, further comprising:

receiving a request from a client based on the schema; and

responsive to receiving the request:

accessing an object of the objects based on the request;

transforming the object into transformed data that expresses any inter-object relationship between the object and any other object of the objects based on the request; and

issuing the transformed data to the client.

8. A method as recited in claim 7, wherein the transformed data expresses the inter-object relationships with respect to other objects in a same dimension or other objects in a different dimension, the same and/or the different dimension being indicated by the request.

9. A method as recited in claim 7, wherein the request comprises a limiting attribute to limit the transformed data by presenting the one or more objects only with respect to the limiting attribute.

10. A method as recited in claim 7, wherein the request queries for information corresponding to an object in the data polyarchy with respect to one or more particular dimensions.

11. A method as recited in claim 7, wherein the request comprises a dimension indicator to specify one or more hierarchies within which the data is to be presented in the transformed data.

12. A method as recited in claim 7, wherein the request further comprises a distinguishing attribute, a classifying attribute, or a locating attribute.

13. A method as recited in claim 7, wherein the request comprises a dimension information modifier to specify a particular direction and a particular depth to retrieve information from the data polyarchy.

14. A method as recited in claim 13, wherein the dimension information modifier is a siblings indication to retrieve all objects with a same parent as a current object in the data polyarchy.

15. A method as recited in claim 7, wherein the request indicates that at least one subset of the objects comprise a similar attribute; and wherein accessing one or more objects further comprises:

retrieving the one or more objects in a manner that is independent of any hierarchical data relationship between the data objects in the at least one subset.

16. A method as recited in claim 15, wherein the similar attribute comprises a logical domain selected from a distinguishing domain, a locating domain or a classifying domain.

17. A method as recited in claim 7, wherein the request corresponds to at least a first and second subset of the objects, the request comprising a logical modifier that specifies an operation, and wherein the method further comprises:

responsive to receiving the request, identifying at least a portion of the first and second subsets of directory objects in the polyarchical data set; and

wherein transforming the one or more objects further comprises performing the operation on the first and second subsets.

18. A method as recited in claim 17, wherein the logical modifier is a Boolean modifier.

19. A method as recited in claim 17, wherein the operation comprises any combination of filtering, union, intersection, join, and/or exclusion operations.

20. A method as recited in claim 7, wherein accessing the object further comprises accessing the object in a manner that is independent of any inter-object relationship between the object and any other object of the objects in a manner that is independent of any definition of a hierarchy in the data polyarchy.

21. A method as recited in claim 7, wherein accessing the object further comprises querying the data polyarchy for the object.

22. A method as recited in claim 7, wherein accessing the object further comprises managing, manipulating, or modifying the object or a relationship between the object and one of more of the other objects.

23. A computer-readable medium having computer-executable instructions comprising instructions for:

dynamically generating a schema to represent multiple hierarchies of inter-object relationships between a plurality of objects in a data polyarchy, the schema being

generated based on values of attributes of the objects, the schema indicating each attribute of interest in the data polyarchy, the schema further indicating any of one or more dimensions of inter-object relationships within which objects that comprise at least a subset of the attributes of interest participate.

24. A computer-readable medium as recited in claim 23, wherein the inter-object relationships comprise a flat relationship, a hierarchical relationship, and multiple intersecting hierarchies of relationships.

25. A computer-readable medium as recited in claim 23, where the objects comprise enterprise resource planning (ERP) objects, directory based objects, or database objects.

26. A computer-readable medium as recited in claim 23, further comprising computer-executable instructions for communicating the schema to a client to indicate how the client is to interface with the objects in the data polyarchy.

27. A computer-readable medium as recited in claim 23, further comprising computer-executable instructions for:

receiving a request from a client based on the schema;

responsive to receiving the request:

accessing at least one object in the data polyarchy based on the request;

transforming the at least one object into transformed data that expresses any inter-object relationship between the at least one object and any other objects of the objects based on the request; and

issuing the transformed data to the client.

28. A computer-readable medium as recited in claim 27, wherein the transformed data expresses the inter-object relationships with respect to other objects in a same dimension or other objects in a different dimension, the same and/or the different dimension being indicated by the request.

29. A computer-readable medium as recited in claim 27, wherein the request comprises a limiting attribute to limit the transformed data by presenting the one or more objects only with respect to the limiting attribute.

30. A computer-readable medium as recited in claim 27, wherein the request queries for information corresponding to an object in the data polyarchy with respect to one or more particular dimensions.

31. A computer-readable medium as recited in claim 27, wherein the request comprises a dimension indicator to specify one or more hierarchies within which the data is to be presented in the transformed data.

32. A computer-readable medium as recited in claim 27, wherein the request further comprises a distinguishing attribute, a classifying attribute, or a locating attribute.

33. A computer-readable medium as recited in claim 27, wherein the request comprises a dimension information modifier to specify a particular direction and a particular depth to retrieve information from the data polyarchy.

34. A computer-readable medium as recited in claim 33, wherein the dimension information modifier is a siblings indication to retrieve all objects with a same parent as a current object in the data polyarchy.

35. A computer-readable medium as recited in claim 27 wherein the request indicates that at least one subset of the objects comprise a similar attribute; and wherein the computer-executable instructions for accessing one or more objects further comprise instructions for:

retrieving the one or more objects in a manner that is independent of any hierarchical data relationship between the data objects in the at least one subset.

36. A computer-readable medium as recited in claim 35, wherein the similar attribute comprises a logical domain selected from a distinguishing domain, a locating domain or a classifying domain.

37. A computer-readable medium as recited in claim 27, wherein the request corresponds to at least a first and second subset of the objects, the request comprising a logical modifier that specifies an operation, and wherein the computer-executable instructions further comprise instructions for:

responsive to receiving the request, identifying at least a portion of the first and second subsets of directory objects in the polyarchical data set; and

wherein transforming the one or more objects further comprises performing the operation on the first and second subsets.

38. A computer-readable medium as recited in claim 37, wherein the logical modifier is a Boolean modifier.

39. A computer-readable medium as recited in claim 37, wherein the operation comprises any combination of filtering, union, intersection, join, and/or exclusion operations.

40. A computer-readable medium as recited in claim 27, wherein accessing the at least one object further comprises accessing the at least one object in a manner that is independent of any inter-object relationship between the at least one object and any other object of the objects in a manner that is independent of any definition of a hierarchy in the data polyarchy.

41. A computer-readable medium as recited in claim 27, wherein accessing the at least one object further comprises querying the data polyarchy for the at least one object.

42. A computer-readable medium as recited in claim 27, wherein accessing the at least one object further comprises managing, manipulating, or modifying the at least one object or a relationship between the at least one object and one of more different objects of the objects.

43. A computer comprising:

a memory comprising the computer-executable instructions; and

a processor coupled to the memory, the processor being configured to fetch and execute the computer-executable instructions for:

dynamically generating a schema to represent multiple hierarchies of inter-object relationships between a plurality of objects in a data polyarchy, the schema being generated based on values of attributes of the objects, the schema indicating each attribute of interest in the data polyarchy, the schema further indicating any of one or more dimensions of inter-object relationships within which objects that comprise at least a subset of the attributes of interest participate.

44. A computer as recited in claim 43, wherein the inter-object relationships comprise a flat relationship, a hierarchical relationship, and multiple intersecting hierarchies of relationships.

45. A computer as recited in claim 43, where the objects comprise enterprise resource planning (ERP) objects, directory based objects, or database objects.

46. A computer as recited in claim 43, wherein the computer-executable instructions further comprise instructions for communicating the schema to a client to indicate how the client is to interface with the objects in the data polyarchy.

47. A computer as recited in claim 43, wherein the computer-executable instructions further comprise instructions for:

receiving a request from a client based on the schema;

responsive to receiving the request:

accessing one or more objects in the data polyarchy based on the request;

transforming the one or more of the objects into transformed data that expresses any inter-object relationships based on the request; and

issuing the transformed data to the client.

48. A computer as recited in claim 47, wherein the transformed data expresses the inter-object relationships with respect to other objects in a same dimension or other objects in a different dimension, the same and/or the different dimension being indicated by the request.

49. A computer as recited in claim 47, wherein the request comprises a limiting attribute to limit the transformed data by presenting the one or more objects only with respect to the limiting attribute.

50. A computer as recited in claim 47, wherein the request queries for information corresponding to an object in the data polyarchy with respect to one or more particular dimensions.

51. A computer as recited in claim 47, wherein the request comprises a dimension indicator to specify one or more hierarchies within which the data is to be presented in the transformed data.

52. A computer as recited in claim 47, wherein the request further comprises a distinguishing attribute, a classifying attribute, or a locating attribute.

53. A computer as recited in claim 47, wherein the request comprises a dimension information modifier to specify a particular direction and a particular depth to retrieve information from the data polyarchy.

54. A computer as recited in claim 53, wherein the dimension information modifier is a siblings indication to retrieve all objects with a same parent as a current object in the data polyarchy.

55. A computer as recited in claim 47, wherein the request indicates that at least one subset of the objects comprise a similar attribute; and wherein the computer-executable instructions for accessing one or more objects further comprise instructions for:

retrieving the one or more objects in a manner that is independent of any hierarchical data relationship between the data objects in the at least one subset.

56. A computer as recited in claim 55, wherein the similar attribute comprises a logical domain selected from a distinguishing domain, a locating domain or a classifying domain.

57. A computer as recited in claim 47, wherein the request corresponds to at least a first and second subset of the objects, the request comprising a logical modifier that specifies an operation, and wherein the computer-executable instructions further comprise instructions for:

responsive to receiving the request, identifying at least a portion of the first and second subsets of directory objects in the polyarchical data set; and

wherein transforming the one or more objects further comprises performing the operation on the first and second subsets.

58. A computer as recited in claim 57, wherein the logical modifier is a Boolean modifier.

59. A computer as recited in claim 57, wherein the operation comprises any combination of filtering, union, intersection, join, and/or exclusion operations.

60. A computer as recited in claim 47, wherein accessing the one or more objects further comprises accessing the one or more objects in a manner that is independent of any inter-object relationship between the one or more objects and any other object of the objects in a manner that is independent of any definition of a hierarchy in the data polyarchy.

61. A computer as recited in claim 47, wherein accessing the one or more objects further comprises querying the data polyarchy for the one or more objects.

62. A computer as recited in claim 47, wherein accessing the one or more objects further comprises managing, manipulating, or modifying the one or more objects or a relationship between an object of the one or more objects and one of more different objects of the objects.

63. A computer comprising:

processing means for dynamically generating a schema to represent multiple hierarchies of inter-object relationships between a plurality of objects in a data polyarchy, the schema being generated based on values of attributes of the objects, the schema indicating each attribute of interest in the data polyarchy, the schema further indicating any of one or more dimensions of inter-object relationships within which objects that comprise at least a subset of the attributes of interest participate.

64. A computer as recited in claim 63, wherein the inter-object relationships comprise a flat relationship, a hierarchical relationship, and multiple intersecting hierarchies of relationships.

65. A computer as recited in claim 63, where the objects comprise enterprise resource planning (ERP) objects, directory based objects, or database objects.

66. A computer as recited in claim 63, further comprising processing means for communicating the schema to a client to indicate how the client is to interface with the objects in the data polyarchy.

67. A computer as recited in claim 63, further comprising processing means for:

receiving a request from a client based on the schema;

responsive to receiving the request:

accessing one or more objects in the data polyarchy based on the request;

transforming the one or more of the objects into transformed data that expresses any inter-object relationships based on the request; and

issuing the transformed data to the client.

68. A computer as recited in claim 67, wherein the transformed data expresses the inter-object relationships with respect to other objects in a same dimension or other

objects in a different dimension, the same and/or the different dimension being indicated by the request.

69. A computer as recited in claim 67, wherein the request comprises a limiting attribute to limit the transformed data by presenting the one or more objects only with respect to the limiting attribute.

70. A computer as recited in claim 67, wherein the request queries for information corresponding to an object in the data polyarchy with respect to one or more particular dimensions.

71. A computer as recited in claim 67, wherein the request comprises a dimension indicator to specify one or more hierarchies within which the data is to be presented in the transformed data.

72. A computer as recited in claim 67, wherein the request further comprises a distinguishing attribute, a classifying attribute, or a locating attribute.

73. A computer as recited in claim 67, wherein the request comprises a dimension information modifier to specify a particular direction and a particular depth to retrieve information from the data polyarchy.

74. A computer as recited in claim 73, wherein the dimension information modifier is a siblings indication to retrieve all objects with a same parent as a current object in the data polyarchy.

75. A computer as recited in claim 67, wherein the request indicates that at least one subset of the objects comprise a similar attribute; and wherein the means for accessing one or more objects further comprise means for:

retrieving the one or more objects in a manner that is independent of any hierarchical data relationship between the data objects in the at least one subset.

76. A computer as recited in claim 75, wherein the similar attribute comprises a logical domain selected from a distinguishing domain, a locating domain or a classifying domain.

77. A computer as recited in claim 67, wherein the request corresponds to at least a first and second subset of the objects, the request comprising a logical modifier that specifies an operation, and wherein the processing means further comprise means for:

responsive to receiving the request, identifying at least a portion of the first and second subsets of directory objects in the polyarchical data set; and

wherein transforming the one or more objects further comprises performing the operation on the first and second subsets.

78. A computer as recited in claim 77, wherein the logical modifier is a Boolean modifier.

79. A computer as recited in claim 77, wherein the operation comprises any combination of filtering, union, intersection, join, and/or exclusion operations.

80. A computer as recited in claim 67, wherein the means for accessing the one or more objects further comprises means for accessing the one or more objects in a manner that is independent of any inter-object relationship between the one or more objects and any other object of the objects in a manner that is independent of any definition of a hierarchy in the data polyarchy.

81. A computer as recited in claim 67, wherein the means for accessing the one or more objects further comprises querying the data polyarchy for the one or more objects.

82. A computer as recited in claim 67, wherein the means for accessing the one or more objects further comprises means for managing, manipulating, or modifying the one or more objects or a relationship between an object of the one or more objects and one of more different objects of the objects.

83. A polyarchical query language data structure comprising:

a first data field to specify a particular schema, the particular schema indicating how to meaningfully present or manage a plurality of objects in a data polyarchy based on values of attributes in the objects; and

a second data field to indicate an attribute of interest; and

a third data field to indicate how one or more objects comprising the attribute of interest are to be presented or managed with respect to one or more participating dimensions of inter-object relationships which are based on the schema.

84. A polyarchical query language data structure as recited in claim 83 further comprising a fourth data field to indicate a physical access strategy with respect to the data polyarchy, the physical access strategy being identified by indicating that the attribute of interest belongs to a distinguishing domain, a classifying domain, or a locating domain.

85. A polyarchical query language data structure as recited in claim 83, wherein the third data field further comprises a modifier to limit the one or more objects.

86. A polyarchical query language data structure as recited in claim 83, wherein the third data field further comprises a logical modifier to limit the one or more objects.

87. A polyarchical query language data structure as recited in claim 83, wherein the third data field further comprises a dimension information indicator for specifying a dimension within which to present the one or more objects.

88. A polyarchical query language data structure as recited in claim 83, wherein the third data field further comprises a dimension information indicator for specifying a particular direction and a particular depth within which to present a data relationship between a complex object of the one or more objects and one or more different objects of the one or more objects.

89. A polyarchical query language data structure as recited in claim 83, wherein each data field is expressed in an XML data format.

90. A polyarchical query language data structure as recited in claim 83, wherein the particular schema provides access to only a first subset of the objects to provide access control to the objects.

91. A computer-readable medium comprising a polyarchical query language data structure as recited in claim 83.

* * * * *

46/3,K/5 (Item 5 from file: 350) <u>Links</u>
Derwent WPIX

010635979   **Image available**
WPI Acc No: 1996-132932/**199614**
XRPX Acc No: N96-111846
  **Reference condition integrating device for database search
  operation - has database reference execution part which, on receiving
  database integrating condition specified by user, refers to
  database**
Patent Assignee: NEC CORP (NIDE)
Inventor: NIHEI K
Number of Countries: 002   Number of Patents: 002
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
JP 8016616	A	19960119	JP 94153198	A	19940705	199614	B
US 5819253	A	19981006	US 95498778	A	19950705	199847	

Priority Applications (No Type Date): JP 94153198 A 19940705
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
JP 8016616	A		6	G06F-017/30	
US 5819253	A			G06F-017/30	

  **has database reference execution part which, on receiving
  database integrating condition specified by user, refers to
  database**

  Abstract (Basic): condition input parts (1). Using these input parts,
    different database reference condition types such as **SQL,**
    menu, natural language type etc are input to the device. If
    needed, the **user** is able to **select** and **combine** the
    different database reference conditions. The **combined** condition
    is input by the **user** to the device using an integrated reference
    condition receptionist part (2). The integrated reference condition...

  ...Any **change** in the integrated reference conditions stored in the
    memory part is performed by the **user** at any time. The integrated
    reference condition generation part (4) then performs the actual
    combination or integration of the database reference conditions
    as per the **user's** specification stored in the memory part. This
    reference condition is then conveyed to an...

  ...of the execution part is displayed on the display part of the device for
    the **user.**

  Title Terms: **USER;**
International Patent Class (Main): **G06F-017/30**

International Patent Class (Additional): **G06F-012/00**
Manual Codes (EPI/S-X): **T01-J05B3**

[54] **DATABASE RETRIEVAL SYSTEM HAVING A PLURALITY OF INDEPENDENT INPUT MODES FOR DATABASE RETRIEVAL CONDITIONS**

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[73] Assignee: **NEC Corporation**, Tokyo, Japan

[21] Appl. No.: **498,778**

[22] Filed: **Jul. 5, 1995**

[30] **Foreign Application Priority Data**

Jul. 5, 1994 [JP] Japan 6-153198

[51] **Int. Cl.⁶** **G06F 17/30**

[52] **U.S. Cl.** **707/2; 707/1; 707/5; 707/7;**
707/531; 364/962; 364/962.3; 364/963;
364/963.1; 364/963.3; 364/974; 364/974.4;
364/DIG. 2

[58] **Field of Search** 395/600, 601;
707/5, 1, 7, 531; 364/962, 962.3, 963, 963.1,
963.3, 974, 974.4, DIG. 4, DIG. 2

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Primary Examiner—Thomas G. Black
*Assistant Examiner*—Jean H. Corriélus
Attorney, Agent, or Firm—Foley & Lardner

[57] **ABSTRACT**

A database retrieval system includes a retrieval condition input unit for inputting database retrieval conditions in a plurality of retrieval modes different from each other, an integration condition acceptance unit for inputting integration conditions for combining a plurality of database retrieval conditions input through the retrieval condition input unit, an integrated retrieval condition generation unit for integrating retrieval conditions input by the retrieval condition input unit according to integration conditions input by the integration condition acceptance unit, and a retrieval execution unit for retrieving a database according to retrieval conditions integrated by the integrated retrieval condition generation unit.

16 Claims, 4 Drawing Sheets
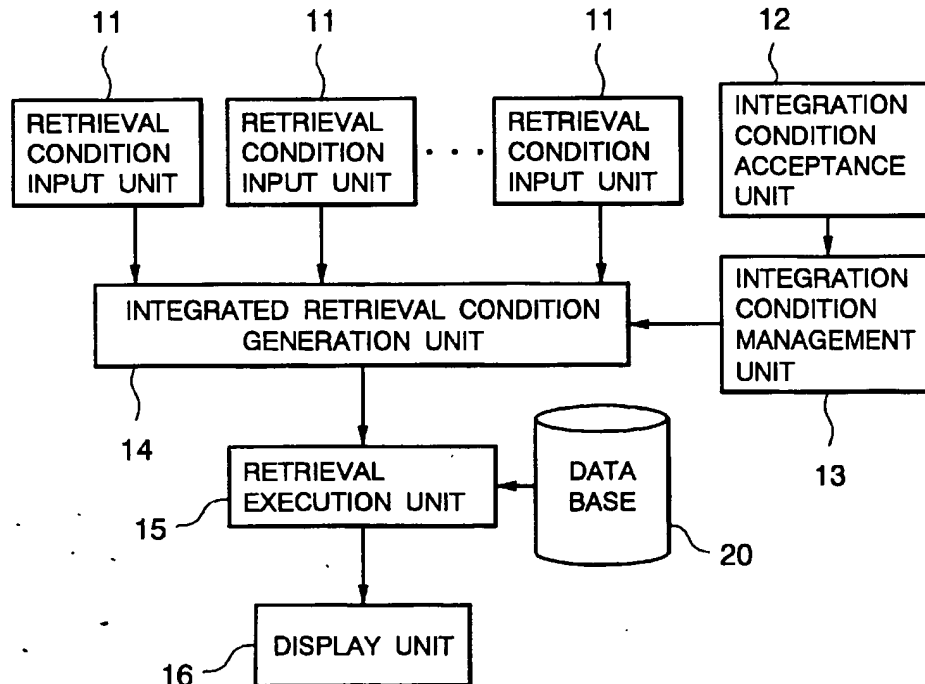
expression nor an integration condition corresponding to the retrieval mode. The foregoing processing will be repeated until information corresponding to all the retrieval mode selection designation units 22 is read (Step 405). Thereafter, the unit 14 combines the stored retrieval conditional expressions based on the stored integration conditions to generate and output an integrated retrieval conditional expression (Step 406).

For example, when integration conditions are set as shown in FIG. 2, an integrated retrieval conditional expression is generated that is a retrieval conditional expression of the retrieval mode a, a retrieval conditional expression of the retrieval mode b, and a retrieval conditional expression of the retrieval mode c.

The retrieval execution unit 15, is implemented by program-controlled CPU and memory, or the like, of a personal computer or a work station. The execution unit 15 retrieves the database 20 according to an integrated retrieval conditional expression generated by the integrated retrieval condition generation unit 14 to output retrieval results. In the present invention conventional manners are applicable, but no specific manner of database retrieval is proposed here. For example, the method recited in the reference literature by Shunsuke Uemura, *Basics of Database System*, Ohmsha, Ltd., can be used.

The display unit 16 is implemented by a display device or the like. The display unit 16 displays retrieval results obtained by the retrieval execution unit 15. The unit 16 is also capable of accommodating the interface panel 21 of the integration condition acceptance unit 12 as described above.

With reference to the flow chart shown in FIG. 5, the flow of data retrieval processing according to the present embodiment will be described.

Referring to FIGS. 1, 2, and 5, first, a user inputs a desired retrieval condition to the retrieval condition input unit 11 by using an input device such as a keyboard (Step 501). Then, the user inputs a desired integration condition to the integration condition acceptance unit 12 by means of the interface panel 21 or the like (Step 502). Upon input of these conditions, the integrated retrieval condition generation unit 14 receives input of a retrieval conditional expression from the retrieval condition input unit 11 and receives input of the integration condition from the integration condition management unit 13 to generate an integrated retrieval conditional expression (Step 503). Then, based on the generated integrated conditional expression, the retrieval execution unit 15 executes retrieval of the database 20 (Step 504) and the display unit 16 displays retrieval results obtained by the retrieval execution unit 15 (Step 505).

Referring to the retrieval results displayed on the display unit 16, the user is allowed to modify integration conditions when finding the retrieval results to be unsatisfactory (Step 506). If the retrieval results are unsatisfactory then it goes back to Step 502, and the user re-enters the integration conditions. When the integration conditions are re-entered, the integration condition management unit 13 detects modification of integration conditions so as to notify the integrated retrieval condition generation unit 14. The integrated retrieval condition generation unit 14 generates an integrated retrieval conditional expression based on the new integration conditions in response to the notification from the integration condition management unit 13 (Step 503). Then, the database 20 is retrieved (Step 504) and retrieval results are displayed (Step 505).

As described in the foregoing, since the database retrieval system of the present invention retrieves a database based on

combined retrieval conditions in a plurality of different retrieval modes, entry of necessary retrieval conditions and their combinations enables the system to cope with the retrieval under multiple retrieval conditions, whereby time and labor necessary for the processing can be reduced.

In addition, referring to displayed retrieval results, a user is allowed to obtain retrieval results one after another based upon different integration conditions. It is therefore possible for the system of the present invention to execute flexible retrieval with ease according to a user's request.

Although the invention has been illustrated and described with respect to exemplary embodiment thereof, it is understood by those skilled in the art that the foregoing and various other changes, omissions and additions may be made therein and thereto, without departing from the spirit and scope of the present invention. Therefore, the present invention is not limited to the specific embodiment set out above, but is comprised of all possible embodiments which are within the scope, and equivalents thereof, with respect to the features set out in the appended claims.

What is claimed is:

1. A database retrieval system comprising:

a plurality of independent retrieval condition input units, each input unit receiving an input of database retrieval conditions and operating in a unique one of a first arbitrary number of different retrieval modes, each input unit generating a retrieval conditional expression based on the input, the number of input units corresponding to the first arbitrary number of retrieval modes;

integration condition input unit for inputting integration conditions for combining the retrieval conditional expressions generated by said independent retrieval condition input units;

retrieval condition integrating unit for integrating a second arbitrary number of the retrieval conditional expression s generated by said independent retrieval condition input units according to the integration conditions, said retrieval condition integrating unit generating an integrated retrieval conditional expression; and

retrieval execution unit for executing retrieval of a database according to the integrated retrieval conditional expression.

2. The database retrieval system as set forth in claim 1, wherein the integration conditions include a condition for designating whether or not retrieval in a predetermined retrieval mode is to be executed and a logical operator indicative of a combination relationship among the retrieval conditional expressions associated with the retrieval modes designated to be executed.

3. The database retrieval system as set forth in claim 2, wherein said retrieval condition integrating unit

refers to the designated condition of the retrieval modes, receives the retrieval conditional expressions associated with the retrieval modes designated to be executed from said retrieval condition input units, and stores the retrieval conditional expressions;

receives and stores said logical operator; and

after receiving all retrieval conditional expressions associated with the retrieval modes designated to be executed, combines the stored retrieval conditional expressions using the logical operator to generate the integrated retrieval conditional expression.

4. The database retrieval system as set forth in claim 1, wherein the retrieval modes include at least one of SQL, a menu interface, and a natural language interface.

5. The database retrieval system as set forth in claim 1, further comprising:

    integration condition management unit connected to receive the integration conditions from said integration condition input unit, and to notify said retrieval condition integrating unit when the integration conditions have been modified from previous integration conditions received from said integration condition input unit.

6. A retrieval condition integrating device in a database retrieval system comprising:

    integration condition input unit for inputting integration conditions for combining a plurality of retrieval conditional expressions to generate an integrated retrieval conditional expression, each retrieval conditional expression being generated using a unique one of a first arbitrary number of different retrieval modes; and

    retrieval condition integrating unit for integrating a second arbitrary number of the retrieval conditional expressions according to the integration conditions.

7. The retrieval condition integrating device in a database retrieval system as set forth in claim 6, wherein the integration conditions include a condition for designating whether or not retrieval in a predetermined retrieval mode is to be executed and a logical operator indicative of a combination relationship among the retrieval conditional expressions associated with the retrieval modes designated to be executed.

8. The retrieval condition integrating device in a database retrieval system as set forth in claim 6, wherein said retrieval condition integrating unit

    refers to the designated condition of the retrieval modes, receives the retrieval conditional expressions associated with the retrieval modes designated to be executed, and stores the retrieval conditional expressions;

    receives and stores said logical operator; and

    after receiving all retrieval conditional expressions associated with the retrieval modes designated to be executed, combines the stored retrieval conditional expressions using the logical operator to generate the integrated retrieval conditional expression.

9. The retrieval condition integrating device in a database retrieval system as set forth in claim 6, wherein the retrieval modes include at least one of SQL, a menu interface, and a natural language interface.

10. The retrieval condition integrating device in a database retrieval system as set forth in claim 6, further comprising:

    integration condition management unit connected to receive the integration conditions from said integration condition input unit, and to notify said retrieval condition integrating unit when the integration conditions have been modified from previous integration conditions received from said integration condition input unit.

11. A database retrieval method comprising the steps of:

  (a) in putting database retrieval conditions and generating retrieval conditional expressions based on the database retrieval conditions, each retrieval conditional expression being generated using a unique one of a first arbitrary number of different retrieval modes;

  (b) inputting integration conditions for combining a second arbitrary number of the retrieval conditional expressions generated in said step (a);

  (c) integrating the retrieval conditional expressions generated in said step (a) according to the integration

conditions to generate an integrated retrieval conditional expression; and

  (d) retrieving a database according to the integrated retrieval conditional expression.

12. The database retrieval method as set forth in claim 11, wherein said step (b) of inputting integration conditions comprises the steps of:

    designating whether or not retrieval in a predetermined retrieval mode is to be executed; and

    setting a logical operator indicative of a combination relationship among the retrieval conditional expressions associated with the retrieval modes designated to be executed.

13. The database retrieval method as set forth in claim 12, wherein said step (c) of integrating retrieval conditions comprises the steps of:

    with reference to the designated condition of a retrieval mode receiving and storing the retrieval conditional expressions associated with the retrieval modes designated to be executed;

    receiving and storing said logical operator; and

    after receiving all retrieval conditional expressions associated with the retrieval modes designated to be executed, combining the stored retrieval conditional expressions using the logical operator to generate the integrated retrieval conditional expression.

14. The database retrieval method as set forth in claim 11, wherein the retrieval modes include at least one of SQL, a menu interface, and a natural language interface.

15. A database retrieval system comprising:

    at least four independent retrieval condition input units, each input unit receiving an input of database retrieval conditions and operating in a unique retrieval mode, each input unit generating a retrieval conditional expression based on the input;

    integration condition input unit for inputting integration conditions for combining the retrieval conditional expressions generated by said independent retrieval condition input units;

    retrieval condition integrating unit for integrating the retrieval conditional expressions generated by said independent retrieval condition input units according to the integration conditions, said retrieval condition integrating unit being capable of integrating more than three retrieval conditional expressions, and said retrieval condition integrating unit generating an integrated retrieval conditional expression; and

    retrieval execution unit for executing retrieval of a database according to the integrated retrieval conditional expression.

16. A method of generating an integrated retrieval conditional expression comprising the steps of:

  (a) providing an integration condition input unit having an arbitrary number of mode selection designation units corresponding to unique retrieval modes;

  (b) inputting retrieval mode information into said mode selection designation units;

  (c) inputting integration conditions corresponding to said retrieval mode information into said integration condition input unit;

  (d) reading the retrieval mode information from one of said mode selection designation units into a retrieval condition integrating unit;

  (e) determining whether the retrieval mode of said one of said mode selection designation units is to be used based on the retrieval mode information;

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(f) if the retrieval mode is to be used, storing in said retrieval condition integrating unit a retrieval conditional expression corresponding to the retrieval mode, then reading the integration condition corresponding to said retrieval mode information from said integration condition input unit, and determining whether all integration conditions entered into said integration condition input unit have been read;

(g) if the retrieval mode is not to be used, determining whether all integration conditions entered into said integration condition input unit have been read;

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(h) if all integration conditions entered into said integration condition input unit have been read, generating an integrated retrieval conditional expression based on the stored retrieval condition expressions and the stored integration conditions; and

(i) if all integration conditions entered into said integration condition input unit have not been read, repeating steps (d)–(g).

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46/3,K/14 (Item 14 from file: 350) **Links**
Derwent WPIX

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016947819    **Image available**
WPI Acc No: 2005-272127/200528
Related WPI Acc No: 2003-746922; 2005-272128
XRPX Acc No: N05-223543
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Computer readable medium storing schema generation program, has specific data field to indicate presentation or management of objects comprising attribute of interest, with respect to participating dimensions of inter-object relationships

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Patent Assignee: MICROSOFT CORP (MICT  )
Inventor: BOOTH J H; CAMERON K; CLEMENT L; MACLEOD S P; ROBERTSON G G
Number of Countries: 001  Number of Patents: 001
Patent Family: ·
Patent No    Kind· .Date    Applicat No    Kind    Date     Week
.US 20050071354 A1  20050331  US 2001995415  A    20011126  200528  B
                              US 2004967435  A    20041018
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Priority Applications (No Type Date): US 2001995415 A 20011126; US
   2004967435 A 20041018
Patent Details:
Patent No  Kind Lan Pg   Main IPC    Filing Notes
US 20050071354 A1     32 G06F-017/00   Cont of application US 2001995415
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Abstract (Basic):
...   A polyarchical **query language** data structure has a
   data **field** to **specify** schema indicating meaningful
   presentation or management of **objects** in data polyarchy based on
   attribute values. in **objects**. Other data **fields** indicate
   an attribute of interest, and indicate presentation or management of
   objects comprising the attribute of interest with respect to
   participating dimensions of inter-**object relationships**
   which are based on schema.
...·.   For storing program of dynamically generated schema representing
   several hierarchies of inter-**object relationships** for
   computing device (claimed) e.g. personal computer (PC), server
   computer, thin **client**, thick **client**, hand-held or laptop
   device, multiprocessor system, microprocessor-based system, **set**
   top box, programmable consumer electronics, network PC, minicomputer
   and mainframe computer, in distributed computing environment...

...The inter-**object relationships** are managed and navigated,
efficiently using data polyarchy schema that is generated to access
elements...
...Title Terms: **FIELD;**
International Patent Class (Main): **G06F-017/00**
Manual Codes (EPI/S-X): **T01-F07**...

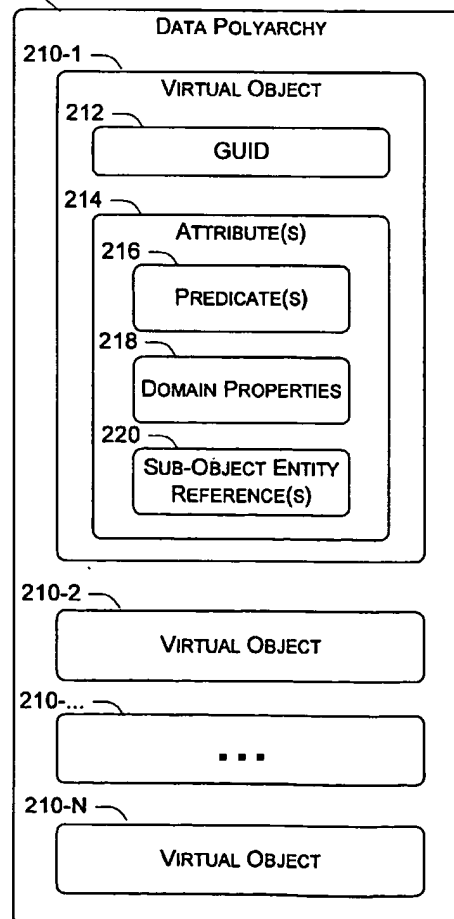...**T01-J05B2B**...

...**T01-J05B3**...

...**T01-S03**

(54) **DYNAMICALLY GENERATED SCHEMA REPRESENTING MULTIPLE HIERARCHIES OF INTER-OBJECT RELATIONSHIPS**

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(21) Appl. No.: 10/967,435

(22) Filed: Oct. 18, 2004

(57) **ABSTRACT**

Systems and methods for dynamically generating a schema
representing multiple hierarchies of inter-object relation-
ships are described. In one aspect, a polyarchical query
language data structure includes first, second, and third data
fields. The first data field is used to specify a particular
schema for presenting or managing a plurality of objects in
a data polyarchy based on values of attributes in the objects.
The second data field is to indicate an attribute of interest.
The third data field indicates how one or more objects that
include the attribute of interest are to be presented or
managed with respect to one or more participating dimen-
sions of inter-object relationships based on the schema.

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DATA POLYARCHY

210-1

VIRTUAL OBJECT

212

GUID

214

ATTRIBUTE(S)

216

PREDICATE(S)

218

DOMAIN PROPERTIES

220

SUB-OBJECT ENTITY
REFERENCE(S)

210-2

VIRTUAL OBJECT

210-...

. . .

210-N

VIRTUAL OBJECT

limited to the specific features or operations described. Rather, the specific features and steps are disclosed as preferred forms of implementing the claimed present invention.

1. A computer-readable medium comprising:

a polyarchical query language data structure comprising:

a first data field to specify a particular schema, the particular schema indicating how to meaningfully present or manage a plurality of objects in a data polyarchy based on values of attributes in the objects; and

a second data field to indicate an attribute of interest; and

a third data field to indicate how one or more objects comprising the attribute of interest are to be presented or managed with respect to one or more participating dimensions of inter-object relationships which are based on the schema.

2. A computer-readable medium as recited in claim 1, wherein the third data field further comprises a modifier to limit the one or more objects.

3. A computer-readable medium as recited in claim 1, wherein the third data field further comprises a logical modifier to limit the one or more objects.

4. A computer-readable medium as recited in claim 1, wherein the third data field further comprises a dimension information indicator for specifying a dimension within which to present the one or more objects.

5. A computer-readable medium as recited in claim 1, wherein the third data field further comprises a dimension information indicator for specifying a particular direction and a particular depth within which to present a data relationship between a complex object of the one or more objects and one or more different objects of the one or more objects.

6. A computer-readable medium as recited in claim 1, wherein each data field is expressed in an XML data format.

7. A computer-readable medium as recited in claim 1, wherein the particular schema provides access to only a first subset of the objects to provide access control to the objects.

8. A computer-readable medium as recited in claim 1, wherein the polyarchical query language data structure further comprises a fourth data field to indicate a physical access strategy with respect to the data polyarchy, the physical access strategy being identified by indicating that the attribute of interest belongs to a distinguishing domain, a classifying domain, or a locating domain.

9. A method comprising:

specifying a particular schema via a polyarchical query language data structure, the polyarchical query language data structure comprising:

a first data field to specify the particular schema, the particular schema indicating how to meaningfully present or manage a plurality of objects in a data polyarchy based on values of attributes in the objects; and

a second data field to indicate an attribute of interest; and

a third data field to indicate how one or more objects comprising the attribute of interest are to be presented or managed with respect to one or more

participating dimensions of inter-object relationships which are based on the schema.

10. A method as recited in claim 19, wherein the third data field further comprises a modifier to limit the one or more objects.

11. A method as recited in claim 19, wherein the third data field further comprises a logical modifier to limit the one or more objects.

12. A method as recited in claim 19, wherein the third data field further comprises a dimension information indicator for specifying a dimension within which to present the one or more objects.

13. A method as recited in claim 19, wherein the third data field further comprises a dimension information indicator for specifying a particular direction and a particular depth within which to present a data relationship between a complex object of the one or more objects and one or more different objects of the one or more objects.

14. A method as recited in claim 19, wherein each data field is expressed in an XML data format.

15. A method as recited in claim 19, wherein the particular schema provides access to only a first subset of the objects to provide access control to the objects.

16. A method as recited in claim 19, wherein the inter-object relationships comprise a flat relationship, a hierarchical relationship, and multiple intersecting hierarchies of relationships.

17. A method as recited in claim 19, where the objects comprise enterprise resource planning (ERP) objects, directory based objects, or database objects.

18. A method as recited in claim 19, wherein the polyarchical query language data structure further comprises a fourth data field to indicate a physical access strategy with respect to the data polyarchy, the physical access strategy being identified by indicating that the attribute of interest belongs to a distinguishing domain, a classifying domain, or a locating domain.

19. A computing device comprising:

a processor; and

a memory coupled to the processor, the memory comprising computer-program instructions executable by the processor for:

specifying a particular schema via a polyarchical query language data structure, the polyarchical query language data structure comprising:

a first data field to specify the particular schema, the particular schema indicating how to meaningfully present or manage a plurality of objects in a data polyarchy based on values of attributes in the objects; and

a second data field to indicate an attribute of interest; and

a third data field to indicate how one or more objects comprising the attribute of interest are to be presented or managed with respect to one or more participating dimensions of inter-object relationships which are based on the schema.

20. A computing device as recited in claim 19, wherein the third data field further comprises a dimension information indicator for specifying a dimension within which to present the one or more objects.

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Derwent WPIX

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015594631     **Image available**
WPI Acc No: 2003-656786/200362
XRPX Acc No: N03-523210
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  Multimedia objects updating method, involves
  accepting agent program to invoke structured query language
  procedure to manipulate multimedia object in particular column and
  in particular table
```
Patent Assignee: ORACLE INT CORP (ORAC-N)
Inventor: DIAMOND D L; JANOSIK J L; OXBURY S J; RUBINO M
Number of Countries: 001  Number of Patents: 001
Patent Family:
Patent No      Kind   Date      Applicat No    Kind   Date      Week
US 6591295     B1     20030708  US 99434243    A      19991105  200362  B
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Priority Applications (No Type Date): US 99434243 A 19991105
Patent Details:
Patent No  Kind Lan Pg   Main IPC     Filing Notes
US 6591295     B1    11  G06F-015/16
```

Abstract (Basic):
...   The method involves executing **graphical user**
    interface clipboard program (117) that accepts an agent program
    (111) **designation** to provide access to **relational**
    database (113). An alphanumeric uniform resource locator is transferred
    to **designate** a structured **query language** (
    SQL) procedure (112) to an application program. The **SQL**
    procedure is invoked using the locator to manipulate a multimedia
    object.
...   An INDEPENDENT CLAIM is also included for an apparatus for
    storing, **updating** and retrieving data into a web server database
    ...

...Used for storing, retrieving and manipulating multimedia data
    objects.
        ...

...The clipboard program provides a **graphical user**
    interface that easily identifies and retrieves a multimedia
    object stored in the database. The clipboard is used to capture
    a multimedia **object** from a device through Internet to store the

new **object** in a **selected** row and column location in a
selected database table, thereby facilitating the use of
objects in web applications...

...The drawing shows a block diagram explaining the interactions between
the principle components of multimedia **objects** storing and
updating apparatus...

...**SQL** procedure (112...

...**Relational** database (113
Title Terms: **OBJECT;**
International Patent Class (Main): **G06F-015/16**
Manual Codes (EPI/S-X): **T01-J10**

(12) **United States Patent**

Diamond et al.

(10) Patent No.: **US 6,591,295 B1**
(45) Date of Patent: **Jul. 8, 2003**

(54) **METHODS AND APPARATUS FOR USING MULTIMEDIA DATA STORED IN A RELATIONAL DATABASE IN WEB APPLICATIONS**

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(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 0 days.

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(22) Filed: . **Nov. 5, 1999**

(51) Int. Cl.[7] G06F 15/16
(52) U.S. Cl. 709/217; 709/219; 709/203; 709/226; 707/3; 707/10; 707/102
(58) Field of Search .... 709/217, 218, 709/219, 226, 203; 707/102, 100, 3, 10, 10 G

(56) **References Cited**

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* cited by examiner

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(74) Attorney, Agent, or Firm—Sanjay Prasad

(57) **ABSTRACT**

A Web agent software interface between a relational database and Web-based application programs employs stored SQL procedures to store, update and retrieve multimedia objects. Web based applications manipulate the data in the relational database table by using URLs consisting of an alphanumeric designation of the host computer and port, the Web agent program associated with the database, the database agent configuration information associated with the database, specific stored SQL procedure for performing the desired data manipulation, and additional parameter data needed to specify and manipulate a particular multimedia object. A clipboard application program provides a graphical user interface which a user may employ to select and view selected multi-media objects stored in selected relational database tables, and to transfer the Web URLs designating those selected objects to other Web applications, such as HTML authoring tools, using convenient "drag-and-drop" and "cut, copy and paste" operations. The Web agent program retrieves data from the database which are designated by such URLs, and accepts POST requests from HTML forms which are activated by the user of a Web browser to load or update data in the relational database.

10 Claims, 4 Drawing Sheets

http://domain:port/intermedia/agent_name/request_mode/proc_name/path_info
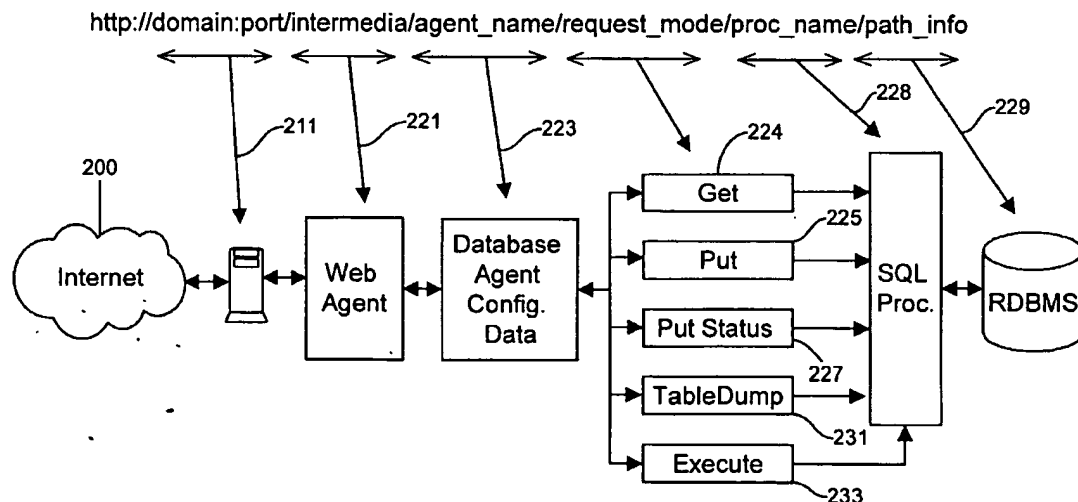
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of the object is displayed and this value is passed via a parameter in the SQL procedure. The object is loaded into the database for persistent storage when user clicks on "Update" in the Database tab.

An existing multimedia object may updated in much the same way. The particular table containing the object is selected using the navigation tree 311, and a row and column is selected in that table using the table view 315 so that the object to be updated is displayed in the object view. The table view may be used to display a subset of the records in any table sorted in a desired order using conventional database query and sort procedures selected via the menu at 310. Once a particular object has been identified in a particular row and column location in the table view, that object may then be copied and pasted into an editor, saved as a file and opened in an editor, or replaced by a newly captured image or by a different image fetched via the Internet as described above. A list of one or more editors which may be invoked is presented at the Edit tab 338. When the updated image is present in the object view, the Update button at the Database tab 337 is clicked to update the object.

To retrieve an object from the database and use that object in an application that supports "drag-and-drop" operations, the object is first displayed in the object view as described above. The user then specifies Table→Drag Mode→Embed from the menu at 310 to place a reference (URL) to the database object in the target application, such as an HTML editor for creating Web pages. The necessary <img src= "url"> tag including the database URL is then automatically inserted into the Web page HTML, and the image is seen in the Web page editor. Alternatively, by selecting "Edit→Drag Mode→Link" a link of the form "anchor text" containing the object's database URL is automatically placed in the Web page HTML. Having selected the drag mode, the icon which represents the object in the table view as seen at 341 may be dragged to and dropped into the target application, such as the Web page authoring tool which constructs the appropriate HTML tag.

Storing Data From Web Pages

The Web agent illustrated at 111 in FIG. 1 may also be used to store data supplied via a Post request method from an HTML form. When the HTML Web page containing the input form is written, the parameters needed are incorporated into the form markup as illustrated by the following example:

```
<form
action="http://nedc.us.ora.com:8007/intermedia/emp_write/mediaput/
PUT_EMP_PIC"
method="post" enctype="multipart/form-data">
<input type="hidden" name="ord_post_put_call" value=
"SET_EMP_PIC">
Employee ID: <input type="text" name="ord_procedure_path"
length=5>
Employee Picture File: <input type="file" name="ord_content">
<input type="Submit" value="Store Employee Picture Now">
</form>
```

FIG. 4 of the drawings illustrates an HTML web page that contains the form specified by the markup above. Note in the HTML above that the "action" specifies the full URL pathname for the SQL procedure PUT_EMP_PIC which has been previously stored for use by the Web agent in storing multimedia objects in a particular database table. In the URL, "emp_write" is the database agent name. The encoding type is specified as "multipart/form-data" and the

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multimedia object is identified with the variable name "ord_content" which specifies the value entered by the browser user for the particular image file that contains the multimedia object to be loaded. The key value of the table row which is to receive the object is specified by the "ord_procedure_path" variable name, and the "ord_post_put_call" variable name is used to set the properties of the data after it is stored in the database. The variable specifies the name ("SET_EMP_PIC") of the SQL procedure that performs post-processing.

Using the method illustrated above, multimedia objects may be inserted or updated into a Web server database from any Web browser. Similarly, as noted earlier, multimedia objects may be imbedded in, or linked from, a Web page display using conventional Web page authoring tools in combination with the automatic URL generation and transfer mechanism provided by the clipboard and Web agent.

Conclusion

It is to be understood that the methods and apparatus which have been described are merely illustrative applications of the principles of the invention. Numerous modifications may be made by those skilled in the art without departing from the true spirit and scope of the invention.

What is claimed is:

1. A method of storing and updating multimedia objects in a relational database and retrieving said objects via the Internet comprising, in combination, the steps of:

connecting a host computer to the Internet,

programming said host computer to function as a relational database capable of storing, updating and retrieving data objects logically stored in one or more database tables,

executing an application program which utilizes multimedia objects,

executing a graphical user interface clipboard program for performing the steps of:

accepting from a user the designation of an agent program which provides access to said relational database, the designation of a particular table in said database, and the designation of a particular column in said particular table,

accepting from said user and storing a SQL procedure at said host computer, said SQL procedure being expressed at least in part in a structured query language for manipulating a multimedia object in said particular column in said particular table, and

transferring an alphanumeric Uniform Resource Locator which designates said SQL procedure from said clipboard program to said further application program, and

executing said agent program for performing the steps of:

accepting a data manipulation request from said application program which includes said alphanumeric Internet Uniform Resource Locator, and

invoking said SQL procedure designated by said alphanumeric Uniform Resource Locator to manipulate said multimedia object.

2. The method set forth in claim 1 further comprising the steps of:

placing a hypertext link in a Web page, said link including said Uniform Resource Locator, and

processing said Web page in a browser application program to activate said hypertext link by transmitting said data manipulation request for processing by said agent program.

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3. The method of claim 1 wherein said step of transferring is performed using an operating system interprocess data transfer procedure.

4. The method of claim 1 wherein said step of transferring is performed by executing a drag-and-drop procedure in 5 which an iconic representation of said multimedia object is visually moved between a screen display area produced by said clipboard and a screen display area produced by said application program.

5. The method of claim 3 wherein transfer procedure is an 10 operating system copy and paste procedure.

6. The method of claim 3 wherein said transfer procedure is an operating system cut and paste procedure.

7. Apparatus for storing, updating and retrieving data comprising, in combination, 15

a host computer including a memory for storing programs and data,

a relational database program stored in and executable by said host computer for storing, updating and retrieving data in one or more database tables, said relational 20 database program including means for interpreting and executing stored data manipulation procedures expressed at least in part in a structured query language,

means for storing at least one of said procedures for 25 manipulating specific data stored in one or more of said database tables, said one of said procedures being designated by a procedure name and including means for processing externally supplied parameter information to locate said specific data in said database tables, 30

a Web agent program stored in and executable by said host computer,

a HTTP interface program for receiving data requests via the Internet from a remotely located computer directed to a resource specified by a multi-part alphanumeric 35 Universal Resource Locator (URL) which consists of:

(a) an identification of the domain name and port of said host computer,

(b) the identification of said Web agent program which operates as an interface to a relational database,

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(c) the specification of said one of said procedures, and

(d) said parameter information,

said Web agent program including means for invoking the operation of said database program to perform said one of said procedures to perform a predetermined data manipulation operation with respect to said specific data, and

a graphical interface application program stored in and executable by said host computer for transferring data from said relational database to a second application program executing on said host computer, said clipboard application program comprising:

means for displaying representations of the content of said database tables,

means for accepting from a user an identification of at least a selected one of said displayed representations to identify said specific data stored in one or more of said database tables, and

means for transferring to said second application program a Uniform Resource Locator which specifies one of said procedures and parameter information and may be transmitted to said interface program to invoke said Web agent program for retrieving said specific data from said database tables using said one of said procedures.

8. The apparatus of claim 7 wherein said means for transferring includes means for performing a drag-and-drop procedure in which an iconic representation of said specific data is visually moved between a screen display area produced by said clipboard application program and a screen display area produced by said second application program.

9. The apparatus of claim 8 wherein said means for transferring further includes means for performing a copy and paste procedure.

10. The apparatus of claim 9 wherein said means for transferring further includes means for performing a cut and paste procedure.

* * * * *

```
015441274      **Image available**
WPI Acc No: 2003-503416/200347
XRAM Acc No: C03-134486
XRPX Acc No: N03-399615
```

Biological data accessing system used for searching database of biological information comprises database graph generation module associated with search engine, and joins module to create joins between database tables based on graph

```
Patent Assignee: ROTH C (ROTH-I)
Inventor: ROTH C
Number of Countries: 001  Number of Patents: 001
Patent Family:
```

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 20030041053	A1	20030227	US 2001938714	A	20010823	200347	B

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Priority Applications (No Type Date): US 2001938714 A 20010823
Patent Details:
Patent No   Kind Lan Pg   Main IPC      Filing Notes
US 20030041053 A1   · 17 G06F-007/00
```

... searching database of biological information comprises database graph generation module associated with search engine, and joins module to create joins between database tables based on graph

Abstract (Basic):
... database graph generation module associated with a search engine to generate a database graph and **joins** module to create **joins** between database tables based on database graph. The server computer runs a structured **query language** search on the database based upon the **joins**.
... graph generation module associated with a search engine to generate a database graph and a **joins** module to create **joins** between database tables based on the database graph. The server computer runs a structured **query language** (**SQL**) search on the database based upon the **joins**.
...

...An INDEPENDENT CLAIM is also included for querying a **relational** database which comprises sending a structured language database query to a search engine, parsing the database and creating a database graph, creating the correct **joins corresponding** to the query, translating the structured query into an **SQL statement** incorporating the **joins** and sending the **SQL statement** to a **relational** database...

...of time to rewriting the code for existing databases or components and
to minimize the **changes** to existing databases needed to
update the system with new functionalities. The system preserves
the property of component reuse, and is...

...a new utility or application is added to the system, or if a component
is **changed,** little or no **lines** of code have to be
changed in other applications or services improving upgrade
performance and reducing maintenance time
Technology Focus:
Preferred Components: The system comprises a second
module that receives the results of the **SQL** search and
translates the search results into a structured language. The
SQL is sent to a **client** computer. The first module
comprises a **user interface** that provides a list of
searchable **fields** within the database and comprises a viewer
module to present search results in a graphical
...Title Terms: **JOIN;**
International Patent Class (Main): **G06F-007/00**
Manual Codes (EPI/S-X): **T01-E**

(19) **United States**

(12) **Patent Application Publication** (10) Pub. No.: US 2003/0041053 A1

Roth (43) Pub. Date: Feb. 27, 2003

(54) **SYSTEM AND METHOD FOR ACCESSING BIOLOGICAL DATA**

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(52) U.S. Cl. 707/3

(57) **ABSTRACT**

A system is presented with a search panel for specifying search criteria for searching a database of biological information. The search panel uses the extensible markup language (XML) to send search requests to the database. A database graph generation module linked to the biological database generates a database graph representing the database schema. Once the database schema is known, another module creates joins between the database tables in order to most effectively join data from one table to another. An SQL statement incorporating the optimized joins is then used to search the biological database.
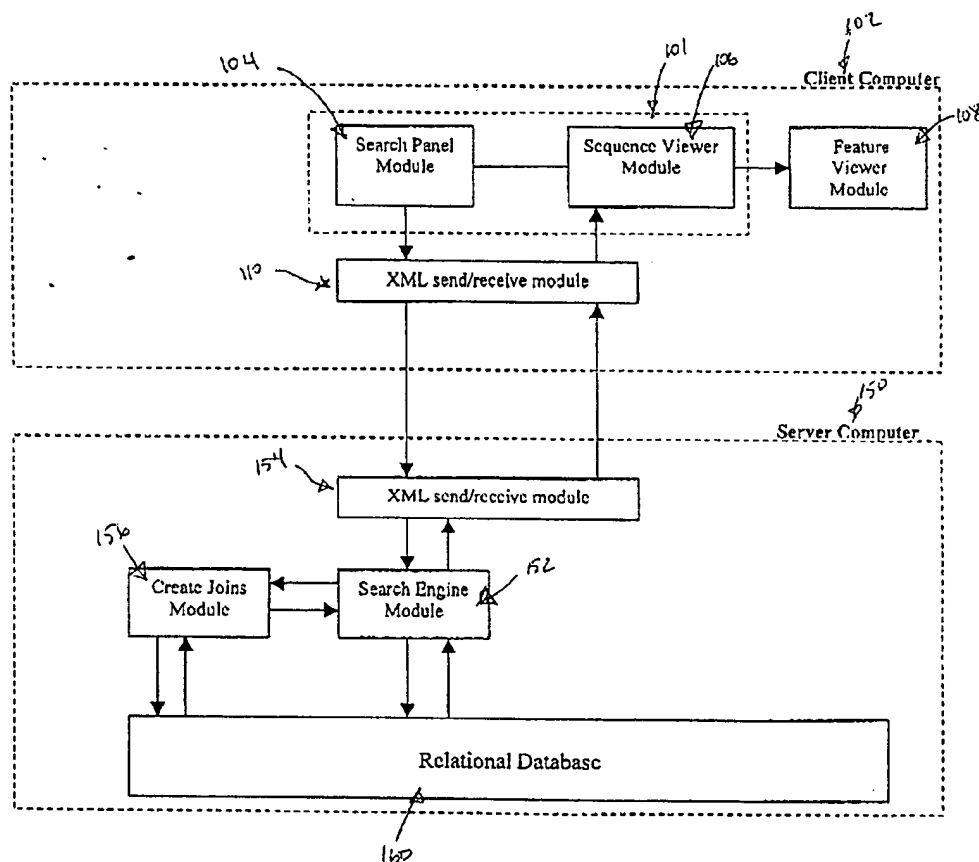
the Sequence Viewer Module 106 for representation to the user. The user can then view and manipulate the results, including invoking a Feature Viewer Module 108 to graphically represent the returned results.

[0090] This specific embodiment describes a method for searching a remotely located database. Searching a remote database in this manner is much slower than locating the Search Engine 520 on the Server Computer 150 along with the Relational Database 160. In this embodiment, the Database Graph Generation Module 502 must communicate and receive a substantial amount of information with the Relational Database 160 through a network system. The time required for this exchange is controlled by the available bandwidth, which will be much slower than if the Search Engine 520 were able to communicate with a Relational Database 160 stored on the same computer. Hence, locating the Search Engine 520 on the Server Computer 150 allows the Database Graph Generation Module 502 to perform its instructions much faster than if it had to communicate across a network protocol such as a TCP protocol.

[0091] In reference to **FIG. 7,** a process flow diagram is depicted. A User 702 accessing a Client Computer utilizes the Search Panel Database Selector 202 field presented within the Search Panel User Interface 704 to choose a database stored on a Server Computer in which to search. Alternatively, a User 702 may specify a Server Computer or a folder stored on a Server Computer in which to conduct the search without specifying a specific database. This way, all databases contained within the specified folder or stored on the specified computer can be searched according to the search criteria, thus allowing a user to retrieve specific results without designating a particular database.

[0092] Once a search location is selected, the user specifies filtering criteria in the Search Panel User Interface 704 presented by the Search Panel Module. An XML File 706 containing the search criteria is created and sent to the Search Tool Module 506 of the Search Engine, which resides on a server computer. The Search Tool Module 506 instructs the Database Graph Generation Module 502 to create a database graph. This is done by parsing the database and accessing the Entry Definitions 710. The Entry Definitions 710 have associated Data Entries 712, which include at least one Primary Key 714 and may contain one or more Secondary Keys 716. The database graph is generated, representing the database schema, from which the Join Path Generation Module 504 is able to construct the correct joins between the database tables and then subsequently calculate the most efficient path between the requested nodes. The efficient path is communicated to the Search Tool Module 506, which then sends the appropriate query, including the requested nodes and the calculated joins between the nodes, to the SQL Statement Generation Module 508. The SQL Statement Generation Module 508 translates the query into an appropriate SQL Statement 720 and sends the query to the Relational Database 160, which is governed by database management software, and returns the requested information in a corresponding SQL Statement 722 to the Search Tool Module 506.

[0093] The Search Tool Module 506 sends the SQL results to an XML send/receive module 154 which translates the SQL results into an appropriate XML file 726. The XML send/receive module 154 then sends the XML Results File

726 to the Search Panel User Interface 704. The graphical results are displayed to the User 702 in the Sequence Viewer Module 106. The User 702 is then able to invoke an external Viewing Application 728 to further view and modify the returned results.

[0094] The foregoing description specifically describes one embodiment and method of practicing the current invention. The invention, however, is not limited to the embodiment described herein. It should be understood that changes may be made to the specific modules or information flow without departing from the scope of the invention, thus, the scope of the present invention is defined solely by the following claims.

What is claimed is:

1. A system for searching a database of biological information, said system comprising:

   a server computer comprising a database of biological information and a first module for receiving a structured language query and transferring said query to a search engine;

   a database graph generation module associated with said search engine configured to generating a database graph; and

   a joins module configured to create joins between database tables based on said database graph, wherein said server computer runs a structured query language (SQL) search on said database based upon said joins.

2. The system of claim 1, comprising a second module that receives the results of said SQL search and translates said search results into a structured language.

3. The system of claim 2, wherein said structured query language is sent to a client computer.

4. The system of claim 1, wherein said first module comprises a user interface that provides a list of searchable fields within said database.

5. The system of claim 1, wherein said first module comprises a viewer module configured to present search results in a graphical format.

6. The system of claim 1, wherein said structured language comprises the extensible markup language (XML), JavaScript, or the hypertext markup language (HTML).

7. A computer system for searching a database of biological information, comprising:

   a database of biological information comprising tables of biological data;

   a search module configured to receive a structured language query and convert said structured language query into a search statement for querying said database of biological information; and

   a joins module configured to determine how to join said tables of biological data in order to provide the results of said query.

8. The search engine of claim 7, further comprising in XML send/receive Module for sending and receiving information to and from a Search Panel Module stored on a Client Computer.

9. The search engine of claim 8, wherein said XML send/receive Module receives an XML structured query from a Client Computer, and delivers said XML structured query to a search tool module.

10. The search engine of claim 7, wherein said Database Graph Generation Module creates a graph of a user-selected database.

11. The search engine of claim 10, wherein said Create Joins Module utilizes said database graph to create joins between database tables.

12. The search engine of claim 11, wherein said Create Joins Module calculates the shortest path between two database nodes thereby optimizing the retrieval of requested database data.

13. The search engine of claim 7, further comprising a SQL statement generation module for translating said XML structured query into an SQL statement and sending said SQL statement to said Relational Database.

14. A method for querying a relational database, comprising the steps of:

sending a structured language database query to a search engine;

parsing the database and creating a database graph;

creating the correct joins corresponding to said query;

translating said structured database query into an SQL statement incorporating said joins; and

sending said SQL statement to a Relational Database.

15. The method of claim 14, including the further step of optimizing said joins by calculating the shortest path between the nodes specified in said query.

16. The method of claim 15, including the further step of receiving requested results from said database search, translating said results into said structured data language and returning said results.

17. The method of claim 16, including the further step of displaying said search results.

18. The method of claim 18, wherein said structured language is the Extensible Markup Language.

* * * * *

46/3,K/16 (Item 16 from file: 350) <u>Links</u>
Derwent WPIX

016798651   **Image available**
WPI Acc No: 2005-122930/200513
XRPX Acc No: N05-106056
  **Structured query language script
  generation method for database management system, involves
  selecting associated template strings based on each
  instruction selected corresponding to type of user
  defined objects in data model**
Patent Assignee: DEFFLER T A (DEFF-I); COMPUTER ASSOC THINK INC (COMP-N)
Inventor: DEFFLER T A
Number of Countries: 108 Number of Patents: 002
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
WO 200508534	A2	20050127	WO 2004US21990	A	20040709	200513	B
US 20050120014	A1	20050602	US 2003486773	P	20030711	200537	
			US 2004888146	A	20040709		

Priority Applications (No Type Date): US 2003486773 P 20030711; US
  2004888146 A 20040709
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
WO 200508534	A2	E	24	G06F-017/30	

    Designated States (National): AE AG AL AM AT AU AZ BA BB BG BR BW BY BZ
    CA CH CN CO CR CU CZ DE DK DM DZ EC EE EG ES FI GB GD GE GH GM HR HU ID
    IL IN IS JP KE KG KP KR KZ LC LK LR LS LT LU LV MA MD MG MK MN MW MX MZ
    NA NI NO NZ OM PG PH PL PT RO RU SC SD SE SG SK SL SY TJ TM TN TR TT TZ
    UA UG US UZ VC VN YU ZA ZM ZW
    Designated States (Regional): AT BE BG BW CH CY CZ DE DK EA EE ES FI FR
    GB GH GM GR HU IE IT KE LS LU MC MW MZ NA NL OA PL PT RO SD SE SI SK SL
    SZ TR TZ UG ZM ZW
US 20050120014 A1   G06F-017/30   Provisional application US 2003486773

  **Structured query language script
  generation method for database management system, involves
  selecting associated template strings based on each
  instruction selected corresponding to type of user
  defined objects in data model**

Abstract (Basic):
...   The associated template **strings** (150) are **selected**
    based on each instruction **selected corresponding** to type
    of **user** defined objects (142) **selected** in data
    model (140): A portion of structured **query
    language (SQL) script** (160) is generated
    automatically based on **selected** template **strings**.
...   1) structured **query language script**

generation program; and...

...2) structured **query language script** generation
system...

...For generating structured **query language (SQL)**
script in database management system (DBMS) in computer...

...Enables **customizing** the **text** generated in any suitable
area of **SQL script,** and reduces debugging time and
administration time of data...

...The figure shows a schematic diagram of the **SQL script**
generation system...

...data **model** (140...

...**user** defined **objects** (142...

...template **strings** (150...

...**SQL script** (160
...Title Terms: **SCRIPT;**
International Patent Class (Main): **G06F-017/30**
Manual Codes (EPI/S-X): **T01-J05B3**...

...**T01-S03**

(19) **United States**

(12) **Patent Application Publication** (10) Pub. No.: US 2005/0120014 A1
   Deffler (43) **Pub. Date: Jun. 2, 2005**

(54) SYSTEM AND METHOD FOR GENERATING
   SQL USING TEMPLATES

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(21) Appl. No.: 10/888,146

(22) Filed: Jul. 9, 2004

   Related U.S. Application Data

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Publication Classification

(51) Int. Cl.[7] G06F 17/30
(52) U.S. Cl. 707/4

(57) **ABSTRACT**

A method for generating a structured query language (SQL)
script based on a template includes selecting one object from
a plurality of objects in a data model. At least one instruction
is selected based, at least in part, on a type of the selected
object and, then, an associated template string is selected
based on each selected instruction. At least a portion of a
SQL script is automatically generated based on the one or
more selected template strings.
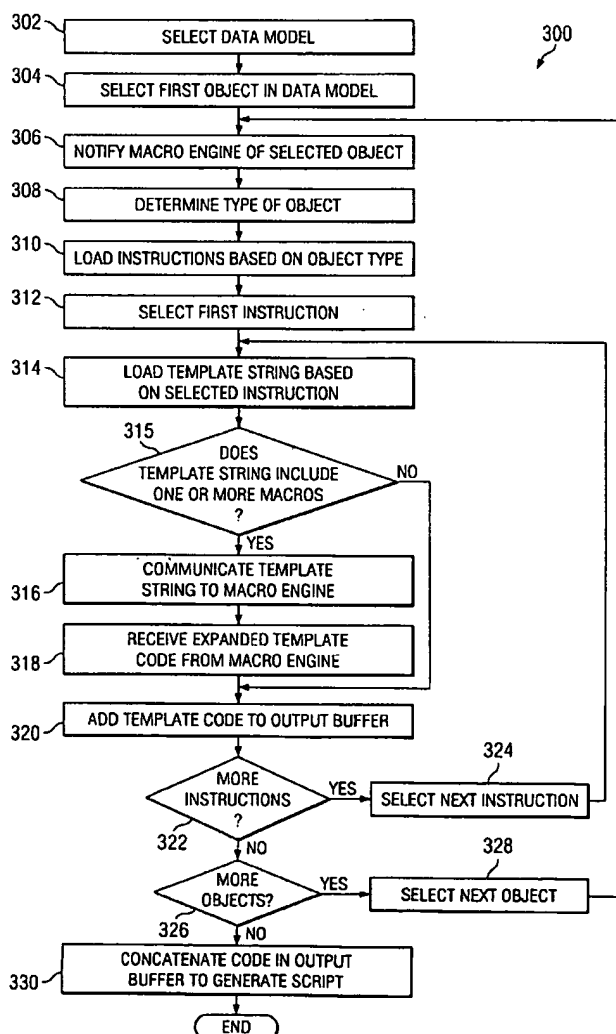
concatenate the portions of template code based, at least in part, on upon an order type (such as bucket identifier) stored in the template string, instruction, or object 142 as appropriate.

[0028] The preceding flowchart and accompanying description illustrate only an exemplary method 300 used by script generator 130 to generate SQL scripts 160 based on templates 150. However, system 100 contemplates script generator 130 using any suitable technique for performing these tasks. Thus, many of the steps in this flowchart may take place simultaneously and/or in different orders than as shown. Moreover, script generator 130 may implement methods with additional steps, fewer steps, and/or different steps, so long as the methods remain appropriate.

[0029] Although this disclosure has been described in terms of certain embodiments and generally associated methods, alterations and permutations of these embodiments and methods will be apparent to those skilled in the art. Accordingly, the above description of example embodiments does not define or constrain this disclosure. Other changes, substitutions, and alterations are also possible without departing from the spirit and scope of this disclosure.

What is claimed is:

1. A method for generating a structured query language (SQL) script based on a template comprises:

selecting one object from a plurality of objects in a data model;

selecting at least one instruction based, at least in part, on a type of the selected object;

selecting an associated template string based on each selected instruction; and

automatically generating at least a portion of a SQL script based on the one or more selected template strings.

2. The method of claim 1, one or more of the plurality of objects comprising a user-defined object.

3. The method of claim 1 further comprising selecting the data model from a plurality of data models.

4. The method of claim 1 further comprising:

determining the existence of one or more macros in one of the selected template strings; and

in response, at least in part, to locating one or more macros in the template string, processing the located one or more macros.

5. The method of claim 4 further comprising expanding the template string associated with the one or more macros based on a result from each macro processing and wherein generating the SQL script is further based on the expanded template string.

6. The method of claim 4, each macro selected from the group comprising:

determining a value of an environment variable;

determining a value of a property associated with the data model;

determining a value of a property associated with the selected object;

evaluating a conditional expression; and

iterating over a subset of the plurality of objects, the subset associated with the selected object.

7. The method of claim 1 further comprising:

selecting a second object from the plurality of objects in the data model;

selecting at least one instruction based, at least in part, on a type of the second object;

selecting an associated template string based on each selected instruction;

concatenating the template strings from the selected objects; and

wherein generating the SQL script based on the one or more selected template strings comprises generating the SQL script based on the concatenated template strings.

8. The method of claim 7, each instruction comprising a template string identifier and a bucket identifier and the method further comprises sorting the concatenated template strings based on the bucket identifier of each associated instruction.

9. The method of claim 1 further comprising dynamically adding user comments to the SQL script based on at least one of the selected template strings.

10. Software for generating a structured query language (SQL) script based on a template operable to:

select one object from a plurality of objects in a data model;

select at least one instruction based, at least in part, on a type of the selected object;

select an associated template string based on each selected instruction; and

automatically generate at least a portion of a SQL script based on the one or more selected template strings.

11. The software of claim 10, one or more of the plurality of objects comprising a user-defined object.

12. The software of claim 10 further operable to select the data model from a plurality of data models.

13. The software of claim 10 further operable to:

determine the existence of one or more macros in one of the selected template strings; and

in response, at least in part, to locating one or more macros in the template string, process the located one or more macros.

14. The software of claim 13 further operable to expand the template string associated with the one or more macros based on a result from each macro processing and wherein generating the SQL script is further based on the expanded template string.

15. The software of claim 13, each macro selected from the group comprising:

determining a value of an environment variable;

determining a value of a property associated with the data model;

determining a value of a property associated with the selected object;

evaluating a conditional expression; and

iterating over a subset of the plurality of objects, the subset associated with the selected object.

16. The software of claim 10 further operable to:

select a second object from the plurality of objects in the data model;

select at least one instruction based, at least in part, on a type of the second object;

select an associated template string based on each selected instruction;

concatenate the template strings from the selected objects; and

wherein the software operable to generate the SQL script based on the one or more selected template strings comprises software operable to generate the SQL script based on the concatenated template strings.

17. The software of claim 16, each instruction comprising a template string identifier and a bucket identifier and the software further operable to sort the concatenated template strings based on the bucket identifier of each associated instruction.

18. The software of claim 10 further operable to dynamically add user comments to the SQL script based on at least one of the selected template strings.

19. A system for generating a structured query language (SQL) script based on a template comprises:

memory operable to store a data model, a plurality of instructions, and a plurality of template strings and the data model comprising a plurality of objects; and

one or more processors operable to:

select one object from the plurality of objects in the data model;

select at least one instruction based, at least in part, on a type of the selected object;

select an associated template string based on each selected instruction; and

automatically generate at least a portion of a SQL script based on the one or more selected template strings.

20. The system of claim 19, one or more of the plurality of objects comprising a user-defined object.

21. The system of claim 19, the one or more processors further operable to select the data model from a plurality of data models.

22. The system of claim 19, the one or more processors further operable to:

determine the existence of one or more macros in one of the selected template strings; and

in response, at least in part, to locating one or more macros in the template string, process the located one or more macros.

23. The system of claim 22, the one or more processors further operable to expand the template string associated

with the one or more macros based on a result from each macro processing and wherein generating the SQL script is further based on the expanded template strings.

24. The system of claim 22, each macro selected from the group comprising:

determining a value of an environment variable;

determining a value of a property associated with the data model;

determining a value of a property associated with the selected object;

evaluating a conditional expression; and

iterating over a subset of the plurality of objects, the subset associated with the selected object.

25. The system of claim 19, the one or more processors further operable to:

select a second object from the plurality of objects in the data model;

select at least one instruction based, at least in part, on a type of the second object;

select an associated template string based on each selected instruction;

concatenate the template strings from the selected objects; and

wherein the processors further operable to generate the SQL script based on the one or more selected template strings comprises the processors further operable to generate the SQL script based on the concatenated template strings.

26. The system of claim 25, each instruction comprising a template string identifier and a bucket identifier and the one or more processors further operable to sort the concatenated template strings based on the bucket identifier of each associated instruction.

27. The system of claim 19, the one or more processors further operable to dynamically add user comments to the SQL script based on at least one of the selected template strings.

28. A system for generating a structured query language (SQL) script based on a template comprises:

means for selecting one object from a plurality of objects in a data model;

means for selecting at least one instruction based, at least in part, on a type of the selected object;

means for selecting an associated template string based on each selected instruction; and

means for automatically generating at least a portion of a SQL script based on the one or more selected template strings.

* * * * *

46/3,K/3 (Item 3 from file: 350) **Links**
Derwent WPIX

```
011030956      **Image available**
WPI Acc No: 1997-008880/199701
XRPX Acc No: N97-008090
  Database structure conversion - involves structured-
  query-language sentence conversion from
  physical structure into logic structure figure on table definition
  information database
Patent Assignee: NEC CORP (NIDE  )
Number of Countries: 001  Number of Patents: 001
Patent Family:
Patent No      Kind  Date     Applicat No    Kind   Date      Week
JP 8278979     A   19961022  JP 9581518      A    19950406  199701  B

Priority Applications (No Type Date): JP 9581518 A 19950406
Patent Details:
Patent No  Kind Lan Pg   Main IPC     Filing Notes
JP 8278979     A       4 G06F-017/30
```

```
  Database structure conversion - ...
...involves structured-query-language sentence
  conversion from physical structure into logic structure
  figure on table definition information database
```

...Abstract (Basic): The method involves automatic bi-directional
  updating of computer system logic structure in database table
  definition information. The computer system logic structure is derived
  from a **physical** structure...

...The **physical** structure is retrieved from an external memory (23)
  by a structured-**query-language** decipher unit (11)
  constituted by an **SQL sentence** which is decomposed into
  several **physical sentence** components. The **sentence**
  components are converted into **logical** structures by a
  structure figure generator (12...
...Title Terms: **CONVERT;**

JAPIO
05323479 **Image available**
METHOD AND DEVICE FOR DATA BASE STRUCTURE CONVERSION

METHOD AND DEVICE FOR DATA BASE STRUCTURE CONVERSION ...
Published: 19961022)

ABSTRACT

**PURPOSE: To eliminate contradiction between the logical structure and physical structure of a data base by
synchronizing both with each other... ...CONSTITUTION: The data base of a computer system is converted
between the logical structure based upon grammar for the logical structure and the physical structure based
upon grammar for the physical structure. Physical structure statements corresponding to the data base in the
physical structure are stored in an external storage device 23. An SQL statement decoding means 11 reads a
physical structure statement out of the external storage device 23 and decomposes the physical structure
statement into plural physical structure statement elements by using the grammar for the physical structure.
A structure diagram generating mean 12 converts the physical structure statement elements into logical
structure statement elements. The logical structure statement elements are processed according to the
grammar for the logical structure.**

METHOD AND DEVICE FOR DATA BASE STRUCTURE CONVERSION

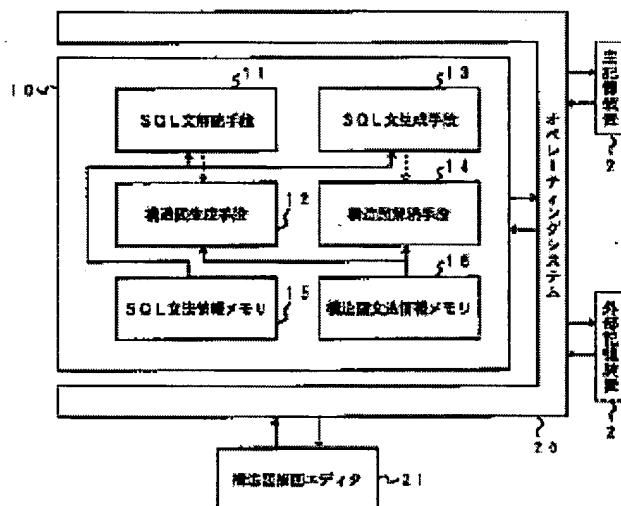|	
Patent number:	JP8278979
Publication date:	1996-10-22
Inventor:	HONDA MASAO
Applicant:	NIPPON ELECTRIC CO
Classification:	
- international:	*G06F12/00; G06F17/30;* G06F12/00; G06F17/30; (IPC1-7): G06F17/30; G06F12/00
- european:	
Application number:	JP19950081518 19950406
Priority number(s):	JP19950081518 19950406

Report a data error here

Abstract of **JP8278979**

PURPOSE: To eliminate contradiction between the logical structure and physical structure of a data base by synchronizing both with each other. CONSTITUTION: The data base of a computer system is converted between the logical structure based upon grammar for the logical structure and the physical structure based upon grammar for the physical structure. Physical structure statements corresponding to the data base in the physical structure are stored in an external storage device 23. An SQL statement decoding means 11 reads a physical structure statement out of the external storage device 23 and decomposes the physical structure statement into plural physical structure statement elements by using the grammar for the physical structure. A structure diagram generating mean 12 converts the physical structure statement elements into logical structure statement elements. The logical structure statement elements are processed according to the grammar for the logical structure.

46/3,K/125 (Item 125 from file: 350) **Links**

THE APPLICANT

009460710   **Image available**
WPI Acc No: 1993-154237/199319
XRPX Acc No: N93-117982
   **Graphical data base query system - includes selector
   for set of displayed parameters in order to define required query**

Patent Assignee: INT BUSINESS MACHINES CORP (IBMC)
Inventor: BANNING K R; JAMES W S; LI S
Number of Countries: 004   Number of Patents: 003
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
EP 541298	A2	19930512	EP 92309943	A	19921029	199319	B
EP 541298	A3	19931215	EP 92309943	A	19921029	199514	
US 5421008	A	19950530	US 91789507	A	19911108	199527	

Priority Applications (No Type Date): US 91789507 A 19911108
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
EP 541298	A2	E	42	G06F-015/403	

   Designated States (Regional): DE FR GB

US 5421008	A		40	G06F-015/40	
EP 541298	A3			G06F-015/403	

...   **includes selector for set of displayed
   parameters in order to define required query**

...Abstract (Basic): The apparatus for directing queries to a database
   containing information in order to access **selected** information
   comprises a display (38) for displaying in graphical form a number of
   parameters relating to queries (51 to 57). A **selection** unit (26)
   selects a **set** of the displayed parameters in order to
   define a required query and an operating unit is responsive to the
   selection of the parameters for performing the required query...

...A number of parameters relating to the query are displayed graphically.
   A **set** of parameters are **selected** in order to define the
   query. The method includes defining **relationships** between the
   selected parameters and defining **logical**
   relationships between them also...

...USE - Graphical database for structured **query language** (
   SQL) database systems...

...Abstract (Equivalent): extracting data from a first database in a
   computer system having a display device and **user** interaction
   device, involves displaying a number of query **object** class
   indicators, creating a number of query **objects** by repeatedly

selecting a query **object** class and entering query
object data for each of the query **objects**, and
linking the number of query **objects** by using the
user interaction device to repeatedly **select** a first
query **object**, a second query **object** and then
specify link parameters, thus creating a **link**
object. The query **objects** and the **link**
objects are stored in a second database...

....The method further involves accepting a query specified in **text,**
parsing the **text** query into a number of query **objects** and
link objects, storing the query **objects** and
link objects in the second database·and displaying the
query **objects** and the **link objects** on the display
device... ·

....ADVANTAGE - Provides **user** with graphical, query software to
directly manipuLate database...
...Title Terms:·**SELECT;**
International Patent Class (Main): **G06F-015/40...**

...**G06F-015/403**
Manual Codes (EPI/S-X): **T01-J05B4...**

...**T01-J12B**

THE APPLICANT (handwritten)

012228463   `**Image available**`.
WPI Acc No: 1999-034570/199903
XRPX Acc No: N99-025913
 **Computerized method for visually depicting join
 relationships in database management system - involves accepting
 user command which enables to choose among join
 relationships visually depicted on display unit**
Patent Assignee: INT BUSINESS MACHINES CORP (IBMC   )
Inventor: LEVINE J M; LI K; MOCEK D J
Number of Countries: 001 Number of Patents: 001
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 5842209	A	19981124	US 96707140	A	19960903	199903	B

Priority Applications (No Type Date): US 96707140 A 19960903
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 5842209	A		22	G06F-017/30	

 **Computerized method for visually depicting join
 relationships in database management system...**

 **...involves accepting user command which enables to choose
 among join relationships visually depicted on display unit**

...Abstract (Basic): The method involves visually depicting two tables,
     fields in the tables, and one or more **join
     relationships** between **fields** in the tables on a display
     unit coupled to a computer. The **join relationships**
     illustrate **join** operations performed between two or more tables
     in the database management system...

...A **user** command is accepted from a **user** input unit into the
     computer. The **user** command enables to **selectively
     choose** among the **join relationships** visually
     depicted on the display unit...

...ADVANTAGE - Provides full **SQL** functionality by allowing
     user to construct search **strings** with **logical
     relationships** between search parameters and create complex
     SQL commands by **linking** simpler **SQL** commands with
     logical conditional **operators**. Is implemented by one or
     more application programs interacting with **user** via
     graphic user interface under the control of the

operating system...
...Title Terms: **JOIN;**
International Patent Class (Main): **G06F-017/30**
Manual Codes (EPI/S-X): **T01-J05B4M...**

...T01-J05B4P...

...T01-J12

Derwent WPIX

JOE APPLICANT

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013958122     **Image available**
WPI Acc No: 2001-442336/200148
XRPX Acc No: N01-327204
```

 . **Database management system for cross-node sharing of
instructions, especially cached dynamic SQL statements in a
multiple relational database management system environment**

Patent Assignee: IBM CANADA LTD (IBMC   ); BIRD P M (BIRD-I); SNOWBELL M J
 (SNOW-I); INT BUSINESS MACHINES CORP (IBMC   )
Inventor: BIRD P M; SNOWBELL M J
Number of Countries: 002   Number of Patents: 004
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
CA 2282999	A1	20010322	CA 2282999	A	19990922	200148	B
US 20020123978	A1	20020905	US 99404438	A	19990922	200269	N
CA 2282999	C	20030527	CA 2282999	A	19990922	200336	
US 6598058	B2	20030722	US 99404438	A	19990922	200354	N

Priority Applications (No Type Date): CA 2282999 A 19990922; US 99404438 A
  19990922
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
CA 2282999	A1	E	62	G06F-017/30	
US 20020123978	A1			G06F-017/30	
CA 2282999	C	E		G06F-017/30	
US 6598058	B2			G06F-017/30	

-**Database management system for cross-node sharing of
instructions, especially cached dynamic SQL statements in a
-multiple relational database management system environment**

Abstract (Basic):
... and tables residing in a data partition. One of the nodes has a
   catalogue containing **metadata** describing tables and other
   objects and **relationships** between them in the data
   processing system. A global instruction cache at each node has a
   statement portion storing instruction **statements** and
   associated **executables**, each of which is designed for a
   particular environment. Instruction cache has dependency portion, which
   lists **objects** and **links** them to **corresponding**
   executable entries that depend on the **objects**. Shipped
   variation portion of the instruction cache has entries each referring
   to one **executable** entry and identifies the node from
   statement associated with **executable** was originally
   shipped.

For **relational** database management systems, particularly for efficiently sharing instructions, especially cached dynamic **SQL statements** (and their sections) between applications connected at different nodes in a multiple node database while...

...Since identification of dynamic **SQL** entry in the shipping process is based on original node information, which **corresponds** to identification of entry at node the request was originally compiled and shipped from, compiled...

...Title Terms: **SQL;**

APPLICANT

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0I6559228      **Image available**
WPI Acc No: 2004-717968/200470
XRPX Acc No: N04-569119
```

Data access providing method involves providing structured query language specification comprising logical fields for specifying criteria for data selection and specification of fields to be returned

```
Patent Assignee: INT BUSINESS MACHINES CORP (IBMC  )
Inventor: DETTINGER R D; LAROCCA J L; STEVENS R J
Number of Countries: 001   Number of Patents: 001
Patent Family:
Patent No      Kind    Date     Applicat No    Kind    Date       Week
US 20040193568 A1   20040930   US 2003401293   A    20030327   200470  B

Priority Applications (No Type Date): US 2003401293 A 20030327
Patent Details:
Patent No   Kind Lan Pg   Main IPC      Filing Notes
US 20040193568 A1      17 G06F-007/00
```

Data access providing method involves providing structured query language specification comprising logical fields for specifying criteria for data selection and specification of fields to be returned

```
Abstract (Basic):   -
```
...   A structured **query language (SQL)** specification comprising **logical fields** for specifying criteria for data selection and specification of **fields** to be returned, is provided for one of the applications. The mapping rules comprising access method, map **logical fields** to particular entity in **physical** data representation.

...   For providing access to data having particular **physical** data representation...

...specification and the underlying data representation. The application defines data query requirement in a mode **abstract** manner, that are bound particular **physical** data representation at runtime...

...Title Terms: **FIELD;**

46/3,K/8 (Item 8 from file: 350) **Links**
Derwent WPIX

APPLICANT

017596935   **Image available**
WP.I Acc No: 2006-108190/200611
Related WPI Acc No: 2003-778362; 2003-778410; 2005-294869
XRPX Acc No: N06-093925
  Data access method for providing access to data having
  particular schema, involves accessing data abstraction
  model having logical field name that is different
  from physical field name of corresponding
  physical data as defined by schema
Patent Assignee: INT. BUSINESS MACHINES CORP (IBMC)
Inventor: DETTINGER R D; JOHNSON P J; STEVENS R J; TONG I; WILL E
Number of Countries: 001 Number of Patents: 001
Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week	
US 20060010127	A1	20060112	US 200283075	A	20020226	200611	B
			US 2005226181	A	20050914		

Priority Applications (No Type Date): US 200283075 A 20020226; US
  2005226181 A 20050914
Patent Details:

Patent No	Kind	Lan	Pg	Main IPC	Filing Notes
US 20060010127	A1		14	G06F-017/30	Cont of application US 200283075

  Data access method for providing access to data having
  particular schema, involves accessing data abstraction
  model having logical field name that is different
  from physical field name of corresponding
  physical data as defined by schema

Abstract (Basic):
...   A data **abstraction model** with **logical**
    field definitions for **modeling** the data having particular
    schema e.g. **SQL** query, **XML query**, and exposing the
    data to users, is accessed in response to an **abstract** query.
    Each of the definitions comprises a **logical field** name
    that is different from the **physical field** name of the
    corresponding physical data as defined by the schema.
...   Enables to access data independent of particular manner in which
    the data is **physically** represented...
...Title Terms: **ABSTRACT;**